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(54) **SOFT ERROR AND TRANSIENT ERROR DETECTION DEVICE AND METHODS THEREFOR**

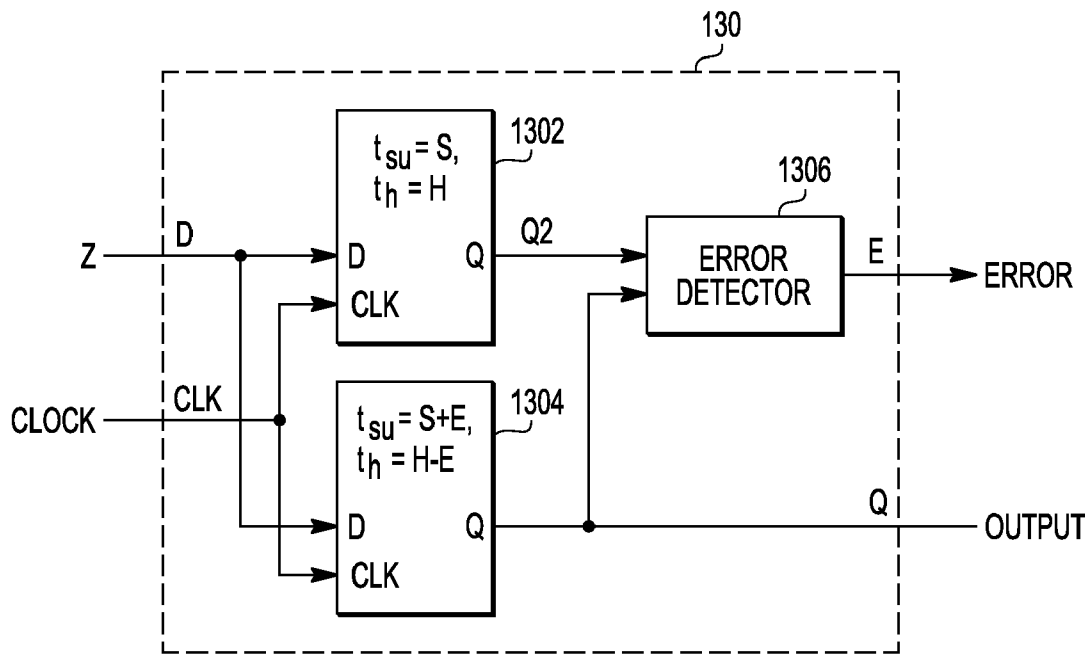(75) Inventors: **William C. Moyer**, Dripping Springs, TX (US); **Troy L. Cooper**, Austin, TX (US)

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(57) **ABSTRACT**

A clock signal is received at a clock node of a latch module, and a data signal is received at a data node of the latch module. The data signal including information to be latched at a first latch of the latch module and at a second latch of the latch module. A first representation of the data signal to a first data node of the first latch is delayed relative to a second representation of the data signal to a corresponding first data node of the second latch to obtain a first timing requirement between the data signal and the clock signal relative to the first latch that is substantially different than a second timing requirement. An error signal is generated in response to different data being latched at the first latch than at the second latch.

130

130

105

110   120   130

INPUT

D   Q   A

LOGIC GATES

Z   D

E — ERROR

Q — OUTPUT

CLK

CLK

CLOCK

100

FIG. 1

130

$t_{su} = S,$
$t_h = H$ 1302

Z   D

D   Q   Q2

ERROR DETECTOR   1306

E   ERROR

CLK

CLOCK   CLK

$t_{su} = S+E,$
$t_h = H-E$ 1304

D   Q

Q   OUTPUT

CLK

130

FIG. 2

FIG. 3

FIG. 4

FIG. 5

FIG. 6

FIG. 7

SOFT ERROR AND TRANSIENT ERROR DETECTION DEVICE AND METHODS THEREFOR

BACKGROUND

[0001]   1. Field of the Disclosure

[0002]   The present disclosure relates generally to data processing devices, and more particularly, to error detection at data processing devices.

[0003]   2. Description of the Related Art

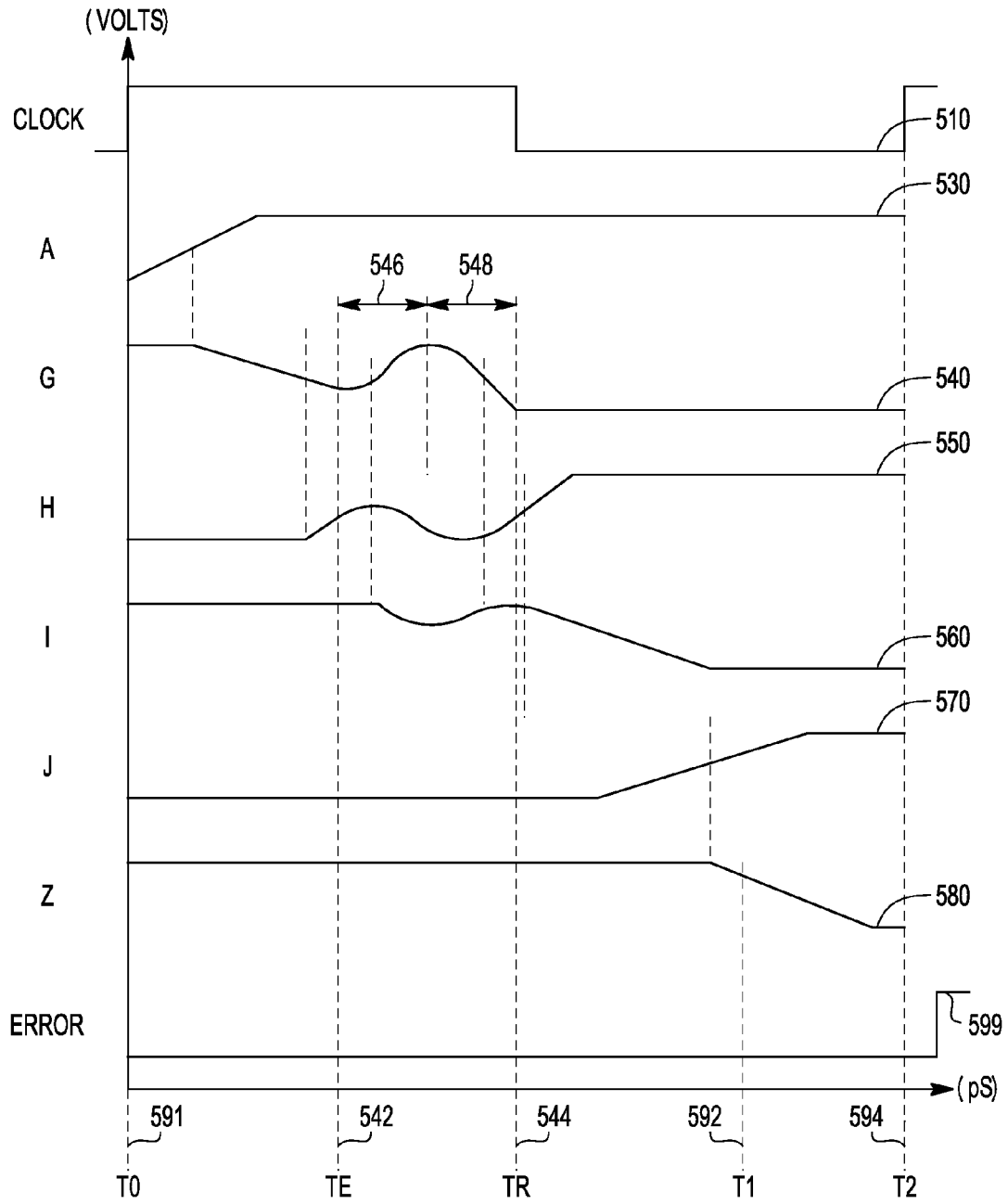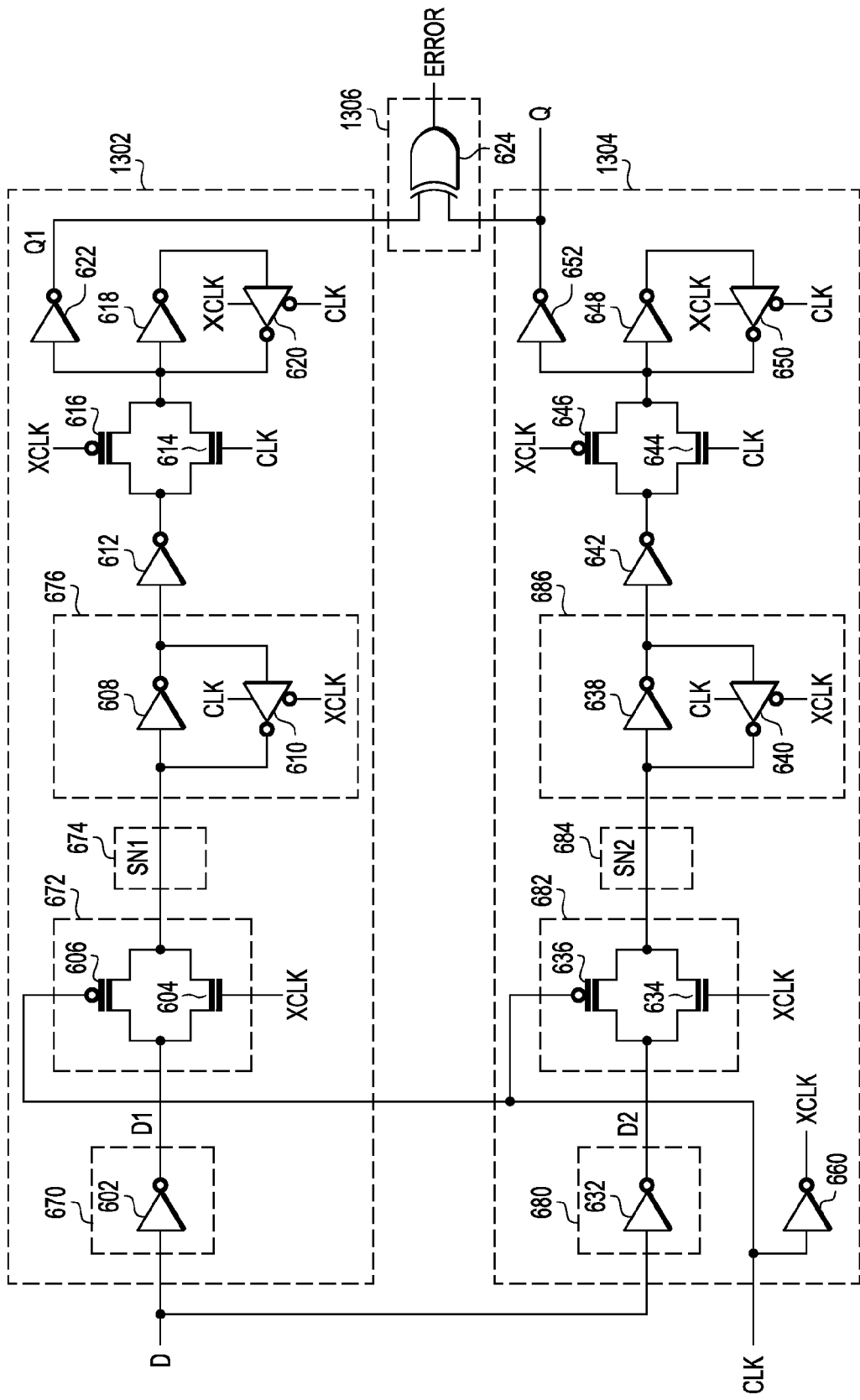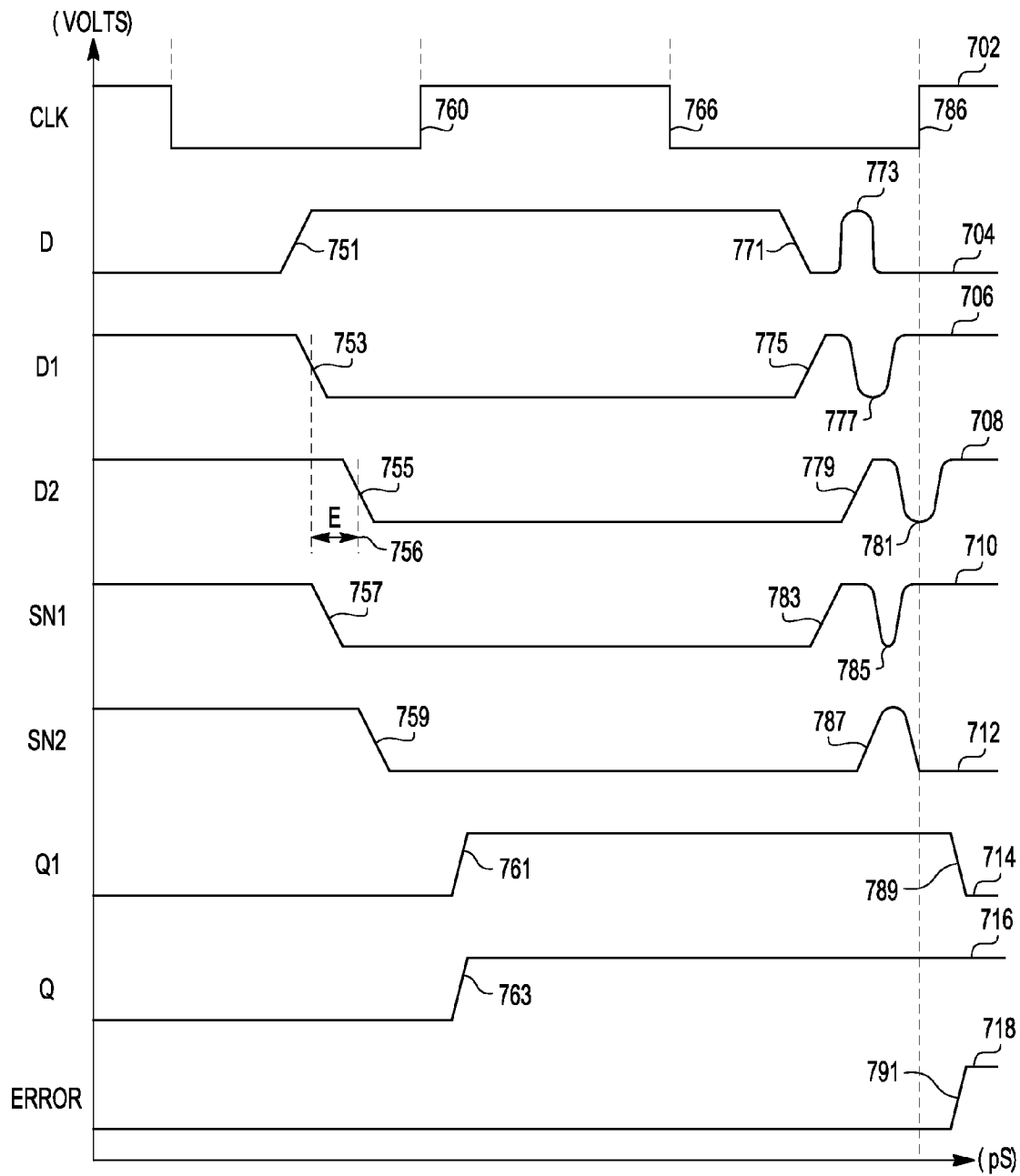[0004]   Advancements in the performance and computational capabilities of integrated circuit (IC) data processing devices are often the result of a corresponding reduction in the size of the transistors, conductors, and other components that make up the devices. Furthermore, the operating voltage of modern IC data processing devices is typically minimized to increase the operating frequency and to reduce power dissipation of the devices. Because of these changes, conductors and storage nodes included at the data processing device have correspondingly low capacitance, and therefore are capable of storing only a small amount of energy. Thus, the device is increasingly susceptible to soft errors. A soft error can result from the interaction of an energetic neutron or alpha particle with the small conductors and charge storage nodes of a data processing device. Neutrons typically originate from cosmic rays, and alpha particles originate from materials present in the integrated circuit package and bonding materials.

[0005]   A combination of the capacitance and the voltage at a node of a device is described by a critical charge parameter, Qcrit, which is the minimum electron charge disturbance needed to change the logic level at the node. Larger Qcrit values are associated with fewer soft errors. Unfortunately, signals at a node with a larger Qcrit value will transition slower and dissipate more power than a node with a relatively smaller Qcrit value. Reduction in the feature size of transistor and conductors at the data processing device and a reduction in the operating voltage decreases Qcrit. Thus, the incidence of soft errors increases as chip technology advances. In a logic circuit, Qcrit is defined as the minimum amount of induced charge required at a circuit node to cause a voltage pulse to propagate from that node to the output of a corresponding logic gate, and be of sufficient duration and magnitude for an error to be subsequently stored at a down-stream latch device. Error correction codes and associated circuitry can be used to detect and sometimes correct soft errors, but is not practical for use at the large number of latches that make up computational logic blocks of a data processing device.

BRIEF DESCRIPTION OF THE DRAWINGS

[0006]   The present disclosure may be better understood, and its numerous features and advantages made apparent to those skilled in the art by referencing the accompanying drawings.

[0007]   FIG. 1 is a block diagram illustrating a data processing device including a latch capable of detecting a soft error in accordance with a specific embodiment of the present disclosure.

[0008]   FIG. 2 is a block diagram illustrating an error-detecting latch of FIG. 1 in accordance with a specific embodiment of the present disclosure.

[0009]   FIG. 3 is a schematic diagram illustrating logic gates and the error-detecting latch of FIG. 1 in accordance with a specific embodiment of the present disclosure.

[0010]   FIG. 4 is a timing diagram illustrating the operation of a synchronous logic circuit of FIG. 1 in accordance with a specific embodiment of the present disclosure.

[0011]   FIG. 5 is a timing diagram illustrating the operation of the synchronous logic circuit of FIG. 1 in response to a soft error event in accordance with a specific embodiment of the present disclosure.

[0012]   FIG. 6 is a schematic diagram illustrating an error-detecting latch, such as the error-detecting latch of FIG. 1 in accordance with a specific embodiment of the present disclosure.

[0013]   FIG. 7 is a timing diagram illustrating the operation of the error detecting latch of FIG. 6 in accordance with a specific embodiment of the present disclosure.

DETAILED DESCRIPTION

[0014]   The present disclosure provides a storage element, such as a latch, operable to detect a soft error or error caused by transient interference upstream from the latch, and to provide a signal indicative of the error. The storage element includes multiple latches that receive substantially the same clock signal, such that operation of the storage element is dependent only upon the propagation delay of a single external clock signal. Delays associated with the data paths of the multiple latches are different to facilitate different setup time and hold times at each latch to detect an error event, such as a soft error.

[0015]   FIG. 1 is a block diagram illustrating a data processing device 100 including a synchronous logic circuit 105 that includes latch 110, logic gates 120, and a storage element, also referred to as error-detecting latch 130. Latch 110 has an input D connected to a node labeled "INPUT," an input CLK connected to a node labeled "CLOCK," and an output Q connected to a node labeled "A." A logic gate 120 has an input connected node A and an output connected to a node labeled "Z." Latch 130 is an error-detecting latch capable of detecting a soft error in accordance with a specific embodiment of the present disclosure. Latch 130 has a data input D connected to node Z, a clock input CLK connected to node CLOCK, a data output Q connected to a node labeled "OUTPUT," and an error detect output E connected to a node labeled "ERROR."

[0016]   Logic gates 120 represent what is commonly referred to as "a cloud of logic" with one or more data inputs and one or more data outputs. A single data input, connected to node A, and a single data output, connected to node Z, are illustrated for clarity. A logic gate 120 includes logic gates configured to implement a desired logic function. A signal received at an input of logic gates 120, such as from node A, is dispatched in response to a particular transition of clock signal CLOCK, and the data at the outputs of logic gates 120 are latched on the next such clock transition at a latch, such as error-detecting latch 130.

[0017]   Error-detecting latch 130 is configured to indicate an error in response to storage nodes at error-detecting latch 130 latching different data. For purposes of illustration herein it is assumed that a soft error is being detected. A soft error can occur when a neutron, an alpha particle, or other ionizing radiation interacts with synchronous logic circuit 105 in a manner that causes the logic state of the storage node to be altered. For example, the interaction of such a particle can result in a temporary anomalous transition of a node at logic gates 120 that can propagate downstream through one or more logic gates and can result in an incorrect logic value being stored at a latch device. In either case, the particle or

radiation results in a soft error event, and the soft error event has an associated duration. Error-detecting latch **130** includes two or more latches, and data received at a storage node of each latch is latched in response to a common transition of the clock signal at the input of error-detecting latch **130**, such as in response to a clock signal transition from a low logic state to a high logic state. The propagation of the data signal from the data input of the error-detecting latch is varied to adjust the setup time and hold time of each latch so that a soft error event with a duration of less than a predetermined period is rejected by at least one of the latches, while the clock signal received at each latch of error detecting latch **130** at substantially the same time. Due to variances in conductors used to propagate signal CLOCK, such as asymmetries in the topology of conductors used to propagate signal CLOCK, a common transition of the clock signal may begin at one latch associated with error detecting latch **130** delayed with respect to when it begins at another latch of error detecting latch **130**. The arrival of the clock signal at each respective latch is considered to be substantially the same to the extent these variances are not intended to delay one clock transition from another. In a specific embodiment, the arrival of the clock signal at each respective latch is considered to be substantially the same when they differ by less than ten picoseconds, less than eight picoseconds, less than 6 picoseconds, or less than four picoseconds, and dependent on the process technology in which the device is fabricated, e.g. a forty-five nanometer technology.

[0018] FIG. 2 is a block diagram illustrating an error-detecting latch **130** of FIG. 1 in accordance with a specific embodiment of the present disclosure. Error-detecting latch **130** includes latch **1302**, latch **1304**, and error detector **1306**. Latch **1302** has a data input D connected to input D of error-detecting latch **130**, a clock input CLK connect to input CLK of error-detecting latch **130**, and an output Q connected to a node labeled "Q2." Latch **1304** has a data input D also connected to node D of error-detecting latch **130**, a clock input CLK also connected to node CLK of error-detecting latch **130**, and an output Q connected to a node labeled "Q." Error detector **1306** has an input connected to node Q, another input connected to node Q2, and an output connected to node E.

[0019] The setup time (t_{su}) between the data signal received at input D of error-detecting latch **130** and the clock signal received at input CLK of error-detecting latch **130** relative to latch **1302** is configured to be equal to S picoseconds (pS), and the hold time (t_h) is configured to be equal to be H pS. The setup time between the data signal received at input D of error-detecting latch **130** and the clock signal received at input CLK of error-detecting latch **130** relative to latch **1304** is configured to be equal to S+E pS, and the hold time (t_h) is configured to be equal to H−E pS, where the value E corresponds to a delay to facilitate detection of a soft error. For example, E can correspond to a portion of the approximate duration of an expected soft error event, or in the case of a soft error event propagated by logic gates, a portion of the increase in the effective propagation delay of a logic path due to an error event at a node associated with the logic path. The value of E can be selected to correspond to a soft error event with a specific maximum duration. The latching of data at latch **1302** and latch **1304** is initiated at the same time by clock signal CLOCK received at node CLK. The operation of error-detecting latch **130** is described in detail with regard to FIGS. 4-6.

[0020] FIG. 3 is a schematic diagram illustrating logic gates **120** and error-detecting latch **130** of FIG. 1 in accordance with a specific embodiment of the present disclosure. A logic gate **120** includes inverter **310**, NOR gate **312**, NAND gate **314**, inverter **316**, NAND gate **318**, and inverter **320**, which are referred to collectively as logic gates **310-320**. Inverter **310** has an input connected to node A and an output connected to node labeled "G." NOR gate **312** has an input connected to node G, an input connected to a node labeled "B," and an output connected to a node labeled "H." NAND gate **314** has an input connected to node H, an input connected to a node labeled "C," and an output connected to a node labeled "I." Inverter **316** has an input connected to node I, and an output connected to a node labeled "K." NAND gate **318** has an input connected to node J, an input connected to a node labeled "K," and an output connected to node Z. Inverter **320** has an input connected to a node labeled "D," and an output connected to node K. Error-detecting latch **130** has a data input connected to node Z, a clock input connected to node CLOCK, an error detect output connected to node ERROR, and a data output connected to node OUTPUT.

[0021] Logic gates **310-320** represent a generalized configuration of combinatorial logic gates that may precede a latch at a synchronous logic circuit. A signal received at node A can propagate downstream to nodes G, H, I, J, and Z when the logic path is enabled by appropriate logic values at nodes B, C, and D. Each respective logic gate has a corresponding propagation delay based predominately on the size of transistors included at that gate, and on the capacitance associated with the output node that the gate controls. In order for the synchronous logic circuit to function correctly, information must propagate to node Z before the information is latched at error-detecting latch **130**. Specifically, data provided from node Z to input D must be valid for a period of time equal to or greater than the setup time of error-detecting latch **130**. In this particular example, the setup time is specified relative to the rising transition of clock signal CLOCK an input CLK.

[0022] Error-detecting latch **130** is configured to store information received at input D in response to clock signal CLOCK, and to provide signal ERROR when a soft error results in invalid information stored at one of the latches included at error-detecting latch **130**. It is statistically unlikely for a soft error event to be of sufficient duration, or for two independent soft errors to occur during the same clock cycle, so that both latches included at error-detecting latch **130** stores the invalid information.

[0023] FIG. 4 is a timing diagram **400** illustrating the operation of synchronous logic circuit **105** of FIG. 1 in accordance with a specific embodiment of the present disclosure. Timing diagram **400** includes a horizontal axis representing time in picoseconds (pS), and a vertical axis representing voltage in volts. Timing diagram **400** includes waveforms **410**, **430**, **440**, **450**, **460**, **470**, **480**, and **499** representing signals at nodes associated with logic gates **120** at FIG. 3. Waveform **410** corresponds to clock signal CLOCK at node CLOCK. Waveform **430** represents a signal at node A, waveform **440** represents a signal at node G, waveform **450** represents a signal at node H, waveform **460** represents a signal at node I, waveform **470** represents a signal at node J, waveform **480** represents a signal at node Z, and waveform **499** represents a signal at node ERROR. Also illustrated are time reference T1 **492** and time reference T2 **494**.

[0024] Timing diagram **400** illustrates the operation of synchronous logic circuit **105** in the absence of a soft error event.

In response to a rising transition of clock signal CLOCK, latch **110** at FIG. **1** begins to drive node A to a logic-high level. Once node A reaches a suitable level, inverter **310** begins to drive node G to a logic-low level. Information continues to propagate downstream successively to nodes H, I, J, and ultimately node Z, which is driven to a logic-low level just prior to time reference T1. Time reference T1 **492** and time reference T2 **494** represent the setup times of the two latches included at error-detecting latch **130**, respectively. In the absence of a soft error event, valid information (a logic-low value) is successfully stored at both latches. Thus, signal ERROR is not asserted by error-detecting latch **130**, and waveform **499** remains at a logic-low level.

[0025] FIG. **5** is a timing diagram **500** illustrating the operation of synchronous logic circuit **105** of FIG. **1** in response to a soft error event in accordance with a specific embodiment of the present disclosure. Timing diagram **500** includes a horizontal axis representing time in picoseconds (pS), and a vertical axis representing voltage in volts. Timing diagram **500** includes waveform **510**, **530**, **540**, **550**, **560**, **570**, **580**, and **599** representing signals at nodes associated with logic gates **120** at FIG. **3**. Waveform **510** corresponds to clock signal CLOCK at node CLOCK. Waveform **530** represents a signal at node A, waveform **540** represents a signal at node G, waveform **550** represents a signal at node H, waveform **560** represents a signal at node I, waveform **570** represents a signal at node J, waveform **580** represents a signal at node Z, and waveform **599** represents a signal at node ERROR. Also illustrated is time reference T1 **592**, time reference T2 **594**, time reference TE **542**, time reference TR **544**, interval **546**, and interval **548**.

[0026] In response to a rising transition of clock signal CLOCK at T0, latch **110** at FIG. **1** begins to drive node A to a logic-high level. Once node A reaches a suitable level, inverter **310** begins to drive node G to a logic-low level. At approximately time reference TE, a soft error occurs at node G. The soft error has a duration that is approximately equal to interval **544**, during which time spurious charge is applied to node G causing the voltage at node G to rise to a logic-high level, for example. It will be appreciated that a soft error need not fully discharge or charge a node to an anomalous level for the soft error to result in invalid information being stored at a downstream latch. Merely delaying the propagation of valid information can cause an error at synchronous logic circuit **105** to occur. During interval **546**, inverter **310** continues to drive node G low, which recovers to the logic-low level at time reference TR **544**.

[0027] NOR gate **312** responds to the anomalous signal at its input connected to node G, which is reflected at its output at node H, and to a lesser extent at node I. The affect of the soft error event is not immediately evident by the shape of waveform **570**, node J, or waveform **580**, node Z, but a final signal transition at these nodes is delayed relative to waveforms **470** and **480**, respectively, at FIG. **4**, as a result of the soft error event. Furthermore, signal **580** is at a logic-high level at time reference T1 **592**, and at a logic-low level at time reference T2 **594**. Thus, the two latches included at error-detecting latch **130** contain conflicting information, resulting in the assertion of signal ERROR at error-detecting latch **130**, as illustrated by waveform **599**. Timing diagram **500** illustrates a soft error event associated with node G. It will be appreciated that error detecting latch **130** is operable to detect a soft error event at any node associated with logic cloud **120**, such as nodes A, B, C, D, G, H, I, J, K, and Z.

[0028] FIG. **6** is a schematic diagram illustrating an error-detecting latch **600**, such as error-detecting latch **130** of FIG. **1**, in accordance with a specific embodiment of the present disclosure. Error-detecting latch **600** is a latch module that includes latch **1302**, latch **1304**, and error detector **1306**. Error-detecting latch **600** has a data input connected to node D, a clock input connected to node CLK, a data output connected to node Q, and an error output connected to node ERROR. A data value at data output node Q is determined based upon a state of latch **1304**, and an error indicator at error output node ERROR is based upon the state of latch **1304** and a state of latch **1302**.

[0029] Latch **1302** includes inverters **602**, **608**, **612**, **618**, and **622**, tri-state inverters **610** and **620**, n-channel metal oxide semiconductor (NMOS) field effect transistors **604** and **614**, and p-channel metal oxide semiconductor (PMOS) field effect transistors **606** and **616**. Inverter **602** has an input connected to input node D to receive a data signal and an output connected to a data node labeled "D1." NMOS transistor **604** has a first current electrode connected to node D1, a control electrode connected to a node labeled "XCLK," and a second current electrode connected to a data node labeled "SN1." PMOS transistor **606** has a first current electrode connected to node D1, a control electrode connected to node CLK, and a second current electrode connected to node SN1. NMOS transistor **604** and PMOS transistor **606** together implement a pass gate. Inverter **608** has an input connected to node SN1 and an output connected to a data node that includes the input of inverter **612** and the input of tri-state inverter **610**. Tri-state inverter **610** has active-high drive enable input connected to node CLK, and an active-low drive enable input connected to node XCLK, and an output connected to node SN1.

[0030] Inverter **612** has an output connected to a data node that includes first current electrode of NMOS transistor **614** and a first current electrode of PMOS transistor **616**. NMOS transistor **614** has a control electrode connected to clock node CLK, and a second current electrode connected to the data node that includes the input of inverter **618** and inverter **622**, and the output of tri-state inverter **620**. PMOS transistor **616** has a control electrode connected to node XCLK, and a second current electrode connected to the input of inverter **618** and inverter **622**, and the output of tri-state inverter **620**. Inverter **618** has an output connected to the input of tri-state inverter **620**. Tri-state inverter **620** has active-high drive enable input connected to node XCLK, and an active-low drive enable input connected to node CLK. Inverter **622** has an output connected to a node labeled "Q1," and to an input of an exclusive-OR gate **624** of error detector **1306**. Inverter **660** has in input connected to node CLK, and an output connected to node XCLK. Exclusive-OR gate **624** has an output connected to node ERROR.

[0031] Latch **1304** includes inverters **632**, **638**, **642**, **648**, and **652**, tri-state inverters **640** and **650**, n-channel metal oxide semiconductor (NMOS) field effect transistors **634** and **644**, and p-channel metal oxide semiconductor (PMOS) field effect transistors **636** and **646**. Inverter **632** has an input connected to input node D to receive a data signal and an output connected to a data node labeled "D2." NMOS transistor **634** has a first current electrode connected to node D2, a control electrode connected to node XCLK, and a second current electrode connected to a data node labeled "SN2." PMOS transistor **636** has a first current electrode connected to node D2, a control electrode connected to node CLK, and a

second current electrode connected to node SN2. Inverter **638** has an input connected to node SN2 and an output connected to the input of inverter **642** and to the input of tri-state inverter **640**. Tri-state inverter **640** has active-high drive enable input connected to node CLK, and an active-low drive enable input connected to node XCLK, and an output connected to node SN2.

[0032] Inverter **642** has an output connected to a data node than includes a first current electrode (source/drain) of NMOS transistor **644** and a first current electrode of PMOS transistor **646**. NMOS transistor **644** has a control electrode connected to clock node CLK, and a second current electrode connected to the data node that includes the input of inverter **648** and inverter **652**, and the output of tri-state inverter **650**. PMOS transistor **646** has a control electrode connected to node XCLK, and a second current electrode connected to the input of inverter **648** and inverter **652**, and the output of tri-state inverter **650**. Inverter **648** has an output connected to the input of tri-state inverter **650**. Tri-state inverter **650** has active-high drive enable input connected to node XCLK, and an active-low drive enable input connected to node CLK. Inverter **652** has an output connected to node Q, and to another input of exclusive-OR gate **624** of error detector **1306**.

[0033] Latch **1302** is configured to store information received at node D and provide the stored information at node Q1. Latch **1302** is referred to as a master-slave latch, and includes a master latch provided by inverter **608** and tri-state inverter **610**, and a slave latch provided by inverter **618** and tri-state inverter **620**. The master latch is updated while clock signal CLK is negated. The input of the slave latch is connected to the output of the master latch and is updated while clock signal CLK is asserted.

[0034] A data signal received at node D is buffered by inverter **602**, and is conducted to storage node SN1 by NMOS transistor **604** and PMOS transistor **606** while clock signal XCLK is asserted. The output of tri-state inverter **610** is maintained at a high-impedance state while clock signal XCLK is asserted. When clock signal XCLK is negated (and clock signal CLK is asserted), storage node SN1 is isolated from the output of inverter **602**, and the output of tri-state inverter **610** is enabled to drive storage node SN1. Also in response to the assertion of clock signal CLK, the slave latch is updated and the stored information is provided at node Q1.

[0035] The operation of latch **1304** is the same as described with reference to latch **1302**. Node Q at the output of latch **1304** is the primary data output of error-detecting latch **130**. Exclusive-OR gate **624** compares the information stored at latch **1302** and latch **1304**, and asserts signal ERROR if the value stored at each of the two latches is not the same, such as described with reference to the timing diagram of FIG. **5**. The setup time and hold time of the master portion of latch **1302** is configured to be different than corresponding timing requirements of the master portion of latch **1304**, as described with reference to FIG. **2**. For example, in an embodiment, the setup time associated with latch **1304** can be configured to be greater than the setup time associated with latch **1302** by an amount greater than approximately one half of the duration of an anticipated soft error event, such as approximately one-half of a soft error event.

[0036] The timing requirements, such as the setup time and hold times in particular, of the master latches of error-detecting latch **1302** can be set by a combination of circuit attributes that affect the delay of data along the data paths that include

the master latches. For example attributes associated with respective portions **670**, **672**, **674** and **676** of latch **1302**, and portions **680**, **682**, **684**, and **686** of latch **1304**. Portion **670** includes inverter **602**, portion **672** include NMOS transistor **604** and PMOS transistor **606**, portion **674** includes storage node SN1, and portion **676** includes inverter **608** and tri-state inverter **610**. Portion **680** includes inverter **632**, portion **682** include NMOS transistor **634** and PMOS transistor **636**, portion **684** includes storage node SN2, and portion **686** includes inverter **638** and tri-state inverter **640**. The setup time and hold time associated with the master latches can be configured by adding delay elements, adjusting the widths and lengths of individual transistors, adjusting the threshold voltage of these transistors, and adjusting the capacitance at nodes associated with these transistors.

[0037] The setup time and hold time of a latch are examples of timing requirements that relate the transition of a data signal received at the input D of the latch **600** to a transition of clock signal at input CLK. A setup timing requirement specifies a minimum time that a data signal must be at a valid level prior to an active transition of signal CLK. A hold timing requirement specifies a minimum time that a data signal must remain at the valid level following the active transition of signal CLK. Setup and hold times can be positive or negative values. The setup and hold times of latch **1302** (e.g., the timing requirement between the data signal at input D and the clock signal at input CLK for latch **1302**) and latch **1304** (e.g., the timing requirement between the data signal at input D and the clock signal at input CLK for latch **1304**) are configured to be substantially different. For example, the data signal along the data path of latch **1302** can be adjusted, relative to the data signal along the data path of latch **1304**, by an amount that is greater than an expected variation between the clock paths of latch **1302** and **1304**. For example, the delay of a data signal can be greater than sixteen picoseconds. In contrast to the master portions of the latches **1302** and **1304**, the slave portion of each latch is designed to be substantially similar. For example, inverters **612**, **618**, **622**, NMOS transistor **614**, PMOS transistor **616**, and tri-state inverter **620** of latch **1302** and their interconnects can be designed to have substantially the same electrical properties as the corresponding devices at latch **1304**. Thus, the slave portion of latch **1302** and the slave portion of latch **1304** have substantially the same timing characteristics, and which are not dependent on external clock propagation delays.

[0038] Previous techniques utilized to implement an error-detecting latch require manipulation of the propagation delay of disparate external clock signals provided to the error-detecting latch using complex clock-tree and clock distribution analysis. Additional devices and node capacitance associated with multiple clock signals can increase the amount of power dissipated by the data processing device. The disclosed error-detecting latch **130** receives only one clock signal (signal CLK) and thus eliminates the need to precisely regulate the delay of multiple clock signals. This technique has the additional benefit that the timing characteristics of error detecting latch **130** are not dependent on multiple external clock propagation delays. Thus, the timing properties of error detecting latch **130** can be characterized independent of its instantiations at a device.

[0039] FIG. **7** is a timing diagram **700** illustrating the operation of error detecting latch **600** of FIG. **6** in accordance with a specific embodiment of the present disclosure. Timing diagram **700** includes a horizontal axis representing time in

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picoseconds (pS), and a vertical axis representing voltage in volts. Timing diagram **700** includes waveforms **702**, **704**, **706**, **708**, **710**, **712**, **714**, **716**, and **718** representing signals at nodes associated with devices included at error detecting latch **130** at FIG. **6**. Waveform **702** corresponds to a clock signal at input node CLK and includes transition **760**, **766**, and **786**. Waveform **704** represents a signal at input node D and includes transitions **751** and **771** and transient pulse **773**. Waveform **706** represents a signal at node D1 based upon the signal at D1, e.g., a representation of the signal at input node D, and includes transitions **753** and **775**, and transient pulse **777**. Waveform **708** represents a signal at node D2 based on the signal at D, e.g., a representation of the signal at input node D, and includes transitions **755** and **779**, interval E **756**, and transient pulse **781**. Waveform **710** represents a signal at node SN1 and includes transitions **757** and **783**, and transient pulse **785**. Waveform **712** represents a signal at node SN2 and includes transitions **759** and **787**. Waveform **714** represents a signal at node Q1 and includes transitions **761** and **789**. Waveform **716** represents a signal at node Q and includes transition **763**. Waveform **718** represents a signal at node ERROR and includes transition **791**.

[0040] The setup time of latch **1302** is configured to be different than the setup time of latch **1304** by configuring the propagation delay of inverter **602** and inverter **632** to differ by a desired amount of time, E **756**. For example, inverter **632** can include three inverters connected in series, while inverter **602** can include a single inverter. A low to high transition **751** of data input signal D is received at the input of inverter **602** and inverter **632**, resulting in corresponding high to low transitions **753** and **755** at nodes D1 and D2 respectively. The propagation delay of inverter **632** is configured to be greater than the propagation delay of inverter **602** by an amount of time illustrated by an interval labeled "E." At the time that the nodes D1 and D2 transitions low, clock signal CLK is negated so that the pass gate implemented by PMOS transistor **606** and NMOS transistor **604** and the pass gate implemented by PMOS transistor **636** and NMOS transistor **634** are both activated, resulting in transition **757** at node SN1 and transition **759** at node SN2. A subsequent low to high transition **760** at node CLK deactivates both of the pass gates associated with NMOS transistor **604** and NMOS transistor **634** to latch identical logic values at latch **1302** and latch **1304**. The low to high transition **760** causes the pass gate implemented by PMOS transistor **616** and NMOS transistor **614** and the pass gate implemented by PMOS transistor **646** and NMOS transistor **644** to be activated, and the information stored at latch **1302** and latch **1304** to be propagated to node Q1 (transition **761**) and node Q (transition **763**), respectively. Because the same information was successfully latched in latch **1302** and latch **1304**, signal ERROR is not asserted.

[0041] During the next clock cycle, a high to low transition **771** of data input signal D is received at the input of inverter **602** and inverter **632**, resulting in corresponding high to low transitions **775** and **779** at nodes D1 and D2 respectively. At this time, clock signal CLK is negated so that the pass gate implemented by PMOS transistor **606** and NMOS transistor **604** and the pass gate implemented by PMOS transistor **636** and NMOS transistor **634** are both activated, resulting in transition **783** at node SN1 and transition **787** at node SN2. Shortly after transition **771** of data input signal D, a soft error at node D causes transient pulse **773**. Transient pulse **773** is propagated through inverter **602** and inverter **632**, resulting in transient pulse **777** at node D1 and transient pulse **781** at node

D2. The pass gate implemented by PMOS transistor **606** and NMOS transistor **604** is activated at the time the transient pulse occurs at input D, and the signal received at node D1, including its subsequent recovery to a logic-high level, is propagated to node SN1. The logic-high level is successfully stored at the master latch of latch **1302** in response to transition **786** of signal CLK, resulting in a high to low transition **789** at node Q1. However, because the signal provided to nodes D2 and SN2 are delayed with respect to nodes D1 and SN1, the signal at node SN2 is unable to recover to a logic-high level before the pass gate implemented by PMOS transistor **636** and NMOS transistor **634** is deactivated by the signal CLK. Thus, a logic-low level is latched at the master latch of latch **1304**, and signal Q remains at a logic-high level. Signal ERROR is asserted to indicate an error condition because latch **1302** and latch **1304** have stored different values.

[0042] In an alternate embodiment, the timing requirements of latch **1302** and latch **1304** can be adjusted by manipulating properties associated with other devices along their respective delay paths. For example, increasing the threshold voltage of NMOS transistor **604** and PMOS transistor **606** relative to the threshold voltage of NMOS transistor **634** and PMOS transistor **636** changes the I/V characteristics of these devices and thus modifies the switching characteristics of inverter **608** and **638**, respectively, to introduce and effective delay. In another embodiment, the capacitance at storage node SN1 can be increased relative the capacitance at storage node SN2 to increase the propagation delay of data from input D to node SN1 of the data path through latch **1302**, and thereby reducing the setup time of the master latch. In yet another embodiment, the switching characteristics of inverter **608** or inverter **638** can be manipulated by adding delay elements, adjusting transistor dimensions, doping levels, or the like, to modify the timing requirements of latch **1302** and latch **1304**.

[0043] Note that not all of the activities or elements described above in the general description are required, that a portion of a specific activity or device may not be required, and that one or more further activities may be performed, or elements included, in addition to those described. Still further, the order in which the activities are listed are not necessarily the order in which they are performed.

[0044] Also, the concepts have been described with reference to specific embodiments. However, one of ordinary skill in the art appreciates that various modifications and changes can be made without departing from the scope of the present disclosure as set forth in the claims below. Accordingly, the specification and figures are to be regarded in an illustrative rather than a restrictive sense, and all such modifications are intended to be included within the scope of the present disclosure. For example, in an embodiment, error-detecting latch **130** can be implemented using standard cells available in an existing standard cell library. Latch **1302** and latch **1304** can be selected from existing master-slave latches based on their timing characteristics and based on the transition time of the corresponding data input signal. Error-detecting latch **600** described at FIG. **6** represents one example of how error-detecting latch **130** can be implemented. Latch **1302** and latch **1304** can be implemented using another type of latch circuit that provides the disclosed timing characteristics. Furthermore, the term latch is used broadly in the disclosure and is applicable to various types of storage elements that store information based upon a clock signal. Types of latches

include level-sensitive latch, edge-triggered latch, or flip-flop such as a JK flip-flop or a D flip-flop.

[0045] An error detecting latch, such as error detecting latch **130**, can include more than two latches, and majority-voting logic can be used to determine the value of data provided at the output of error detecting latch **130** based on respective values stored at each latch. Signal ERROR is asserted if all latches included at error detecting latch **130** have not stored the same value. Furthermore, three latches can be included at error detecting latch **130** and each latch can include different setup and hold times relative to the other latches.

[0046] Benefits, other advantages, and solutions to problems have been described above with regard to specific embodiments. However, the benefits, advantages, solutions to problems, and any feature(s) that may cause any benefit, advantage, or solution to occur or become more pronounced are not to be construed as a critical, required, or essential feature of any or all the claims.

What is claimed is:

1. A method comprising:

receiving a clock signal at a clock node of a latch module; and

receiving, at a data node of the latch module, a data signal including information to be latched at a first latch of the latch module and at a second latch of the latch module; and

delaying a first representation of the data signal to a first data node of the first latch relative to a second representation of the data signal to a corresponding first data node of the second latch to obtain a first timing requirement between the data signal and the clock signal relative to the first latch that is substantially different than a second timing requirement between the data signal and the clock signal relative to the second latch, wherein the timing requirement is at least one of a setup time and a hold time.

2. The method of claim **1** further comprising:

determining a data value at a data output of the latch module based upon a state of the first latch.

3. The method of claim **2**, wherein the data value at the data output of the latch module is based only upon the state of the first latch.

4. The method of claim **1** further comprising:

determining a data value at a data output of the latch module based upon a state of the first latch and a state of the second latch.

5. The method of claim **1** further comprising:

determining an error indicator at an output of the latch module indicative of a data error at the latch module, the error indicator based upon a state of the first latch and a state of the second latch.

6. The method of claim **5** further comprising:

determining a data value at a data output of the latch module based upon a state of the first latch.

7. The method of claim **5**, wherein the data value at the data output of the latch module is based only upon the state of the first latch.

8. The method of claim **5** further comprising:

determining a data value at a data output of the latch module based upon a state of the first latch and a state of the second latch.

9. The method of claim **1**, wherein the first data node of the first latch includes a first current electrode of a first pass gate,

the first current electrode of between a second current electrode of the first pass gate and the data node of the latch module.

10. The method of claim **1**, wherein delaying further comprises delaying the first representation of the data signal between a first current electrode of a first pass gate and a second current electrode of the first pass gate, the first current electrode of the first pass.

11. A device comprising:

a data input to receive a data signal;

a clock input to receive a clock signal to latch information based on the data signal;

a first signal path including the data input and a first latch, the first signal path having a first delay between the data input and a first data node for a given transition at the data input; and

a second signal path including the data input and a second latch, the second signal path including a first data node at a same location relative to the second latch as the first data node of the first signal path is located relative to the first latch, the second signal path having a second delay between the data input and the first node of the second signal path for the given transition, the second delay being substantially different that the first delay.

12. The device of claim **1** further comprising a data output coupled to the first latch to provide a latched data value.

13. The device of claim **1** wherein the latched data value is based on the first latch but not the second latch.

14. The device of claim **13** further comprising an error output coupled to the first latch and to the second latch to provide an indicator that the latched data value is to be considered invalid.

15. The device of claim **11**, wherein the first signal path further comprises

a first delay element comprising a first input coupled to the data input, and an output;

a first pass gate comprising a first current electrode coupled to the output of the first delay element, and a second current electrode; and

wherein the second signal path further comprises

a second delay element comprising a first input coupled to the data input, and an output, the second delay element having a substantially larger delay than the first delay element;

a second pass gate comprising a first current electrode coupled to the output of the first delay element, and a second current electrode.

16. The device of claim **11**, wherein a delay through the second latch for the given transition is larger than the delay through the first latch for the given transition.

17. The device of claim **11**, wherein the first signal path further comprises

a first pass gate comprising a first current electrode coupled to receive information from the data input, and a second current electrode coupled to provide the information the first latch, the first pass gate having a third delay between the first current electrode and the second current electrode for the given signal; and

a second pass gate comprising a first current electrode coupled to receive information from the data input, and a second current electrode coupled to provide information to the second latch, the second pass gate having a fourth delay between the first current electrode and the second current electrode for the given signal, the differ-

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ence between the first delay and the second delay is substantially the same as the difference between the third delay and the third delay.

18. A method comprising:

receiving a signal transition at a data node of a latch module;

providing a first representation of the signal transition to a source/drain of a first transistor at a first time;

providing a second representation of the signal transition to a source/drain of a second transistor at a second time;

receiving, at a third time, an initial portion of the clock transition at a gate of a first transistor to latch information at a first storage node of the latch module; and

receiving, at substantially the third time, the initial portion of the clock transition at a gate of a second transistor to latch information at a second storage node of the latch module.

19. The method of claim 18, wherein providing the second representation includes providing the second representation of the signal transition to the source/drain of the second transistor in parallel with providing the first representation of the signal transition to the source/drain of the first transistor.

20. The method of claim 20 further comprising indicating an error condition based upon information latched at the first storage node and at the second storage node.

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