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(54) **OFFLOADING FUSED KERNEL EXECUTION TO A GRAPHICS PROCESSOR**

(71) Applicant: **Intel Corporation**, Santa Clara, CA (US)

(72) Inventors: **Konstantin S. Bobrovsky**, Lozhok (RU); **Sergey N. Ermolaev**, Berdsk (RU); **Serguei N. Dmitriev**, Novosibirsk (RU); **Knud J. Kirkegaard**, Santa Clara, CA (US)

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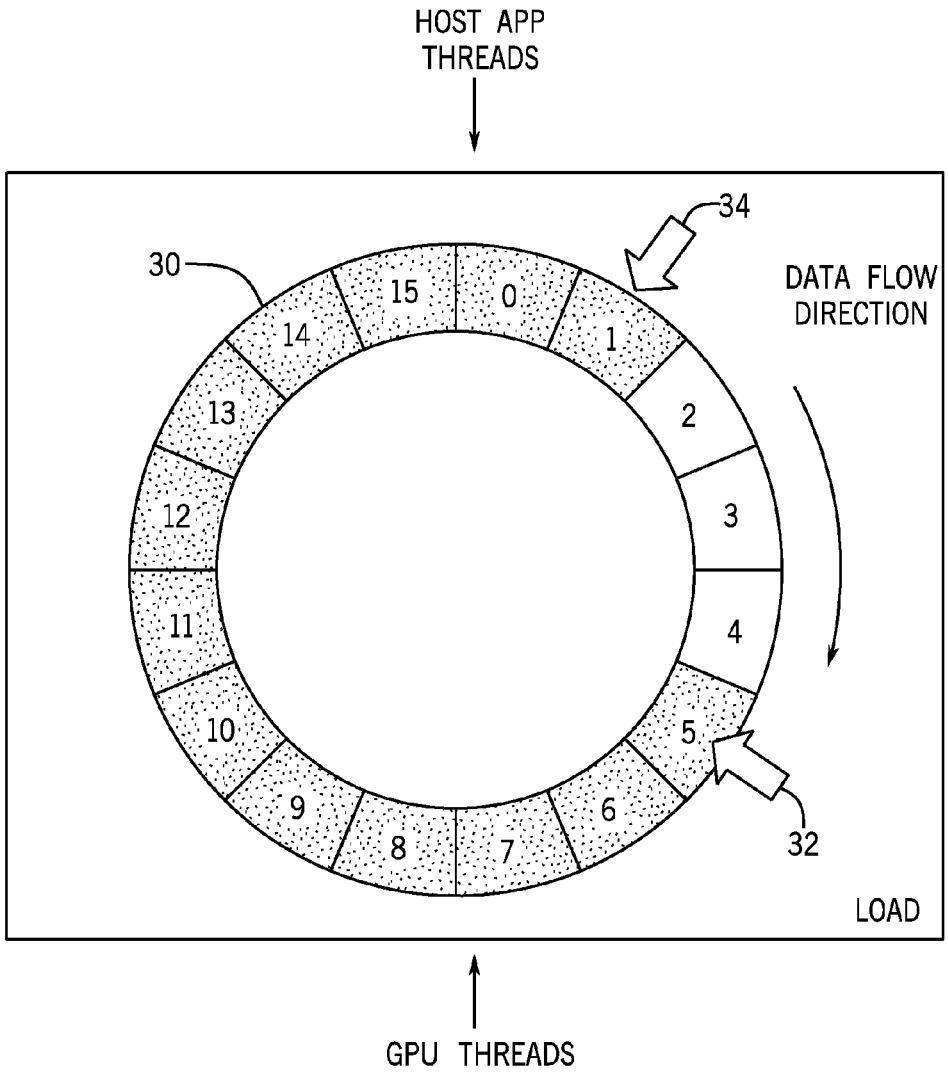
CPC . ***G06T 1/20*** (2013.01); ***G06T 1/60*** (2013.01)

(57) **ABSTRACT**

Execution of a first kernel may be offloaded from a central processing unit to a graphics processing unit using a ring task buffer with a fixed number of task slots incurring full overhead of runtime driver interaction. Execution of a second kernel is offloaded using said ring task buffer, so at least two kernels may be offloaded from a central processing unit to a graphics processing unit via said ring task buffer, while incurring about the same offloading overhead as would be incurred from offloading a single kernel, in some embodiments. Multiple kernels are automatically grouped together by a compiler and linker.



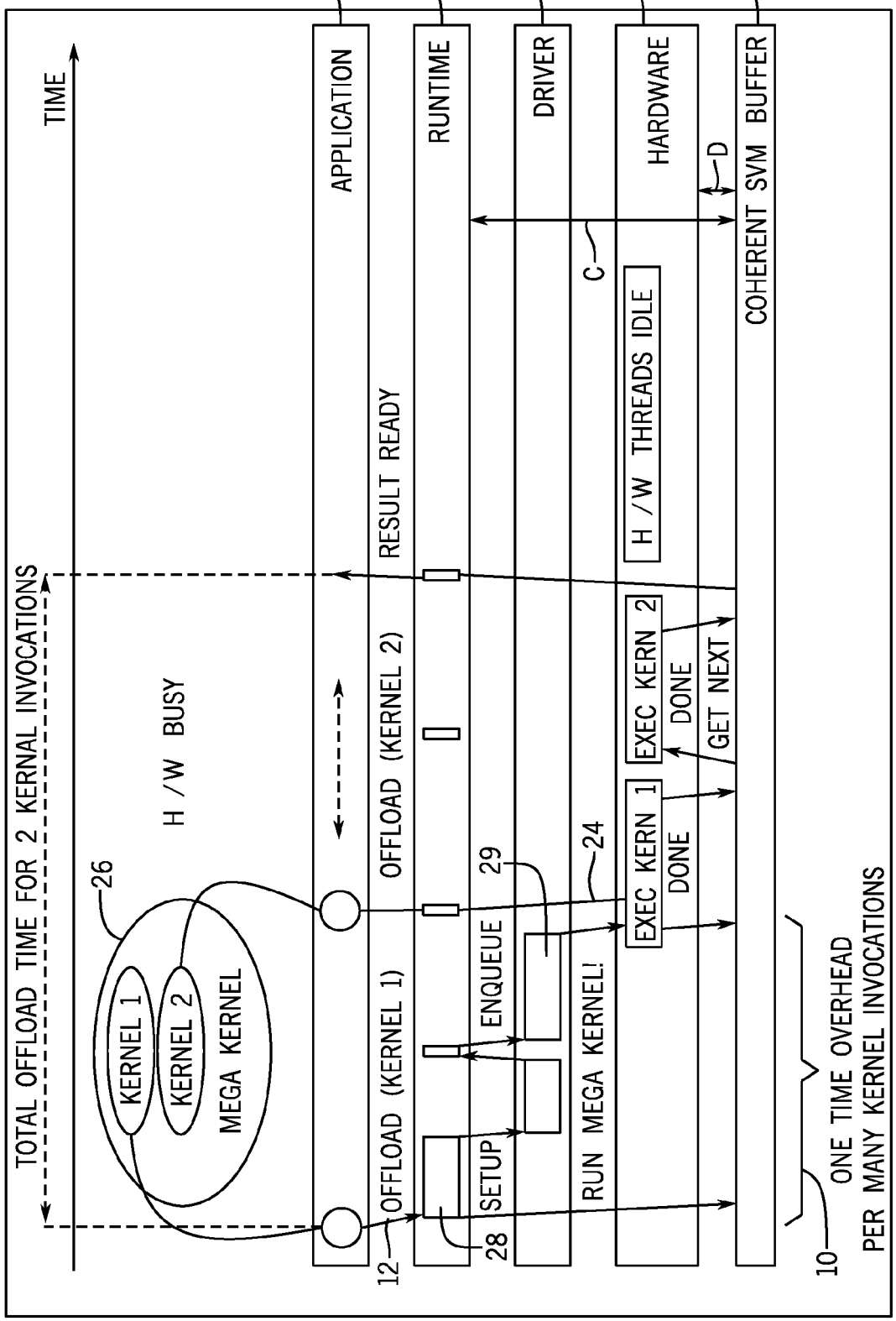

TOTAL OFFLOAD TIME FOR 2 KERNAL INVOCATIONS
TIME
H /W BUSY
KERNEL 1
KERNEL 2
MEGA KERNEL
26
APPLICATION
RUNTIME
DRIVER
HARDWARE
COHERENT SVM BUFFER
14
16
18
20
22
OFFLOAD (KERNEL 1)
OFFLOAD (KERNEL 2)
RESULT READY
SETUP
ENQUEUE
RUN MEGA KERNEL!
H /W THREADS IDLE
EXEC KERN 1
DONE
EXEC KERN 2
DONE
GET NEXT
C
D
12
28
29
24
10
ONE TIME OVERHEAD
PER MANY KERNEL INVOCATIONS

FIG. 1

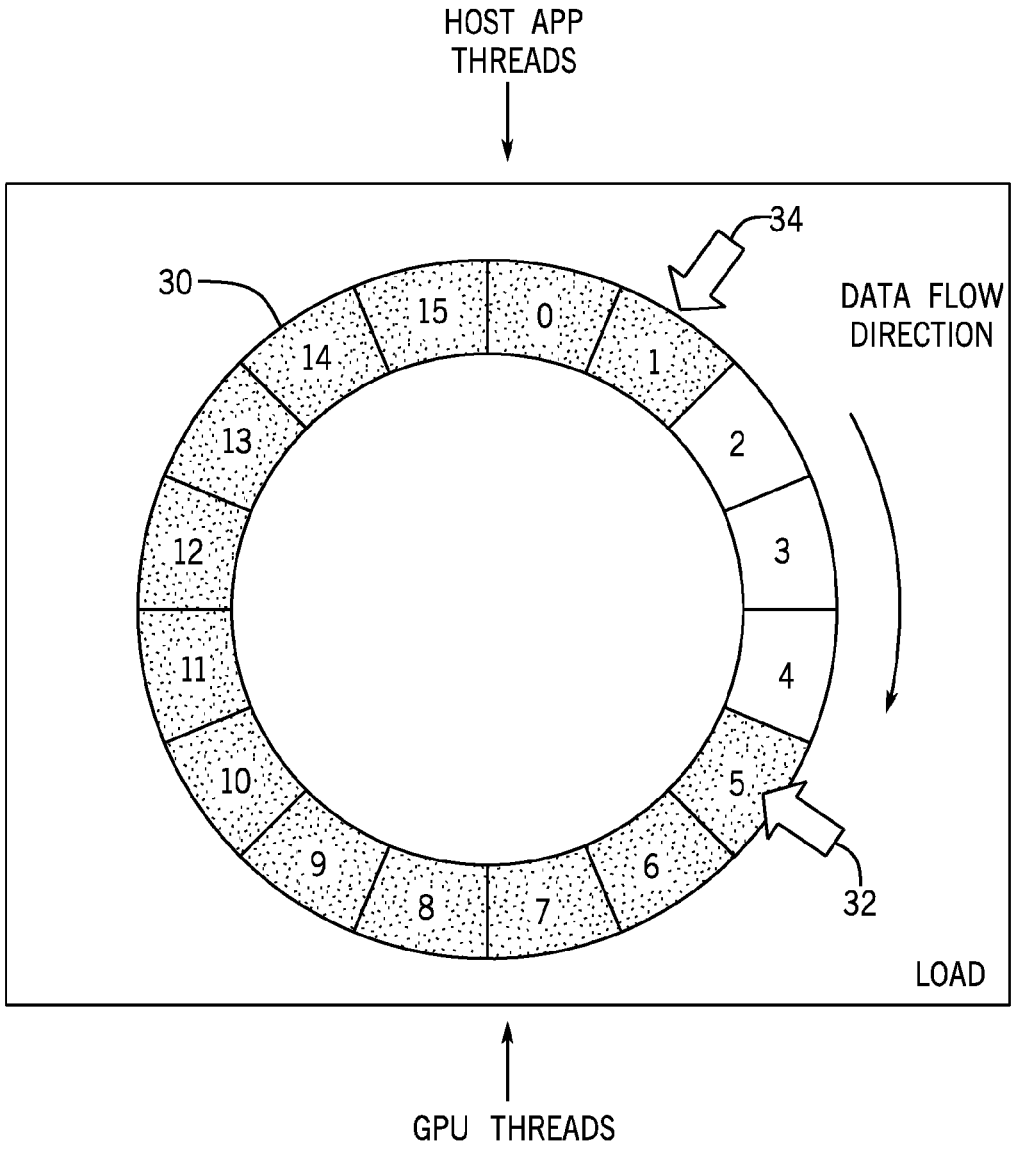

HOST APP
THREADS
34
30
DATA FLOW
DIRECTION
15
0
1
14
2
13
3
12
4
11
5
10
6
9
8
7
32
LOAD
GPU THREADS

FIG. 2

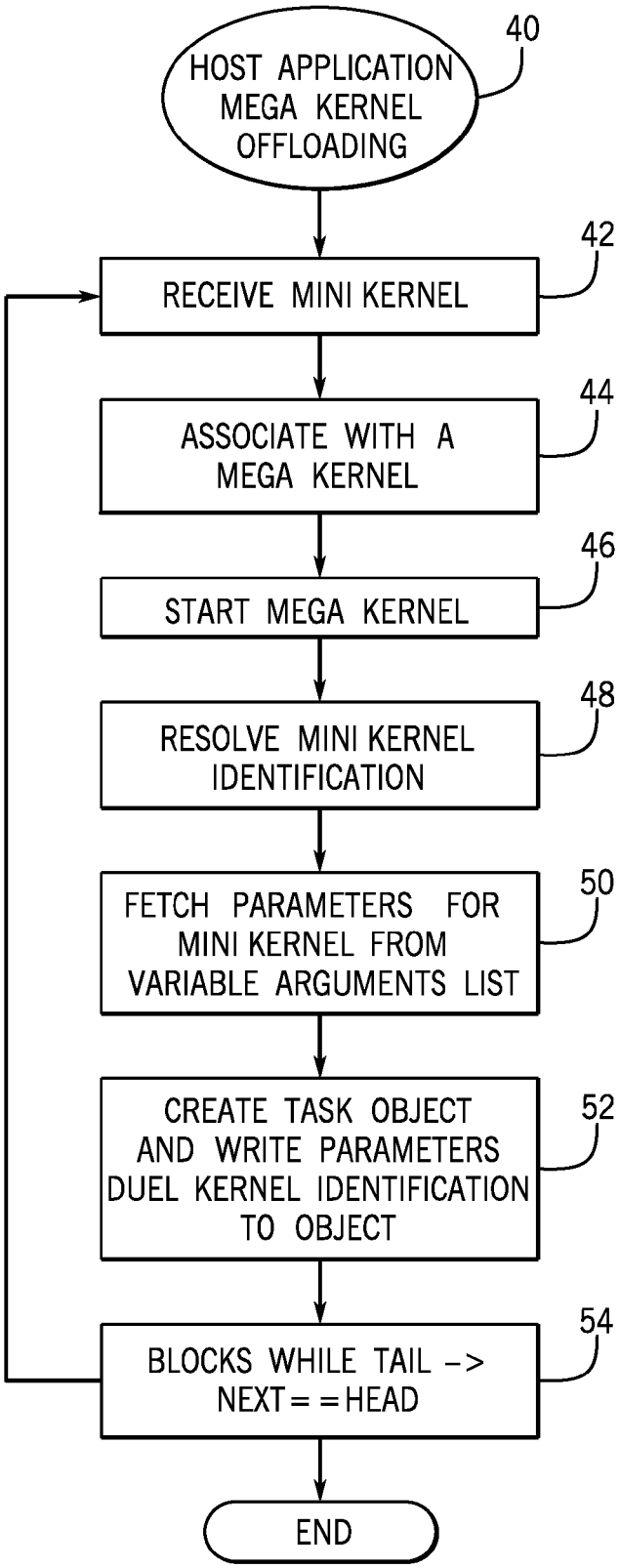

HOST APPLICATION MEGA KERNEL OFFLOADING
40
RECEIVE MINI KERNEL
42
ASSOCIATE WITH A MEGA KERNEL
44
START MEGA KERNEL
46
RESOLVE MINI KERNEL IDENTIFICATION
48
FETCH PARAMETERS FOR MINI KERNEL FROM VARIABLE ARGUMENTS LIST
50
CREATE TASK OBJECT AND WRITE PARAMETERS DUEL KERNEL IDENTIFICATION TO OBJECT
52
BLOCKS WHILE TAIL -> NEXT = = HEAD
54
END

FIG. 3

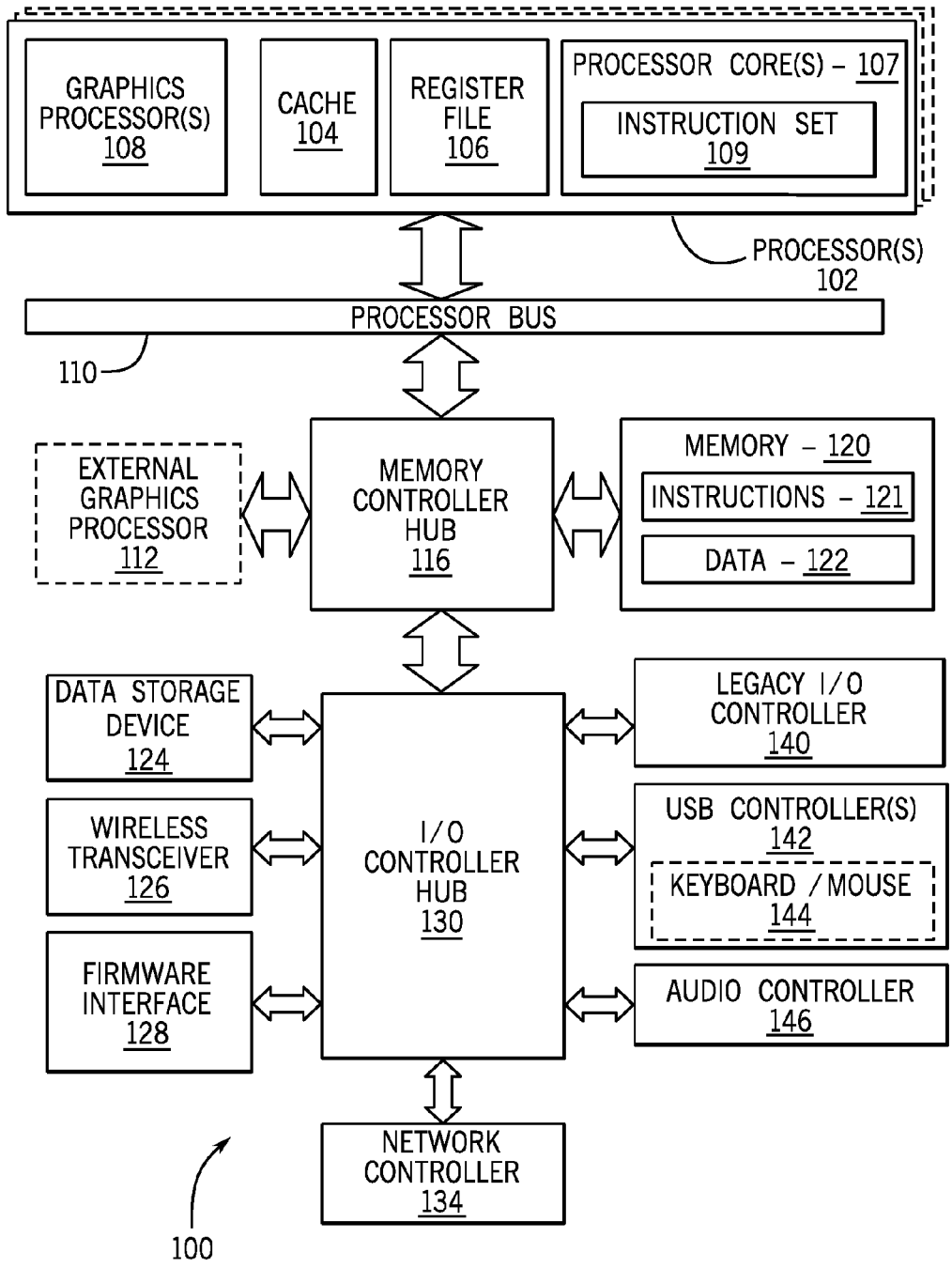

GRAPHICS PROCESSOR(S) 108
CACHE 104
REGISTER FILE 106
PROCESSOR CORE(S) – 107
INSTRUCTION SET 109
PROCESSOR(S) 102
PROCESSOR BUS
110
EXTERNAL GRAPHICS PROCESSOR 112
MEMORY CONTROLLER HUB 116
MEMORY – 120
INSTRUCTIONS – 121
DATA – 122
DATA STORAGE DEVICE 124
WIRELESS TRANSCEIVER 126
FIRMWARE INTERFACE 128
I/O CONTROLLER HUB 130
LEGACY I/O CONTROLLER 140
USB CONTROLLER(S) 142
KEYBOARD / MOUSE 144
AUDIO CONTROLLER 146
NETWORK CONTROLLER 134
100

FIG. 4

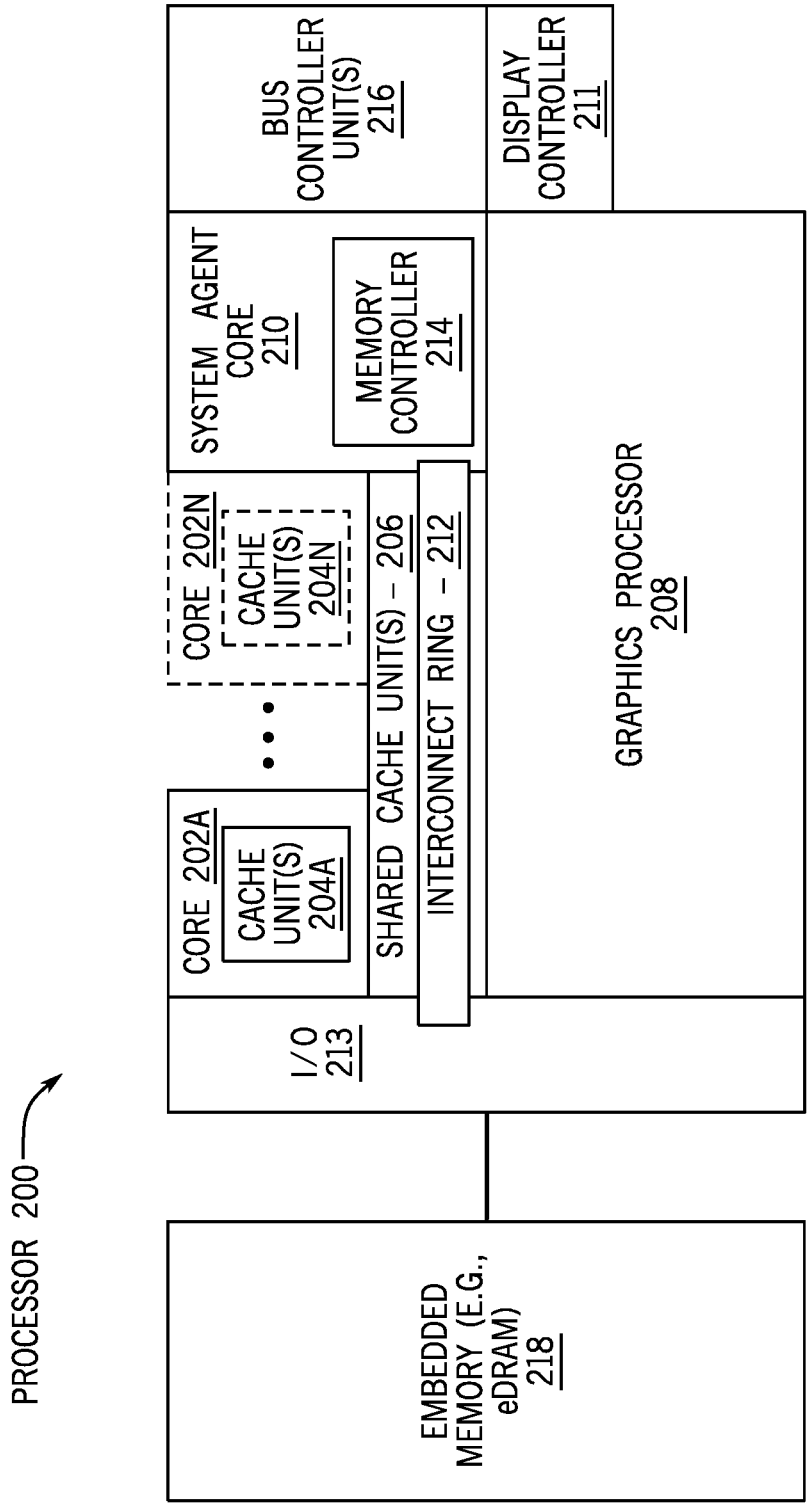

PROCESSOR 200
EMBEDDED MEMORY (E.G., eDRAM) 218
I/O 213
CORE 202A
CACHE UNIT(S) 204A
CORE 202N
CACHE UNIT(S) 204N
SHARED CACHE UNIT(S) – 206
INTERCONNECT RING – 212
SYSTEM AGENT CORE 210
MEMORY CONTROLLER 214
GRAPHICS PROCESSOR 208
BUS CONTROLLER UNIT(S) 216
DISPLAY CONTROLLER 211

FIG. 5

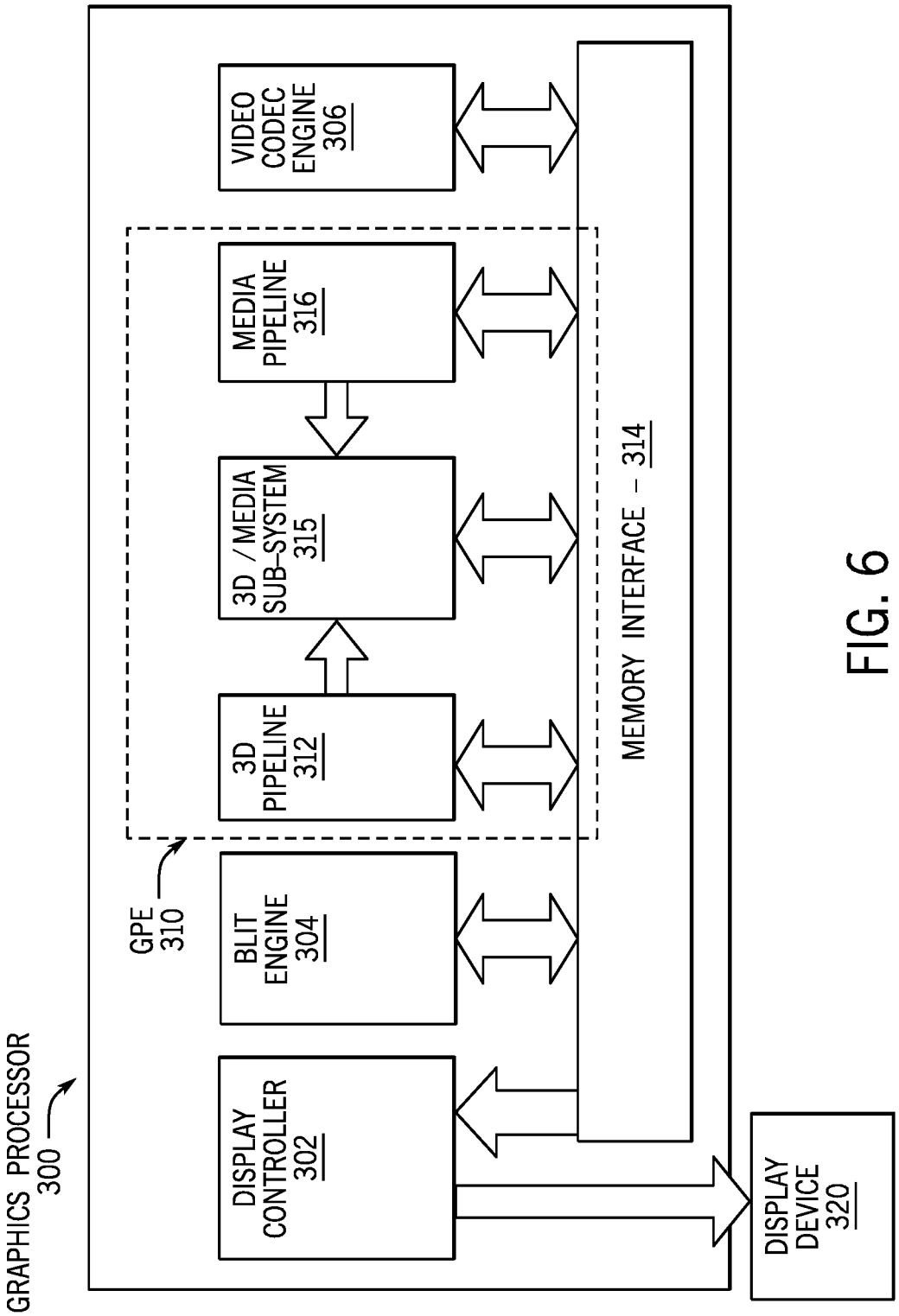

GRAPHICS PROCESSOR 300
DISPLAY CONTROLLER 302
BLIT ENGINE 304
GPE 310
3D PIPELINE 312
3D /MEDIA SUB–SYSTEM 315
MEDIA PIPELINE 316
VIDEO CODEC ENGINE 306
MEMORY INTERFACE – 314
DISPLAY DEVICE 320

FIG. 6

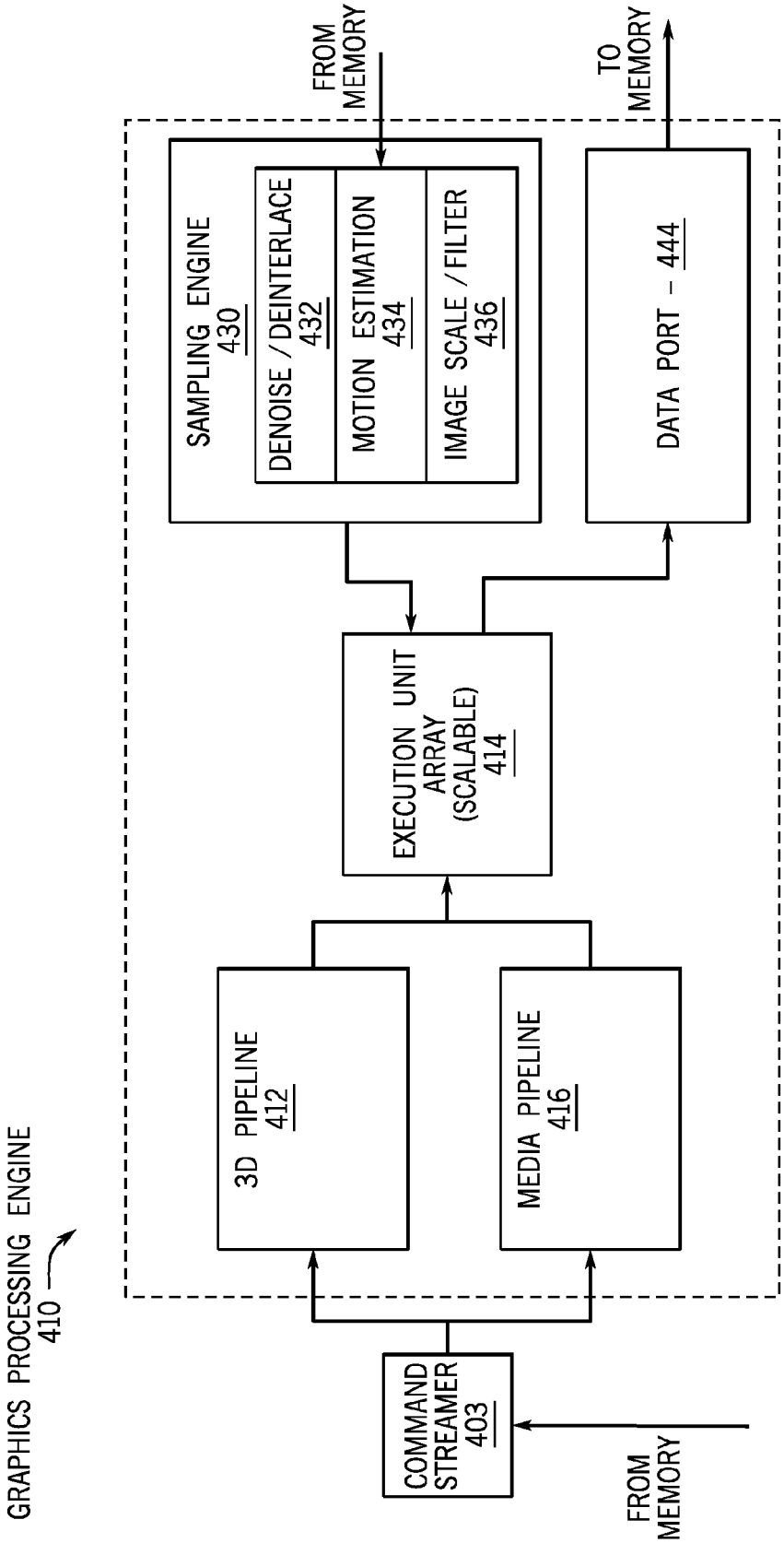

GRAPHICS PROCESSING ENGINE
410
FROM MEMORY
TO MEMORY
SAMPLING ENGINE
430
DENOISE / DEINTERLACE
432
MOTION ESTIMATION
434
IMAGE SCALE / FILTER
436
DATA PORT – 444
EXECUTION UNIT ARRAY (SCALABLE)
414
3D PIPELINE
412
MEDIA PIPELINE
416
COMMAND STREAMER
403
FROM MEMORY

FIG. 7

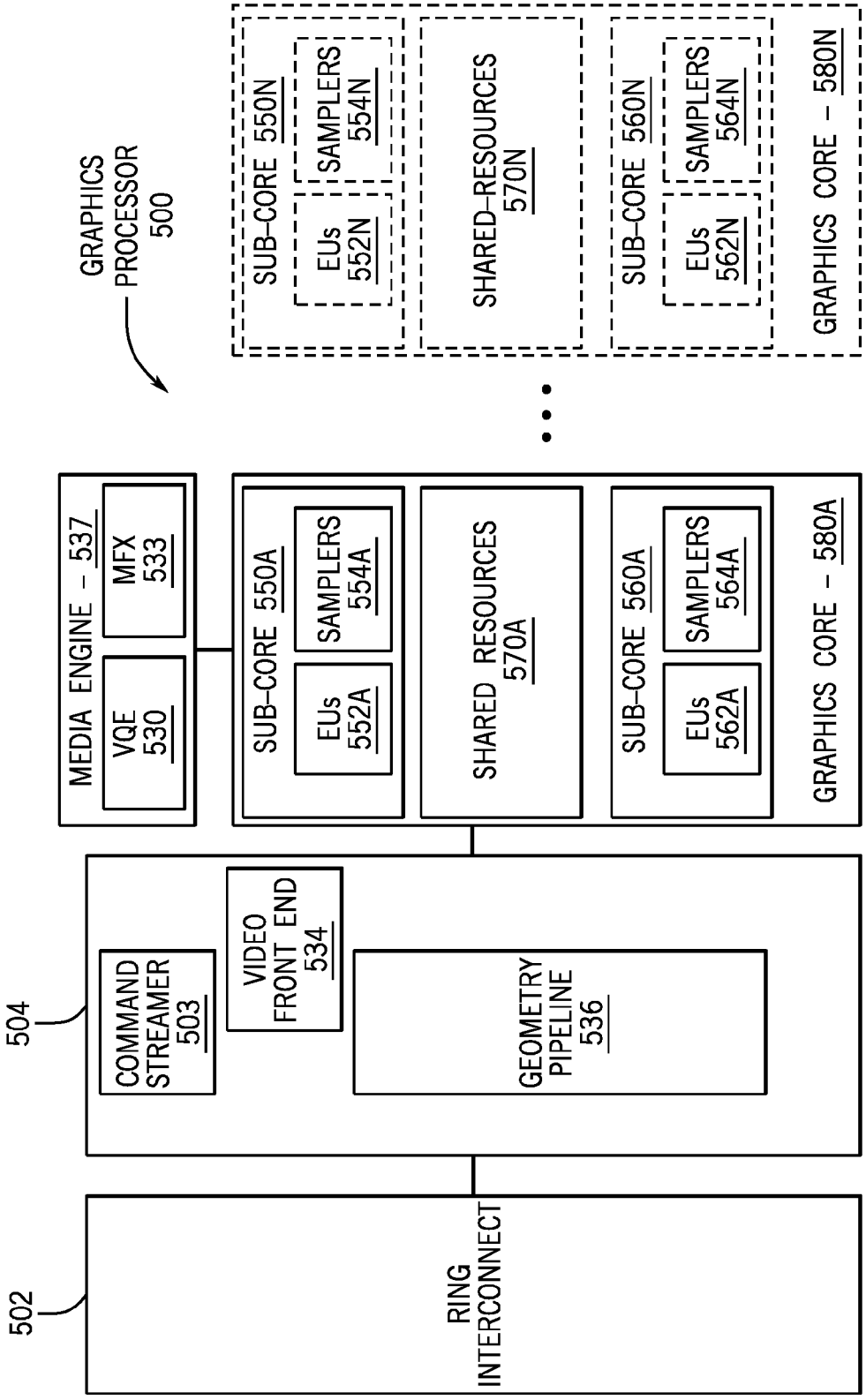

GRAPHICS PROCESSOR 500
SUB-CORE 550N
EUs 552N
SAMPLERS 554N
SHARED-RESOURCES 570N
SUB-CORE 560N
EUs 562N
SAMPLERS 564N
GRAPHICS CORE - 580N
MEDIA ENGINE - 537
VQE 530
MFX 533
SUB-CORE 550A
EUs 552A
SAMPLERS 554A
SHARED RESOURCES 570A
SUB-CORE 560A
EUs 562A
SAMPLERS 564A
GRAPHICS CORE - 580A
504
COMMAND STREAMER 503
VIDEO FRONT END 534
GEOMETRY PIPELINE 536
502
RING INTERCONNECT

FIG. 8

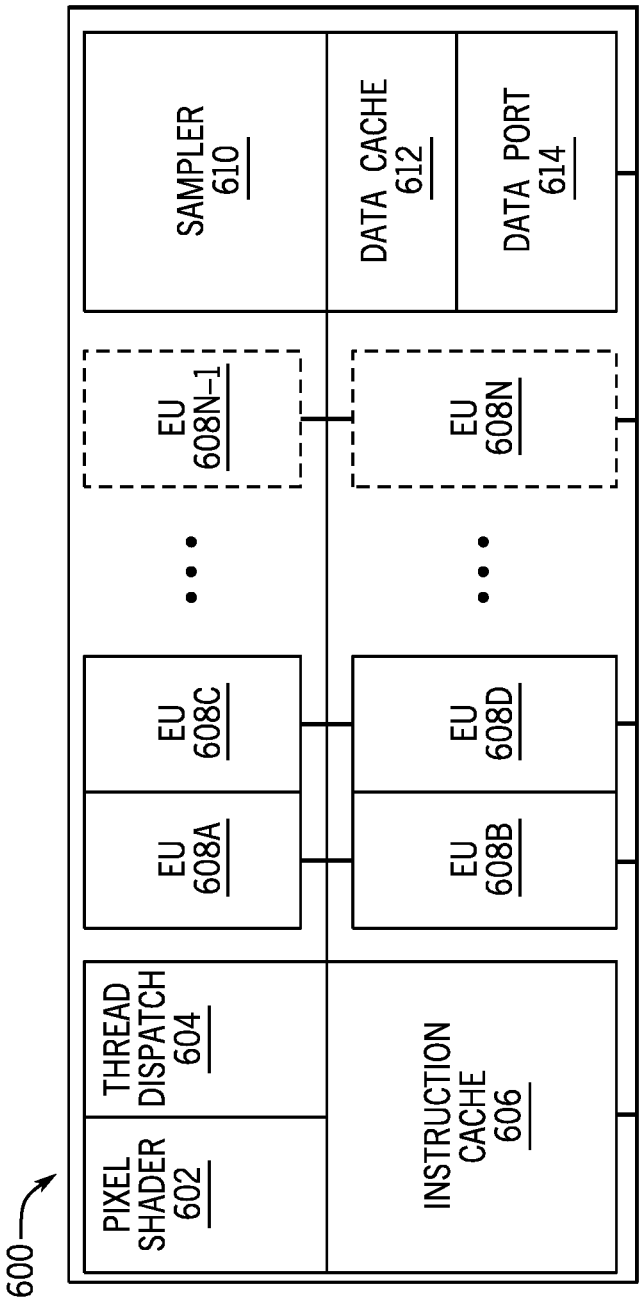

600
PIXEL SHADER 602
THREAD DISPATCH 604
INSTRUCTION CACHE 606
EU 608A
EU 608C
EU 608N–1
EU 608B
EU 608D
EU 608N
SAMPLER 610
DATA CACHE 612
DATA PORT 614

FIG. 9

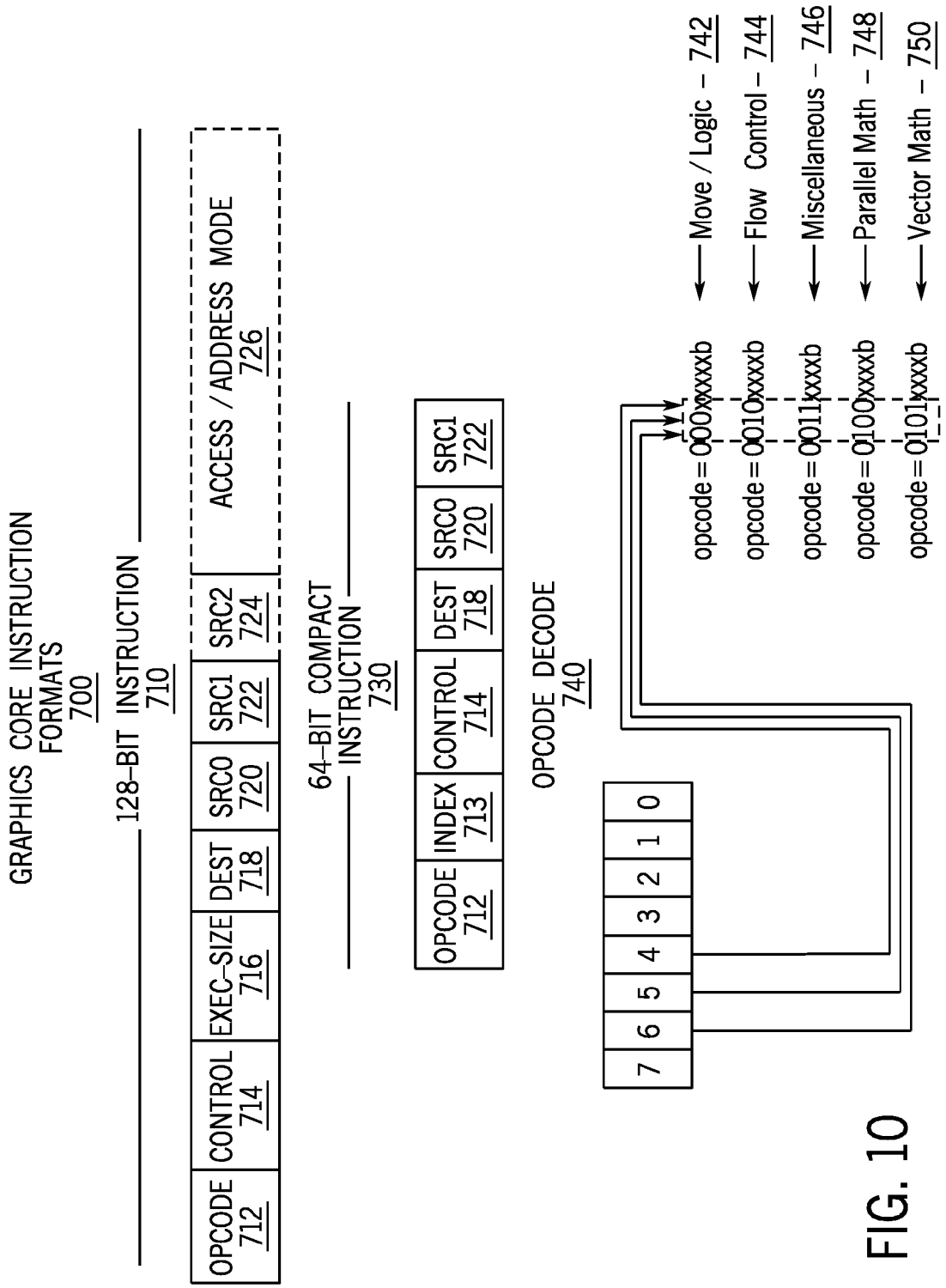

GRAPHICS CORE INSTRUCTION FORMATS 700
128-BIT INSTRUCTION 710
OPCODE 712
CONTROL 714
EXEC-SIZE 716
DEST 718
SRC0 720
SRC1 722
SRC2 724
ACCESS / ADDRESS MODE 726
64-BIT COMPACT INSTRUCTION 730
OPCODE 712
INDEX 713
CONTROL 714
DEST 718
SRC0 720
SRC1 722
OPCODE DECODE 740
7 6 5 4 3 2 1 0
opcode = 000xxxxb
opcode = 001Oxxxxb
opcode = 0011xxxb
opcode = 0100xxxb
opcode = 0101xxxb
Move / Logic – 742
Flow Control – 744
Miscellaneous – 746
Parallel Math – 748
Vector Math – 750

FIG. 10

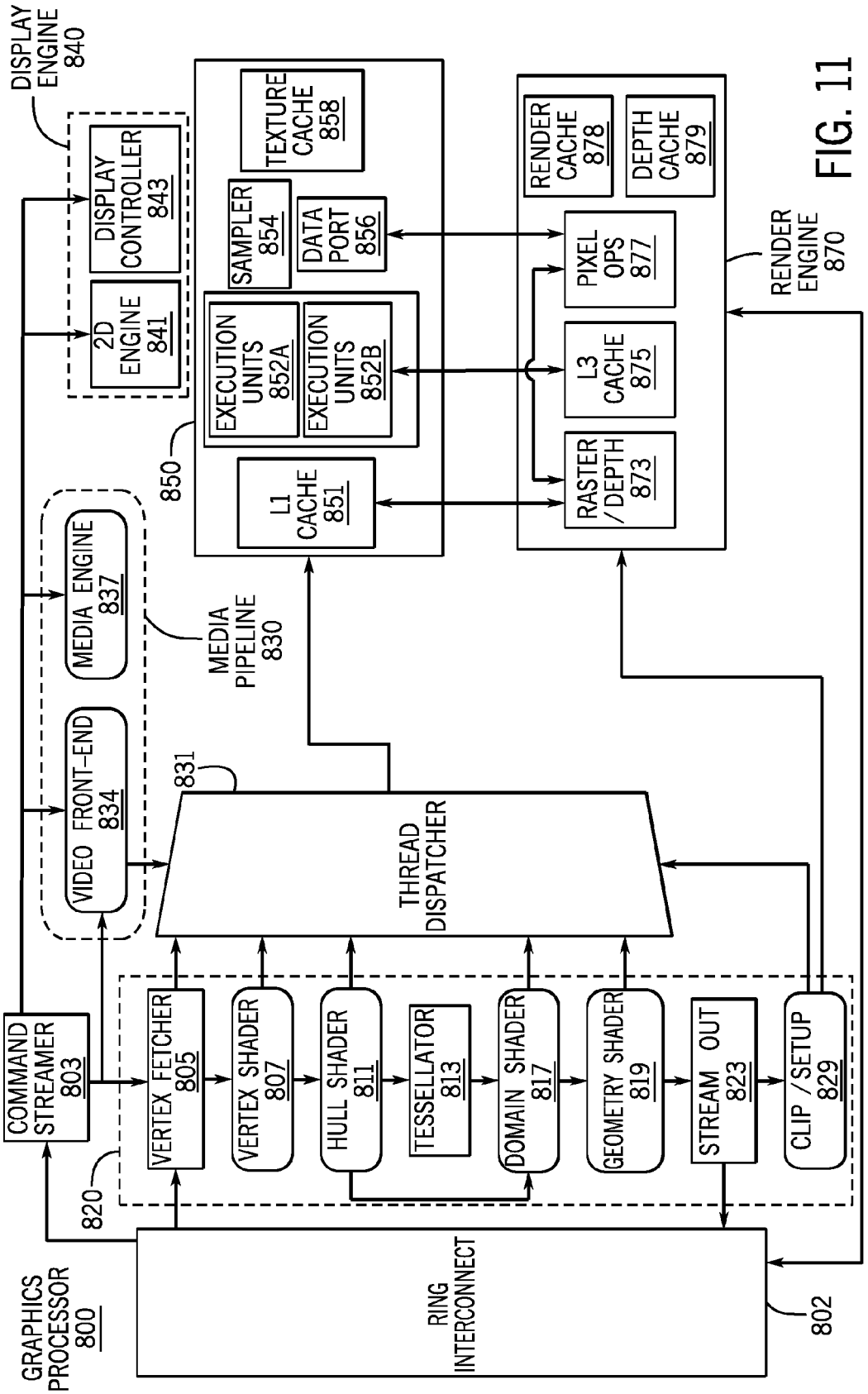

GRAPHICS PROCESSOR 800
RING INTERCONNECT
802
820
COMMAND STREAMER 803
VERTEX FETCHER 805
VERTEX SHADER 807
HULL SHADER 811
TESSELLATOR 813
DOMAIN SHADER 817
GEOMETRY SHADER 819
STREAM OUT 823
CLIP /SETUP 829
THREAD DISPATCHER
831
MEDIA PIPELINE 830
VIDEO FRONT-END 834
MEDIA ENGINE 837
DISPLAY ENGINE 840
2D ENGINE 841
DISPLAY CONTROLLER 843
850
L1 CACHE 851
EXECUTION UNITS 852A
EXECUTION UNITS 852B
SAMPLER 854
DATA PORT 856
TEXTURE CACHE 858
RENDER ENGINE 870
RASTER /DEPTH 873
L3 CACHE 875
PIXEL OPS 877
RENDER CACHE 878
DEPTH CACHE 879

FIG. 11

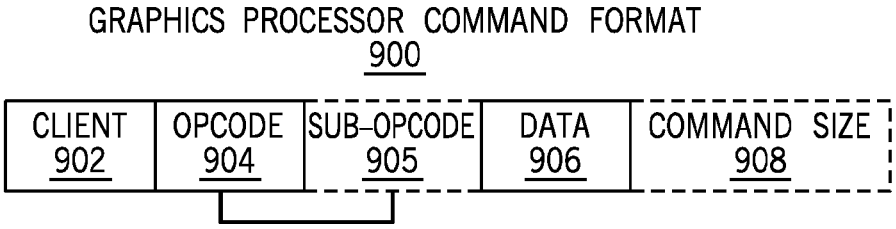


FIG. 12A

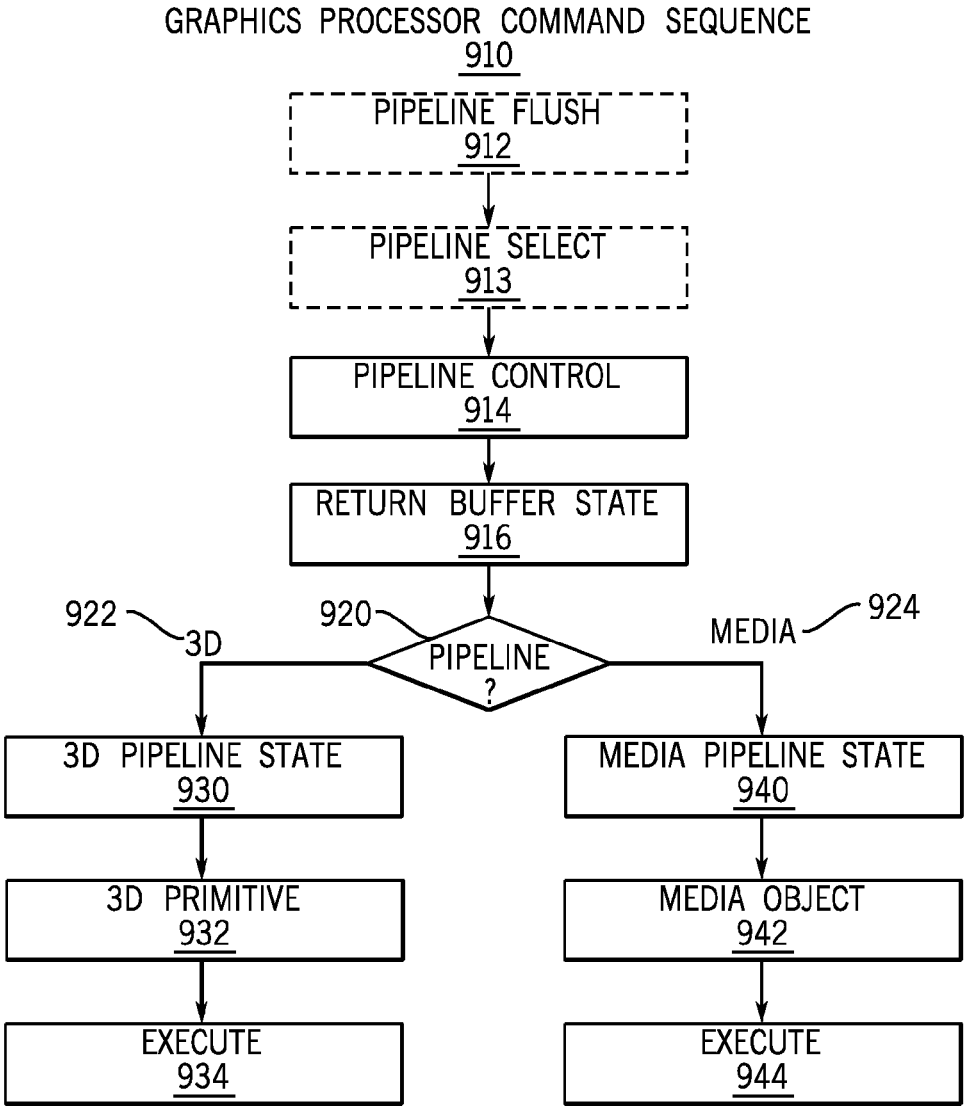


FIG. 12B

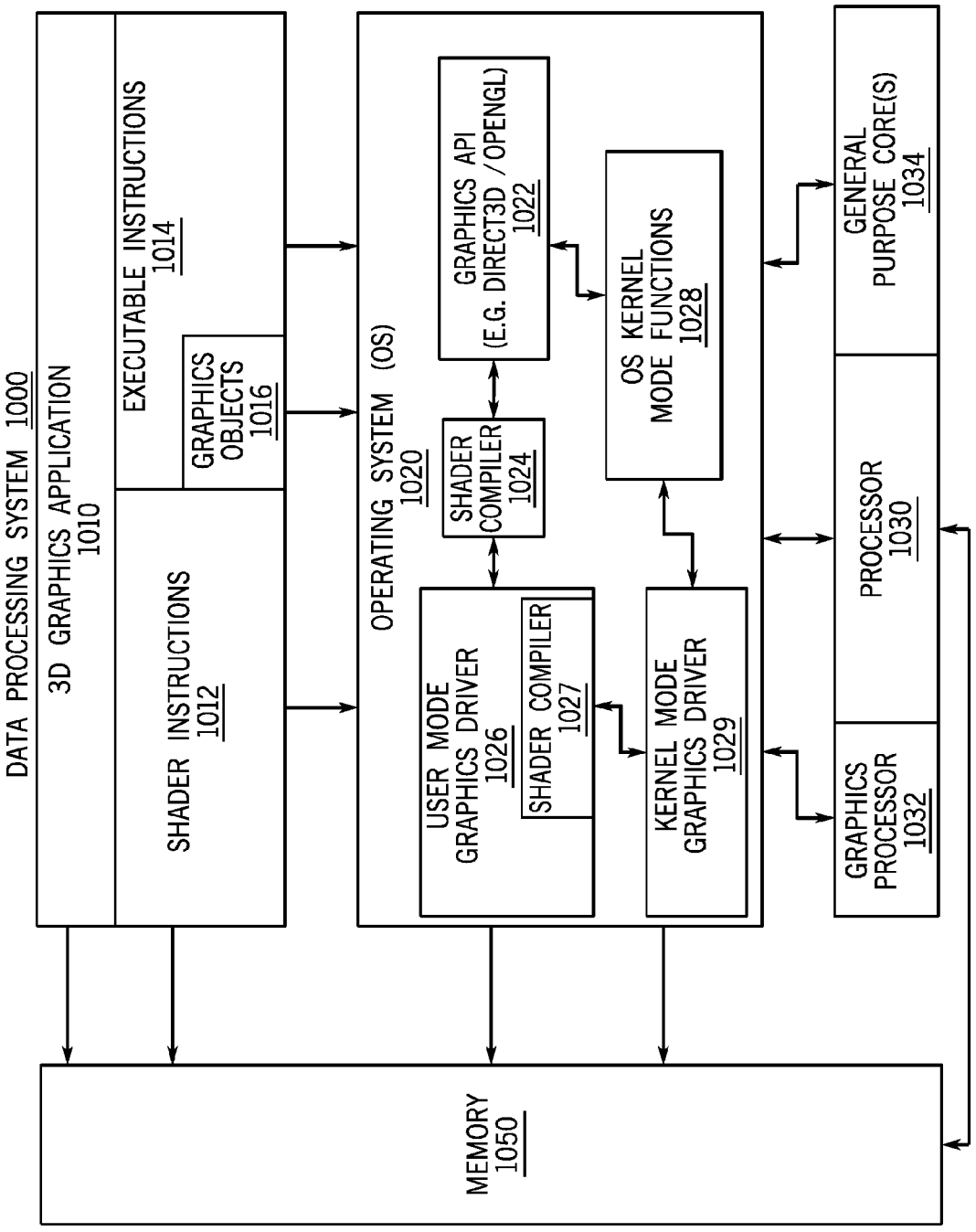

DATA PROCESSING SYSTEM 1000
3D GRAPHICS APPLICATION 1010
EXECUTABLE INSTRUCTIONS 1014
GRAPHICS OBJECTS 1016
SHADER INSTRUCTIONS 1012
OPERATING SYSTEM (OS) 1020
GRAPHICS API (E.G. DIRECT3D /OPENGL) 1022
SHADER COMPILER 1024
OS KERNEL MODE FUNCTIONS 1028
USER MODE GRAPHICS DRIVER 1026
SHADER COMPILER 1027
KERNEL MODE GRAPHICS DRIVER 1029
GENERAL PURPOSE CORE(S) 1034
PROCESSOR 1030
GRAPHICS PROCESSOR 1032
MEMORY 1050

FIG. 13

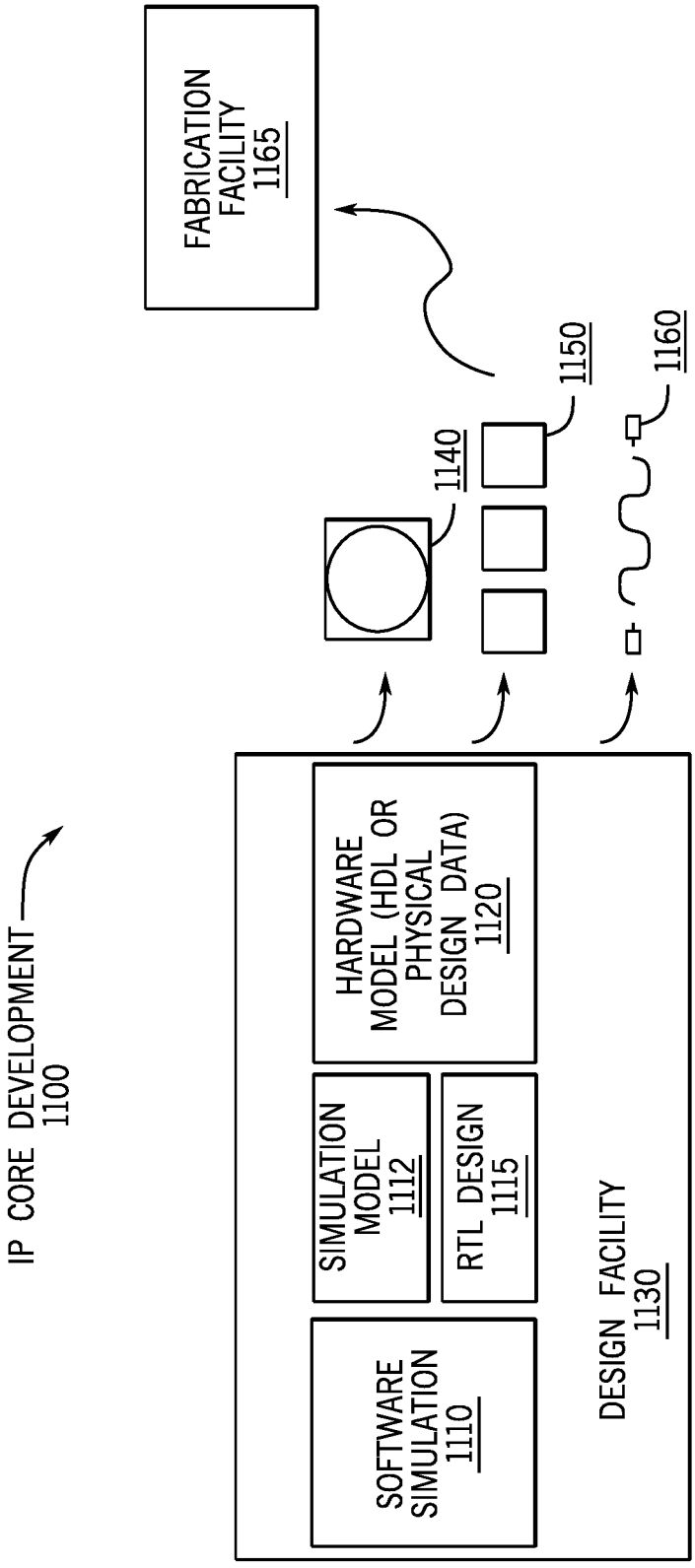

IP CORE DEVELOPMENT
1100
SOFTWARE SIMULATION
1110
SIMULATION MODEL
1112
RTL DESIGN
1115
HARDWARE MODEL (HDL OR PHYSICAL DESIGN DATA)
1120
DESIGN FACILITY
1130
1140
1150
1160
FABRICATION FACILITY
1165

FIG. 14

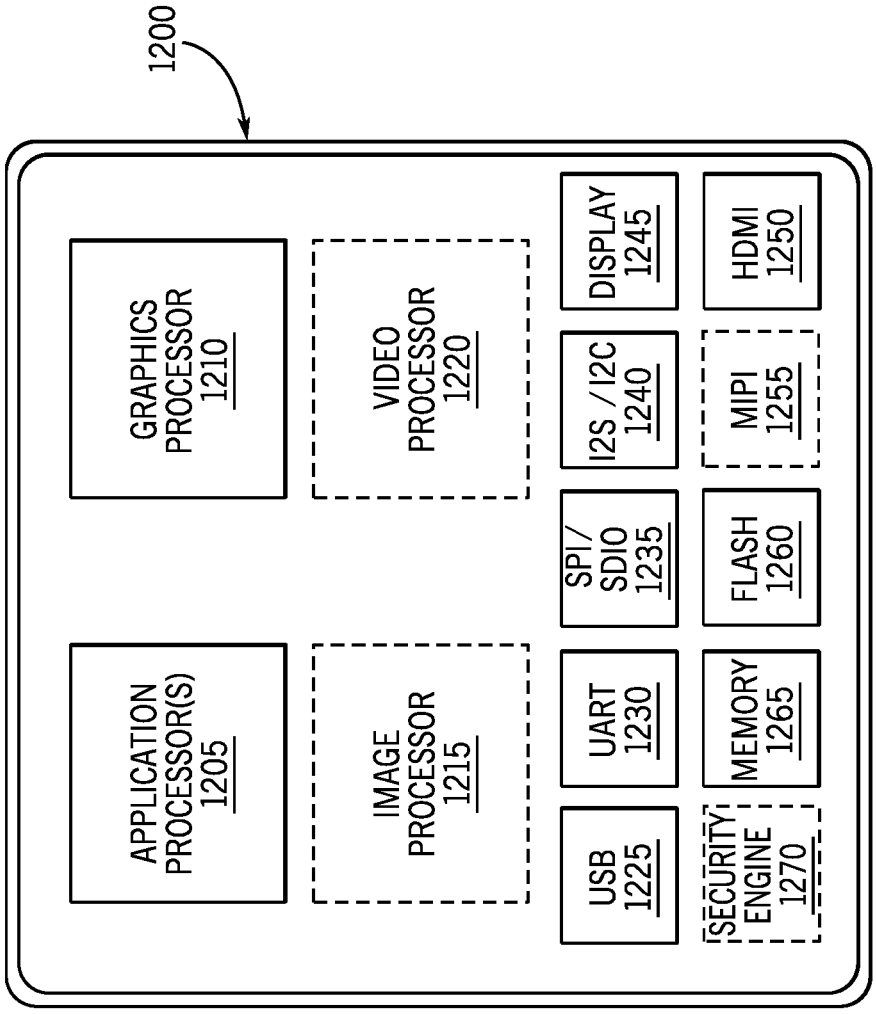

1200
GRAPHICS PROCESSOR 1210
VIDEO PROCESSOR 1220
DISPLAY 1245
HDMI 1250
I2S / I2C 1240
MIPI 1255
SPI / SDIO 1235
FLASH 1260
APPLICATION PROCESSOR(S) 1205
IMAGE PROCESSOR 1215
UART 1230
MEMORY 1265
USB 1225
SECURITY ENGINE 1270

FIG. 15

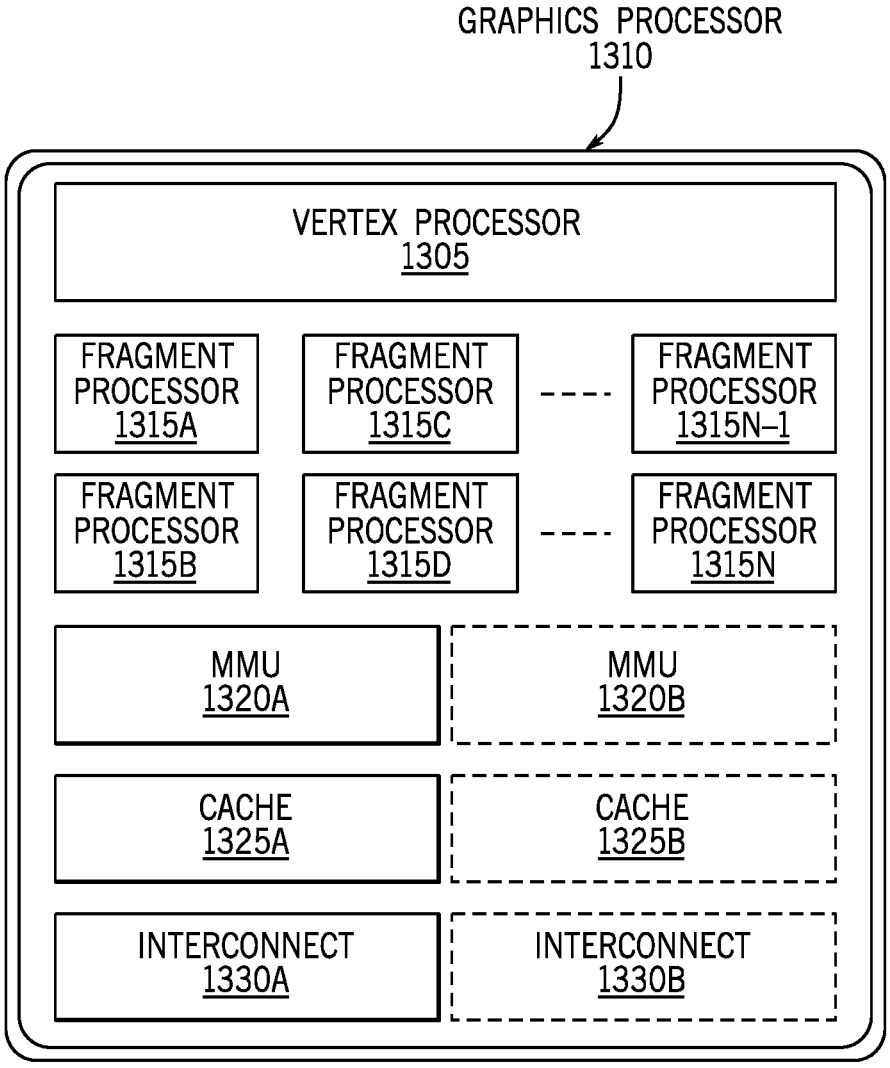

GRAPHICS PROCESSOR
1310
VERTEX PROCESSOR
1305
FRAGMENT PROCESSOR
1315A
FRAGMENT PROCESSOR
1315C
FRAGMENT PROCESSOR
1315N–1
FRAGMENT PROCESSOR
1315B
FRAGMENT PROCESSOR
1315D
FRAGMENT PROCESSOR
1315N
MMU
1320A
MMU
1320B
CACHE
1325A
CACHE
1325B
INTERCONNECT
1330A
INTERCONNECT
1330B

FIG. 16

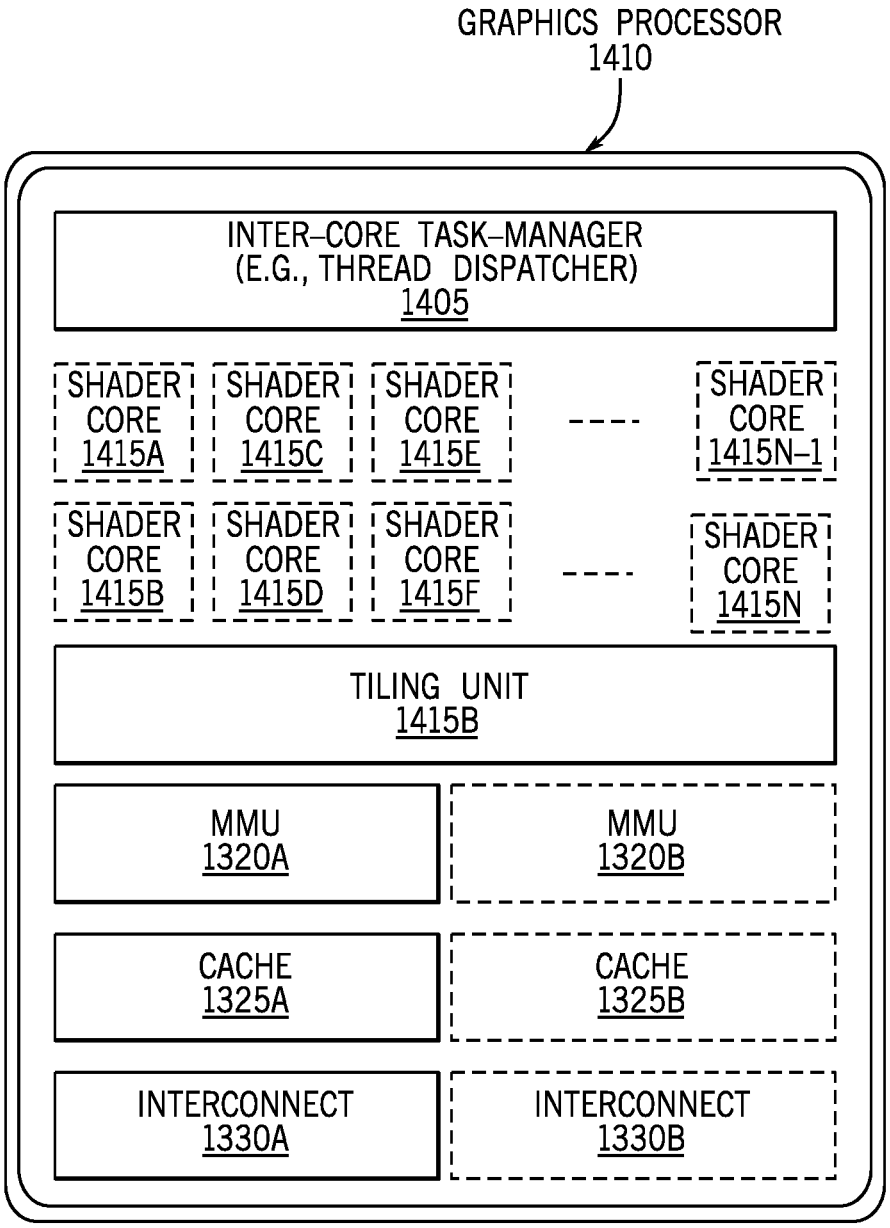

GRAPHICS PROCESSOR
1410
INTER-CORE TASK-MANAGER
(E.G., THREAD DISPATCHER)
1405
SHADER CORE 1415A
SHADER CORE 1415C
SHADER CORE 1415E
SHADER CORE 1415N-1
SHADER CORE 1415B
SHADER CORE 1415D
SHADER CORE 1415F
SHADER CORE 1415N
TILING UNIT
1415B
MMU
1320A
MMU
1320B
CACHE
1325A
CACHE
1325B
INTERCONNECT
1330A
INTERCONNECT
1330B

FIG. 17

OFFLOADING FUSED KERNEL EXECUTION TO A GRAPHICS PROCESSOR

BACKGROUND

[0001] A compute kernel is an application function whose execution is offloaded to a general-purpose graphics computing unit (GPGPU) device available in the system. Integrated graphics is an example of a GPGPU device and may be a graphics processing unit integrated with a central processing unit. Standard schemes of offloading kernel execution to integrated graphics assume interaction between the heterogeneous application and the underlying software stack. The software stack has various runtimes and User-Mode and Kernel-Mode drivers, used for offloading, that are at the bottom of the stack.

[0002] Kernel offload traditionally incurs significant overhead from the various software layers and protection ring transitions of the stack. Therefore, offloading of kernels of comparable or less execution time on a central processing unit (CPU) cannot amortize the offload overhead.

BRIEF DESCRIPTION OF THE DRAWINGS

[0003] Some embodiments are described with respect to the following figures:

[0004] FIG. **1** is a sequence diagram for one embodiment;

[0005] FIG. **2** is a ring buffer schematic for one embodiment;

[0006] FIG. **3** is a flow chart for one embodiment;

[0007] FIG. **4** is a block diagram of a processing system according to one embodiment;

[0008] FIG. **5** is a block diagram of a processor according to one embodiment;

[0009] FIG. **6** is a block diagram of a graphics processor according to one embodiment;

[0010] FIG. **7** is a block diagram of a graphics processing engine according to one embodiment;

[0011] FIG. **8** is a block diagram of another embodiment of a graphics processor;

[0012] FIG. **9** is a depiction of thread execution logic according to one embodiment;

[0013] FIG. **10** is a block diagram of a graphics processor instruction format according to some embodiments;

[0014] FIG. **11** is a block diagram of another embodiment of a graphics processor;

[0015] FIG. **12**A is a block diagram of a graphics processor command format according to some embodiments;

[0016] FIG. **12**B is a block diagram illustrating a graphics processor command sequence according to some embodiments;

[0017] FIG. **13** is a depiction of an exemplary graphics software architecture according to some embodiments;

[0018] FIG. **14** is a block diagram illustrating an IP core development system according to some embodiments;

[0019] FIG. **15** is a block diagram showing an exemplary system on chip integrated circuit according to some embodiments;

[0020] FIG. **16** is a block diagram of a graphics processor in a system on a chip according to one embodiment; and

[0021] FIG. **17** is a block diagram of another graphics processor according to one embodiment.

DETAILED DESCRIPTION

[0022] This kernel offloading overhead may be avoided without losing generality of offload features. Also small kernel offloading can still be interleaved with CPU code using the offload results. Multiple kernels are automatically grouped together by a compiler and linker.

[0023] Hardware threads run continuously and use coherent Shared Virtual Memory (SVM) buffers for direct communication between the host application and the graphics processing unit (GPU) hardware threads. "Coherent" means that different agents (CPU, GPU) accessing the buffer immediately see updates to the buffer (writes) made by other agents. Application GPU kernels ("mini kernels") are merged into a single combined or "mega kernel" which runs continuously by every GPU hardware thread.

[0024] The host launches the mega kernel once via the high-overhead driver mechanism. Then the host submits offload tasks directly to the coherent SVM buffer, bypassing the driver. The mega kernel is basically a wrapper around a switch control statement on different mini kernels. A switch statement is as follows:

```
switch (kernel_id)  {
case 0
     call kernel_0( );
     break;
case 1"
     call kernel_1( );
     break;
...
}
```

[0025] There is a dedicated GPU "master" thread efficiently distinguished via its hardware identifier (e.g. 0) which does all the communication with the host and dispatches tasks to other "worker" threads.

[0026] As shown in FIG. **1**, driver interaction overhead **10** is present only when enqueuing the mega kernel on the first mini kernel offload **12** from an application **14** to runtime **16**. Subsequent offloads of mini kernels **24**, which are part of the running mega kernel **26**, bypass the driver **18** and simply write to the coherent SVM communication buffer **22**.

[0027] The mega kernel is setup at **28** during runtime. The setup includes setting up threads, arguments, and internal control structures. It also includes launching the mega kernel on the GPGPU device.

[0028] The hardware blocks marked "DONE" indicate each mini kernel finished execution. "Done" means signaling kernel execution completion via writing to the coherent SVM buffer by the CPU leader thread. "Get next" means fetching next kernel for execution by the leader thread from the ring buffer allocated in the coherent SVM buffer. In an optimized case there are no control transfers out of the hardware between the two "DONE" blocks, and the time taken to switch between mini kernels is smaller. In FIG. **1**, the x-axis is time. The hardware is "busy" whenever it executes a mini kernel (to avoid clutter "busy" is shown only for the first kernel), as indicated in FIG. **1**.

[0029] The unlabeled arrows in FIG. **1** indicate passing of control, except unlabeled arrows to/from the coherent SVM buffer. Those arrows, as well as the arrows C and D, from hardware to the coherent SVM buffer and from the coherent SVM buffer to runtime, indicate passing of data.

[0030] A key element of the communication is a ring task buffer **22**, shown in FIG. **2**, with a fixed number of tasks slots **30** (to eliminate memory management). The system maintains two pointers—head **32** and tail **34**—to tasks in the ring. When head==tail there is no task to execute. The host inserts a new task at the tail, moving the tail to the next ring element. If the next element is head (tail→next==head), this means that no free task slots are available, and the host blocks until GPU completes the current task and moves the head to the next element. GPU leader thread waits until head!=tail and dispatches the task pointed to head. After task completion, it moves the head to the next element. Shaded ring buffer slots in FIG. **2** depict tasks for kernel(s) execution which are not yet complete. Moving the head to the next element "clears" the shading.

[0031] Pseudo-code for a mega kernel is as follows:

```
while (true) {
    If (<current thread is the leader>) {
        spin while (head==tail);
        if (head is special "EXIT' task) {
            write CMD_EXIT command for soldier threads
            break; //exists the outermost loop
        }
        copy parameters from head task to <parameter area>;
        copy kernel_id from head tasks to <kernel_id area>
        reset number of busy threads to N
        store-release memory barrier
        write CMD_RUN command for soldier threads
        command = CMD_RUN
    }
        else        {
        command = read command from the leader
    }
        if (command==CMD_RUN)   {
            kernel_id = read <kernel_id area>
            switch (kernel_id) {
            case 1: call kernel1; break;
            case 2: call kernel2; break;
        ...
        }
        if (current thread is the leader>)        {
            spin while busy thread count is greater than 1
            write CMD_WAIT command for soldier threads
            store-release memory barrier
            head=head->next
            // untame waiting soldier threads;
            write 0 to busy thread count
        }
        else        {
            automatically decrement busy thread count
            spin while busy thread count is greater than 0
        }
    }
    else if (command == CMD_EXIT) {
    }
    yield ( );
}
```

[0032] The mini kernels are now called from within the mega kernel, unlike the usual case where they were called from the host. This requires (a) construction of the mega kernel and (b) special compilation of mini kernels so that they read their parameters from a memory buffer instead of using a traditional mechanism for passing parameters from host to a GPU kernel. This may be done in runtime setup **28** in FIG. **1**.

[0033] The set of mini kernels participating in the mega kernel is decided at link time (after compile time but before runtime) based on which kernels are marked as mini. A linker replaces call targets within the wrapper function's switch statement with real mini kernels and embeds a map from a mega kernel to its constituent mini kernels. This information is used by the runtime when offloading a kernel by name. There can be more than one mega kernel.

[0034] On the host side, the offload interface does not change. It remains the same as for the "direct kernel" offload model, where application program interface (API) call _GFX_enqueue (mini_kernel_host_pointer, kernel_parameters) is used to offload a mini kernel as indicated at enqueue **29** in FIG. **1**. Upon this call, the host application:

[0035] finds a mega kernel this mini kernel belongs to;

[0036] starts the mega kernel if it is not yet running;

[0037] resolves kernel identification for the mini kernel (for example, simply the ordinal number of the mini kernel within the mega kernel);

[0038] fetches parameters from the variable argument list;

[0039] creates a lightweight task object and writes the parameters and the kernel identification to it; and

[0040] blocks while tail→next==head.

[0041] A "cooperative preemption" of the mega kernel task ensures display responsiveness and lower power consumption. The host starts an "interrupter" thread whose purpose is to periodically enqueue a special 'exit' task to make the mega kernel finishes execution and exits. With the next mini kernel offload, the mega-kernel will restart. But in between, the operating system driver has a chance to schedule another GPU task, such as display update.

[0042] To minimize power consumption and keep the GPU idle (not executing the mega-kernel wait-for-task spin loop) when there are no GPU tasks, a mega kernel start/stop application program interface (API) is exposed to the users so that they can decide when to engage the continuous offload scheme.

[0043] Offloading in this fashion makes it profitable (in terms of performance) to offload much smaller kernels (1 ms. or less) than current schemes which suffer in performance due to overhead with existing offload middleware.

[0044] Without this offloading technique, several existing algorithms may require rewrite to create longer running kernels that absorb the offload overhead, making beneficial offload programming or algorithm changes to existing programs, such as JPEG compression and discrete cosine transform (DCT).

[0045] The sequence **40**, shown in FIG. **3**, for host application mega-kernel offloading may be implemented in software, firmware and/or hardware. In software and firmware embodiments it may be implemented by computer executed instructions stored in one or more non-transitory computer readable media such as magnetic, optical, or semiconductor storage.

[0046] This offloading sequence **40** begins by receiving a mini kernel as indicated in block **42**. The mini kernel is associated with a mega-kernel as indicated in block **44**. The mega-kernel is started in block **46**. The mini kernel's identification is resolved in block **48**. Then in block **50** the parameters for the mini kernel are fetched from a variable arguments list.

[0047] Then a task object is created. The parameters and the kernel identification are written to that object as indicated in block **52**. Then the host application blocks while the tile→next==head as indicated in block **54**.

[0048] FIG. **4** is a block diagram of a processing system **100**, according to an embodiment. In various embodiments

the system **100** includes one or more processors **102** and one or more graphics processors **108**, and may be a single processor desktop system, a multiprocessor workstation system, or a server system having a large number of processors **102** or processor cores **107**. In one embodiment, the system **100** is a processing platform incorporated within a system-on-a-chip (SoC) integrated circuit for use in mobile, handheld, or embedded devices.

[0049] The processing system including a graphics processing unit may be an integrated circuit. An integrated circuit means a single integrated silicon die. The die contains the graphics processing unit and parallel interconnected geometry processing fixed-function units.

[0050] An embodiment of system **100** can include, or be incorporated within a server-based gaming platform, a game console, including a game and media console, a mobile gaming console, a handheld game console, or an online game console. In some embodiments system **100** is a mobile phone, smart phone, tablet computing device or mobile Internet device. Data processing system **100** can also include, couple with, or be integrated within a wearable device, such as a smart watch wearable device, smart eyewear device, augmented reality device, or virtual reality device. In some embodiments, data processing system **100** is a television or set top box device having one or more processors **102** and a graphical interface generated by one or more graphics processors **108**.

[0051] In some embodiments, the one or more processors **102** each include one or more processor cores **107** to process instructions which, when executed, perform operations for system and user software. In some embodiments, each of the one or more processor cores **107** is configured to process a specific instruction set **109**. In some embodiments, instruction set **109** may facilitate Complex Instruction Set Computing (CISC), Reduced Instruction Set Computing (RISC), or computing via a Very Long Instruction Word (VLIW). Multiple processor cores **107** may each process a different instruction set **109**, which may include instructions to facilitate the emulation of other instruction sets. Processor core **107** may also include other processing devices, such a Digital Signal Processor (DSP).

[0052] In some embodiments, the processor **102** includes cache memory **104**. Depending on the architecture, the processor **102** can have a single internal cache or multiple levels of internal cache. In some embodiments, the cache memory is shared among various components of the processor **102**. In some embodiments, the processor **102** also uses an external cache (e.g., a Level-3 (L3) cache or Last Level Cache (LLC)) (not shown), which may be shared among processor cores **107** using known cache coherency techniques. A register file **106** is additionally included in processor **102** which may include different types of registers for storing different types of data (e.g., integer registers, floating point registers, status registers, and an instruction pointer register). Some registers may be general-purpose registers, while other registers may be specific to the design of the processor **102**.

[0053] In some embodiments, processor **102** is coupled with a processor bus **110** to transmit communication signals such as address, data, or control signals between processor **102** and other components in system **100**. In one embodiment the system **100** uses an exemplary 'hub' system architecture, including a memory controller hub **116** and an Input Output (I/O) controller hub **130**. A memory controller hub **116** facilitates communication between a memory device and other components of system **100**, while an I/O Controller Hub (ICH) **130** provides connections to I/O devices via a local I/O bus. In one embodiment, the logic of the memory controller hub **116** is integrated within the processor.

[0054] Memory device **120** can be a dynamic random access memory (DRAM) device, a static random access memory (SRAM) device, flash memory device, phase-change memory device, or some other memory device having suitable performance to serve as process memory. In one embodiment the memory device **120** can operate as system memory for the system **100**, to store data **122** and instructions **121** for use when the one or more processors **102** executes an application or process. Memory controller hub **116** also couples with an optional external graphics processor **112**, which may communicate with the one or more graphics processors **108** in processors **102** to perform graphics and media operations.

[0055] In some embodiments, ICH **130** enables peripherals to connect to memory device **120** and processor **102** via a high-speed I/O bus. The I/O peripherals include, but are not limited to, an audio controller **146**, a firmware interface **128**, a wireless transceiver **126** (e.g., Wi-Fi, Bluetooth), a data storage device **124** (e.g., hard disk drive, flash memory, etc.), and a legacy I/O controller **140** for coupling legacy (e.g., Personal System 2 (PS/2)) devices to the system. One or more Universal Serial Bus (USB) controllers **142** connect input devices, such as keyboard and mouse **144** combinations. A network controller **134** may also couple with ICH **130**. In some embodiments, a high-performance network controller (not shown) couples with processor bus **110**. It will be appreciated that the system **100** shown is exemplary and not limiting, as other types of data processing systems that are differently configured may also be used. For example, the I/O controller hub **130** may be integrated within the one or more processor **102**, or the memory controller hub **116** and I/O controller hub **130** may be integrated into a discreet external graphics processor, such as the external graphics processor **112**.

[0056] FIG. **5** is a block diagram of an embodiment of a processor **200** having one or more processor cores **202**A-**202**N, an integrated memory controller **214**, and an integrated graphics processor **208**. Those elements of FIG. **5** having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such. Processor **200** can include additional cores up to and including additional core **202**N represented by the dashed lined boxes. Each of processor cores **202**A-**202**N includes one or more internal cache units **204**A-**204**N. In some embodiments each processor core also has access to one or more shared cached units **206**.

[0057] The internal cache units **204**A-**204**N and shared cache units **206** represent a cache memory hierarchy within the processor **200**. The cache memory hierarchy may include at least one level of instruction and data cache within each processor core and one or more levels of shared mid-level cache, such as a Level 2 (L2), Level 3 (L3), Level 4 (L4), or other levels of cache, where the highest level of cache before external memory is classified as the LLC. In some embodiments, cache coherency logic maintains coherency between the various cache units **206** and **204**A-**204**N.

[0058] In some embodiments, processor **200** may also include a set of one or more bus controller units **216** and a system agent core **210**. The one or more bus controller units **216** manage a set of peripheral buses, such as one or more Peripheral Component Interconnect buses (e.g., PCI, PCI Express). System agent core **210** provides management functionality for the various processor components. In some embodiments, system agent core **210** includes one or more integrated memory controllers **214** to manage access to various external memory devices (not shown).

[0059] In some embodiments, one or more of the processor cores **202**A-**202**N include support for simultaneous multi-threading. In such embodiment, the system agent core **210** includes components for coordinating and operating cores **202**A-**202**N during multi-threaded processing. System agent core **210** may additionally include a power control unit (PCU), which includes logic and components to regulate the power state of processor cores **202**A-**202**N and graphics processor **208**.

[0060] In some embodiments, processor **200** additionally includes graphics processor **208** to execute graphics processing operations. In some embodiments, the graphics processor **208** couples with the set of shared cache units **206**, and the system agent core **210**, including the one or more integrated memory controllers **214**. In some embodiments, a display controller **211** is coupled with the graphics processor **208** to drive graphics processor output to one or more coupled displays. In some embodiments, display controller **211** may be a separate module coupled with the graphics processor via at least one interconnect, or may be integrated within the graphics processor **208** or system agent core **210**.

[0061] In some embodiments, a ring based interconnect unit **212** is used to couple the internal components of the processor **200**. However, an alternative interconnect unit may be used, such as a point-to-point interconnect, a switched interconnect, or other techniques, including techniques well known in the art. In some embodiments, graphics processor **208** couples with the ring interconnect **212** via an I/O link **213**.

[0062] The exemplary I/O link **213** represents at least one of multiple varieties of I/O interconnects, including an on package I/O interconnect which facilitates communication between various processor components and a high-performance embedded memory module **218**, such as an eDRAM module. In some embodiments, each of the processor cores **202**A-**202**N and graphics processor **208** use embedded memory modules **218** as a shared Last Level Cache.

[0063] In some embodiments, processor cores **202**A-**202**N are homogenous cores executing the same instruction set architecture. In another embodiment, processor cores **202**A-**202**N are heterogeneous in terms of instruction set architecture (ISA), where one or more of processor cores **202**A-**202**N execute a first instruction set, while at least one of the other cores executes a subset of the first instruction set or a different instruction set. In one embodiment processor cores **202**A-**202**N are heterogeneous in terms of microarchitecture, where one or more cores having a relatively higher power consumption couple with one or more power cores having a lower power consumption. Additionally, processor **200** can be implemented on one or more chips or as an SoC integrated circuit having the illustrated components, in addition to other components.

[0064] FIG. **6** is a block diagram of a graphics processor **300**, which may be a discrete graphics processing unit, or may be a graphics processor integrated with a plurality of processing cores. In some embodiments, the graphics processor communicates via a memory mapped I/O interface to registers on the graphics processor and with commands placed into the processor memory. In some embodiments, graphics processor **300** includes a memory interface **314** to access memory. Memory interface **314** can be an interface to local memory, one or more internal caches, one or more shared external caches, and/or to system memory.

[0065] In some embodiments, graphics processor **300** also includes a display controller **302** to drive display output data to a display device **320**. Display controller **302** includes hardware for one or more overlay planes for the display and composition of multiple layers of video or user interface elements. In some embodiments, graphics processor **300** includes a video codec engine **306** to encode, decode, or transcode media to, from, or between one or more media encoding formats, including, but not limited to Moving Picture Experts Group (MPEG) formats such as MPEG-2, Advanced Video Coding (AVC) formats such as H.264/MPEG-4 AVC, as well as the Society of Motion Picture & Television Engineers (SMPTE) 421M/VC-1, and Joint Photographic Experts Group (JPEG) formats such as JPEG, and Motion JPEG (MJPEG) formats.

[0066] In some embodiments, graphics processor **300** includes a block image transfer (BLIT) engine **304** to perform two-dimensional (2D) rasterizer operations including, for example, bit-boundary block transfers. However, in one embodiment, 2D graphics operations are performed using one or more components of graphics processing engine (GPE) **310**. In some embodiments, GPE **310** is a compute engine for performing graphics operations, including three-dimensional (3D) graphics operations and media operations.

[0067] In some embodiments, GPE **310** includes a 3D pipeline **312** for performing 3D operations, such as rendering three-dimensional images and scenes using processing functions that act upon 3D primitive shapes (e.g., rectangle, triangle, etc.). The 3D pipeline **312** includes programmable and fixed function elements that perform various tasks within the element and/or spawn execution threads to a 3D/Media sub-system **315**. While 3D pipeline **312** can be used to perform media operations, an embodiment of GPE **310** also includes a media pipeline **316** that is specifically used to perform media operations, such as video post-processing and image enhancement.

[0068] In some embodiments, media pipeline **316** includes fixed function or programmable logic units to perform one or more specialized media operations, such as video decode acceleration, video de-interlacing, and video encode acceleration in place of, or on behalf of video codec engine **306**. In some embodiments, media pipeline **316** additionally includes a thread spawning unit to spawn threads for execution on 3D/Media sub-system **315**. The spawned threads perform computations for the media operations on one or more graphics execution units included in 3D/Media sub-system **315**.

[0069] In some embodiments, 3D/Media subsystem **315** includes logic for executing threads spawned by 3D pipeline **312** and media pipeline **316**. In one embodiment, the pipelines send thread execution requests to 3D/Media subsystem **315**, which includes thread dispatch logic for arbitrating and dispatching the various requests to available thread execution resources. The execution resources include an array of

graphics execution units to process the 3D and media threads. In some embodiments, 3D/Media subsystem **315** includes one or more internal caches for thread instructions and data. In some embodiments, the subsystem also includes shared memory, including registers and addressable memory, to share data between threads and to store output data.

[0070] FIG. **7** is a block diagram of a graphics processing engine **410** of a graphics processor in accordance with some embodiments. In one embodiment, the graphics processing engine (GPE) **410** is a version of the GPE **310** shown in FIG. **7**. Elements of FIG. **7** having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such. For example, the 3D pipeline **312** and media pipeline **316** of FIG. **6** are illustrated. The media pipeline **316** is optional in some embodiments of the GPE **410** and may not be explicitly included within the GPE **410**. For example and in at least one embodiment, a separate media and/or image processor is coupled to the GPE **410**.

[0071] In some embodiments, GPE **410** couples with or includes a command streamer **403**, which provides a command stream to the 3D pipeline **312** and/or media pipelines **316**. In some embodiments, command streamer **403** is coupled with memory, which can be system memory, or one or more of internal cache memory and shared cache memory. In some embodiments, command streamer **403** receives commands from the memory and sends the commands to 3D pipeline **312** and/or media pipeline **316**. The commands are directives fetched from a ring buffer, which stores commands for the 3D pipeline **312** and media pipeline **316**. In one embodiment, the ring buffer can additionally include batch command buffers storing batches of multiple commands. The commands for the 3D pipeline **312** can also include references to data stored in memory, such as but not limited to vertex and geometry data for the 3D pipeline **312** and/or image data and memory objects for the media pipeline **316**. The 3D pipeline **312** and media pipeline **316** process the commands and data by performing operations via logic within the respective pipelines or by dispatching one or more execution threads to a graphics core array **414**.

[0072] In various embodiments the 3D pipeline **312** can execute one or more shader programs, such as vertex shaders, geometry shaders, pixel shaders, fragment shaders, compute shaders, or other shader programs, by processing the instructions and dispatching execution threads to the graphics core array **414**. The graphics core array **414** provides a unified block of execution resources. Multi-purpose execution logic (e.g., execution units) within the graphic core array **414** includes support for various 3D API shader languages and can execute multiple simultaneous execution threads associated with multiple shaders.

[0073] In some embodiments the graphics core array **414** also includes execution logic to perform media functions, such as video and/or image processing. In one embodiment, the execution units additionally include general-purpose logic that is programmable to perform parallel general purpose computational operations, in addition to graphics processing operations. The general purpose logic can perform processing operations in parallel or in conjunction with general purpose logic within the processor core(s) **107** of FIG. **4** or core **202**A-**202**N as in FIG. **5**.

[0074] Output data generated by threads executing on the graphics core array **414** can output data to memory in a unified return buffer (URB) **418**. The URB **418** can store data for multiple threads. In some embodiments the URB **418** may be used to send data between different threads executing on the graphics core array **414**. In some embodiments the URB **418** may additionally be used for synchronization between threads on the graphics core array and fixed function logic within the shared function logic **420**.

[0075] In some embodiments, graphics core array **414** is scalable, such that the array includes a variable number of graphics cores, each having a variable number of execution units based on the target power and performance level of GPE **410**. In one embodiment the execution resources are dynamically scalable, such that execution resources may be enabled or disabled as needed.

[0076] The graphics core array **414** couples with shared function logic **420** that includes multiple resources that are shared between the graphics cores in the graphics core array. The shared functions within the shared function logic **420** are hardware logic units that provide specialized supplemental functionality to the graphics core array **414**.

[0077] In various embodiments, shared function logic **420** includes but is not limited to sampler **421**, math **422**, and inter-thread communication (ITC) **423** logic. Additionally, some embodiments implement one or more cache(s) **425** within the shared function logic **420**. A shared function is implemented where the demand for a given specialized function is insufficient for inclusion within the graphics core array **414**. Instead a single instantiation of that specialized function is implemented as a stand-alone entity in the shared function logic **420** and shared among the execution resources within the graphics core array **414**. The precise set of functions that are shared between the graphics core array **414** and included within the graphics core array **414** varies between embodiments.

[0078] FIG. **8** is a block diagram of another embodiment of a graphics processor **500**. Elements of FIG. **8** having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

[0079] In some embodiments, graphics processor **500** includes a ring interconnect **502**, a pipeline front-end **504**, a media engine **537**, and graphics cores **580**A-**580**N. In some embodiments, ring interconnect **502** couples the graphics processor to other processing units, including other graphics processors or one or more general-purpose processor cores. In some embodiments, the graphics processor is one of many processors integrated within a multi-core processing system.

[0080] In some embodiments, graphics processor **500** receives batches of commands via ring interconnect **502**. The incoming commands are interpreted by a command streamer **503** in the pipeline front-end **504**. In some embodiments, graphics processor **500** includes scalable execution logic to perform 3D geometry processing and media processing via the graphics core(s) **580**A-**580**N. For 3D geometry processing commands, command streamer **503** supplies commands to geometry pipeline **536**. For at least some media processing commands, command streamer **503** supplies the commands to a video front end **534**, which couples with a media engine **537**. In some embodiments, media engine **537** includes a Video Quality Engine (VQE) **530** for video and image post-processing and a multi-format encode/

decode (MFX) **533** engine to provide hardware-accelerated media data encode and decode. In some embodiments, geometry pipeline **536** and media engine **537** each generate execution threads for the thread execution resources provided by at least one graphics core **580**A.

[0081] In some embodiments, graphics processor **500** includes scalable thread execution resources featuring modular cores **580**A-**580**N (sometimes referred to as core slices), each having multiple sub-cores **550**A-**550**N, **560**A-**560**N (sometimes referred to as core sub-slices). In some embodiments, graphics processor **500** can have any number of graphics cores **580**A through **580**N. In some embodiments, graphics processor **500** includes a graphics core **580**A having at least a first sub-core **550**A and a second sub-core **560**A. In other embodiments, the graphics processor is a low power processor with a single sub-core (e.g., **550**A). In some embodiments, graphics processor **500** includes multiple graphics cores **580**A-**580**N, each including a set of first sub-cores **550**A-**550**N and a set of second sub-cores **560**A-**560**N. Each sub-core in the set of first sub-cores **550**A-**550**N includes at least a first set of execution units **552**A-**552**N and media/texture samplers **554**A-**554**N. Each sub-core in the set of second sub-cores **560**A-**560**N includes at least a second set of execution units **562**A-**562**N and samplers **564**A-**564**N. In some embodiments, each sub-core **550**A-**550**N, **560**A-**560**N shares a set of shared resources **570**A-**570**N. In some embodiments, the shared resources include shared cache memory and pixel operation logic. Other shared resources may also be included in the various embodiments of the graphics processor.

[0082] FIG. **9** illustrates thread execution logic **600** including an array of processing elements employed in some embodiments of a GPE. Elements of FIG. **9** having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

[0083] In some embodiments, thread execution logic **600** includes a shader processor **602**, a thread dispatcher **604**, instruction cache **606**, a scalable execution unit array including a plurality of execution units **608**A-**608**N, a sampler **610**, a data cache **612**, and a data port **614**. In one embodiment the scalable execution unit array can dynamically scale by enabling or disabling one or more execution units (e.g., any of execution unit **608**A, **608**B, **608**C, **608**D, through **608**N-**1** and **608**N) based on the computational requirements of a workload. In one embodiment the included components are interconnected via an interconnect fabric that links to each of the components. In some embodiments, thread execution logic **600** includes one or more connections to memory, such as system memory or cache memory, through one or more of instruction cache **606**, data port **614**, sampler **610**, and execution units **608**A-**608**N. In some embodiments, each execution unit (e.g. **608**A) is a stand-alone programmable general purpose computational unit that is capable of executing multiple simultaneous hardware threads while processing multiple data elements in parallel for each thread. In various embodiments, the array of execution units **608**A-**608**N is scalable to include any number individual execution units.

[0084] In some embodiments, the execution units **608**A-**608**N are primarily used to execute shader programs. A shader processor **602** can process the various shader programs and dispatch execution threads associated with the shader programs via a thread dispatcher **604**. In one embodiment the thread dispatcher includes logic to arbitrate thread initiation requests from the graphics and media pipelines and instantiate the requested threads on one or more execution unit in the execution units **608**A-**608**N. For example, the geometry pipeline (e.g., **536** of FIG. **8**) can dispatch vertex, tessellation, or geometry shaders to the thread execution logic **600** (FIG. **9**) for processing. In some embodiments, thread dispatcher **604** can also process runtime thread spawning requests from the executing shader programs.

[0085] In some embodiments, the execution units **608**A-**608**N support an instruction set that includes native support for many standard 3D graphics shader instructions, such that shader programs from graphics libraries (e.g., Direct 3D and OpenGL) are executed with a minimal translation. The execution units support vertex and geometry processing (e.g., vertex programs, geometry programs, vertex shaders), pixel processing (e.g., pixel shaders, fragment shaders) and general-purpose processing (e.g., compute and media shaders). Each of the execution units **608**A-**608**N is capable of multi-issue single instruction multiple data (SIMD) execution and multi-threaded operation enables an efficient execution environment in the face of higher latency memory accesses. Each hardware thread within each execution unit has a dedicated high-bandwidth register file and associated independent thread-state. Execution is multi-issue per clock to pipelines capable of integer, single and double precision floating point operations, SIMD branch capability, logical operations, transcendental operations, and other miscellaneous operations. While waiting for data from memory or one of the shared functions, dependency logic within the execution units **608**A-**608**N causes a waiting thread to sleep until the requested data has been returned. While the waiting thread is sleeping, hardware resources may be devoted to processing other threads. For example, during a delay associated with a vertex shader operation, an execution unit can perform operations for a pixel shader, fragment shader, or another type of shader program, including a different vertex shader.

[0086] Each execution unit in execution units **608**A-**608**N operates on arrays of data elements. The number of data elements is the "execution size," or the number of channels for the instruction. An execution channel is a logical unit of execution for data element access, masking, and flow control within instructions. The number of channels may be independent of the number of physical Arithmetic Logic Units (ALUs) or Floating Point Units (FPUs) for a particular graphics processor. In some embodiments, execution units **608**A-**608**N support integer and floating-point data types.

[0087] The execution unit instruction set includes SIMD instructions. The various data elements can be stored as a packed data type in a register and the execution unit will process the various elements based on the data size of the elements. For example, when operating on a 256-bit wide vector, the 256 bits of the vector are stored in a register and the execution unit operates on the vector as four separate 64-bit packed data elements (Quad-Word (QW) size data elements), eight separate 32-bit packed data elements (Double Word (DW) size data elements), sixteen separate 16-bit packed data elements (Word (W) size data elements), or thirty-two separate 8-bit data elements (byte (B) size data elements). However, different vector widths and register sizes are possible.

[0088] One or more internal instruction caches (e.g., **606**) are included in the thread execution logic **600** to cache thread instructions for the execution units. In some embodiments, one or more data caches (e.g., **612**) are included to cache thread data during thread execution. In some embodiments, a sampler **610** is included to provide texture sampling for 3D operations and media sampling for media operations. In some embodiments, sampler **610** includes specialized texture or media sampling functionality to process texture or media data during the sampling process before providing the sampled data to an execution unit.

[0089] During execution, the graphics and media pipelines send thread initiation requests to thread execution logic **600** via thread spawning and dispatch logic. Once a group of geometric objects has been processed and rasterized into pixel data, pixel processor logic (e.g., pixel shader logic, fragment shader logic, etc.) within the shader processor **602** is invoked to further compute output information and cause results to be written to output surfaces (e.g., color buffers, depth buffers, stencil buffers, etc.). In some embodiments, a pixel shader or fragment shader calculates the values of the various vertex attributes that are to be interpolated across the rasterized object. In some embodiments, pixel processor logic within the shader processor **602** then executes an application programming interface (API)-supplied pixel or fragment shader program. To execute the shader program, the shader processor **602** dispatches threads to an execution unit (e.g., **608**A) via thread dispatcher **604**. In some embodiments, pixel shader **602** uses texture sampling logic in the sampler **610** to access texture data in texture maps stored in memory. Arithmetic operations on the texture data and the input geometry data compute pixel color data for each geometric fragment, or discards one or more pixels from further processing.

[0090] In some embodiments, the data port **614** provides a memory access mechanism for the thread execution logic **600** output processed data to memory for processing on a graphics processor output pipeline. In some embodiments, the data port **614** includes or couples to one or more cache memories (e.g., data cache **612**) to cache data for memory access via the data port.

[0091] FIG. **10** is a block diagram illustrating a graphics processor instruction formats **700** according to some embodiments. In one or more embodiment, the graphics processor execution units support an instruction set having instructions in multiple formats. The solid lined boxes illustrate the components that are generally included in an execution unit instruction, while the dashed lines include components that are optional or that are only included in a sub-set of the instructions. In some embodiments, instruction format **700** described and illustrated are macro-instructions, in that they are instructions supplied to the execution unit, as opposed to micro-operations resulting from instruction decode once the instruction is processed.

[0092] In some embodiments, the graphics processor execution units natively support instructions in a 128-bit instruction format **710**. A 64-bit compacted instruction format **730** is available for some instructions based on the selected instruction, instruction options, and number of operands. The native 128-bit instruction format **710** provides access to all instruction options, while some options and operations are restricted in the 64-bit instruction format **730**. The native instructions available in the 64-bit instruction format **730** vary by embodiment. In some embodiments, the instruction is compacted in part using a set of index values in an index field **713**. The execution unit hardware references a set of compaction tables based on the index values and uses the compaction table outputs to reconstruct a native instruction in the 128-bit instruction format **710**.

[0093] For each format, instruction opcode **712** defines the operation that the execution unit is to perform. The execution units execute each instruction in parallel across the multiple data elements of each operand. For example, in response to an add instruction the execution unit performs a simultaneous add operation across each color channel representing a texture element or picture element. By default, the execution unit performs each instruction across all data channels of the operands. In some embodiments, instruction control field **714** enables control over certain execution options, such as channels selection (e.g., predication) and data channel order (e.g., swizzle). For instructions in the 128-bit instruction format **710** an exec-size field **716** limits the number of data channels that will be executed in parallel. In some embodiments, exec-size field **716** is not available for use in the 64-bit compact instruction format **730**.

[0094] Some execution unit instructions have up to three operands including two source operands, src0 **720**, src1 **722**, and one destination **718**. In some embodiments, the execution units support dual destination instructions, where one of the destinations is implied. Data manipulation instructions can have a third source operand (e.g., SRC2 **724**), where the instruction opcode **712** determines the number of source operands. An instruction's last source operand can be an immediate (e.g., hard-coded) value passed with the instruction.

[0095] In some embodiments, the 128-bit instruction format **710** includes an access/address mode field **726** specifying, for example, whether direct register addressing mode or indirect register addressing mode is used. When direct register addressing mode is used, the register address of one or more operands is directly provided by bits in the instruction.

[0096] In some embodiments, the 128-bit instruction format **710** includes an access/address mode field **726**, which specifies an address mode and/or an access mode for the instruction. In one embodiment the access mode is used to define a data access alignment for the instruction. Some embodiments support access modes including a 16-byte aligned access mode and a 1-byte aligned access mode, where the byte alignment of the access mode determines the access alignment of the instruction operands. For example, when in a first mode, the instruction may use byte-aligned addressing for source and destination operands and when in a second mode, the instruction may use 16-byte-aligned addressing for all source and destination operands.

[0097] In one embodiment, the address mode portion of the access/address mode field **726** determines whether the instruction is to use direct or indirect addressing. When direct register addressing mode is used bits in the instruction directly provide the register address of one or more operands. When indirect register addressing mode is used, the register address of one or more operands may be computed based on an address register value and an address immediate field in the instruction.

[0098] In some embodiments instructions are grouped based on opcode **712** bit-fields to simplify Opcode decode **740**. For an 8-bit opcode, bits **4**, **5**, and **6** allow the execution unit to determine the type of opcode. The precise opcode

grouping shown is merely an example. In some embodiments, a move and logic opcode group **742** includes data movement and logic instructions (e.g., move (mov), compare (cmp)). In some embodiments, move and logic group **742** shares the five most significant bits (MSB), where move (mov) instructions are in the form of 0000xxxxb and logic instructions are in the form of 0001xxxxb. A flow control instruction group **744** (e.g., call, jump (jmp)) includes instructions in the form of 0010xxxxb (e.g., 0x20). A miscellaneous instruction group **746** includes a mix of instructions, including synchronization instructions (e.g., wait, send) in the form of 0011xxxxb (e.g., 0x30). A parallel math instruction group **748** includes component-wise arithmetic instructions (e.g., add, multiply (mul)) in the form of 0100xxxxb (e.g., 0x40). The parallel math group **748** performs the arithmetic operations in parallel across data channels. The vector math group **750** includes arithmetic instructions (e.g., dp4) in the form of 0101xxxxb (e.g., 0x50). The vector math group performs arithmetic such as dot product calculations on vector operands.

[0099] FIG. **11** is a block diagram of another embodiment of a graphics processor **800**. Elements of FIG. **11** having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

[0100] In some embodiments, graphics processor **800** includes a graphics pipeline **820**, a media pipeline **830**, a display engine **840**, thread execution logic **850**, and a render output pipeline **870**. In some embodiments, graphics processor **800** is a graphics processor within a multi-core processing system that includes one or more general purpose processing cores. The graphics processor is controlled by register writes to one or more control registers (not shown) or via commands issued to graphics processor **800** via a ring interconnect **802**. In some embodiments, ring interconnect **802** couples graphics processor **800** to other processing components, such as other graphics processors or general-purpose processors. Commands from ring interconnect **802** are interpreted by a command streamer **803**, which supplies instructions to individual components of graphics pipeline **820** or media pipeline **830**.

[0101] In some embodiments, command streamer **803** directs the operation of a vertex fetcher **805** that reads vertex data from memory and executes vertex-processing commands provided by command streamer **803**. In some embodiments, vertex fetcher **805** provides vertex data to a vertex shader **807**, which performs coordinate space transformation and lighting operations to each vertex. In some embodiments, vertex fetcher **805** and vertex shader **807** execute vertex-processing instructions by dispatching execution threads to execution units **852**A-**852**B via a thread dispatcher **831**.

[0102] In some embodiments, execution units **852**A-**852**B are an array of vector processors having an instruction set for performing graphics and media operations. In some embodiments, execution units **852**A-**852**B have an attached L1 cache **851** that is specific for each array or shared between the arrays. The cache can be configured as a data cache, an instruction cache, or a single cache that is partitioned to contain data and instructions in different partitions.

[0103] In some embodiments, graphics pipeline **820** includes tessellation components to perform hardware-accelerated tessellation of 3D objects. In some embodiments, a programmable hull shader **811** configures the tessellation operations. A programmable domain shader **817** provides back-end evaluation of tessellation output. A tessellator **813** operates at the direction of hull shader **811** and contains special purpose logic to generate a set of detailed geometric objects based on a coarse geometric model that is provided as input to graphics pipeline **820**. In some embodiments, if tessellation is not used, tessellation components (e.g., hull shader **811**, tessellator **813**, and domain shader **817**) can be bypassed.

[0104] In some embodiments, complete geometric objects can be processed by a geometry shader **819** via one or more threads dispatched to execution units **852**A-**852**B, or can proceed directly to the clipper **829**. In some embodiments, the geometry shader operates on entire geometric objects, rather than vertices or patches of vertices as in previous stages of the graphics pipeline. If the tessellation is disabled the geometry shader **819** receives input from the vertex shader **807**. In some embodiments, geometry shader **819** is programmable by a geometry shader program to perform geometry tessellation if the tessellation units are disabled.

[0105] Before rasterization, a clipper **829** processes vertex data. The clipper **829** may be a fixed function clipper or a programmable clipper having clipping and geometry shader functions. In some embodiments, a rasterizer and depth test component **873** in the render output pipeline **870** dispatches pixel shaders to convert the geometric objects into their per pixel representations. In some embodiments, pixel shader logic is included in thread execution logic **850**. In some embodiments, an application can bypass the rasterizer and depth test component **873** and access un-rasterized vertex data via a stream out unit **823**.

[0106] The graphics processor **800** has an interconnect bus, interconnect fabric, or some other interconnect mechanism that allows data and message passing amongst the major components of the processor. In some embodiments, execution units **852**A-**852**B and associated cache(s) **851**, texture and media sampler **854**, and texture/sampler cache **858** interconnect via a data port **856** to perform memory access and communicate with render output pipeline components of the processor. In some embodiments, sampler **854**, caches **851**, **858** and execution units **852**A-**852**B each have separate memory access paths.

[0107] In some embodiments, render output pipeline **870** contains a rasterizer and depth test component **873** that converts vertex-based objects into an associated pixel-based representation. In some embodiments, the rasterizer logic includes a windower/masker unit to perform fixed function triangle and line rasterization. An associated render cache **878** and depth cache **879** are also available in some embodiments. A pixel operations component **877** performs pixel-based operations on the data, though in some instances, pixel operations associated with 2D operations (e.g. bit block image transfers with blending) are performed by the 2D engine **841**, or substituted at display time by the display controller **843** using overlay display planes. In some embodiments, a shared L3 cache **875** is available to all graphics components, allowing the sharing of data without the use of main system memory.

[0108] In some embodiments, graphics processor media pipeline **830** includes a media engine **837** and a video front end **834**. In some embodiments, video front end **834** receives pipeline commands from the command streamer **803**. In some embodiments, media pipeline **830** includes a separate

command streamer. In some embodiments, video front-end **834** processes media commands before sending the command to the media engine **837**. In some embodiments, media engine **837** includes thread spawning functionality to spawn threads for dispatch to thread execution logic **850** via thread dispatcher **831**.

[0109] In some embodiments, graphics processor **800** includes a display engine **840**. In some embodiments, display engine **840** is external to processor **800** and couples with the graphics processor via the ring interconnect **802**, or some other interconnect bus or fabric. In some embodiments, display engine **840** includes a 2D engine **841** and a display controller **843**. In some embodiments, display engine **840** contains special purpose logic capable of operating independently of the 3D pipeline. In some embodiments, display controller **843** couples with a display device (not shown), which may be a system integrated display device, as in a laptop computer, or an external display device attached via a display device connector.

[0110] In some embodiments, graphics pipeline **820** and media pipeline **830** are configurable to perform operations based on multiple graphics and media programming interfaces and are not specific to any one application programming interface (API). In some embodiments, driver software for the graphics processor translates API calls that are specific to a particular graphics or media library into commands that can be processed by the graphics processor. In some embodiments, support is provided for the Open Graphics Library (OpenGL), Open Computing Language (OpenCL), and/or Vulkan graphics and compute API, all from the Khronos Group. In some embodiments, support may also be provided for the Direct3D library from the Microsoft Corporation. In some embodiments, a combination of these libraries may be supported. Support may also be provided for the Open Source Computer Vision Library (OpenCV). A future API with a compatible 3D pipeline would also be supported if a mapping can be made from the pipeline of the future API to the pipeline of the graphics processor.

[0111] FIG. **12**A is a block diagram illustrating a graphics processor command format **900** according to some embodiments. FIG. **12**B is a block diagram illustrating a graphics processor command sequence **910** according to an embodiment. The solid lined boxes in FIG. **12**A illustrate the components that are generally included in a graphics command while the dashed lines include components that are optional or that are only included in a sub-set of the graphics commands. The exemplary graphics processor command format **900** of FIG. **12**A includes data fields to identify a target client **902** of the command, a command operation code (opcode) **904**, and the relevant data **906** for the command. A sub-opcode **905** and a command size **908** are also included in some commands.

[0112] In some embodiments, client **902** specifies the client unit of the graphics device that processes the command data. In some embodiments, a graphics processor command parser examines the client field of each command to condition the further processing of the command and route the command data to the appropriate client unit. In some embodiments, the graphics processor client units include a memory interface unit, a render unit, a 2D unit, a 3D unit, and a media unit. Each client unit has a corresponding processing pipeline that processes the commands. Once the command is received by the client unit, the client unit reads the opcode **904** and, if present, sub-opcode **905** to determine the operation to perform. The client unit performs the command using information in data field **906**. For some commands an explicit command size **908** is expected to specify the size of the command. In some embodiments, the command parser automatically determines the size of at least some of the commands based on the command opcode. In some embodiments commands are aligned via multiples of a double word.

[0113] The flow diagram in FIG. **12**B shows an exemplary graphics processor command sequence **910**. In some embodiments, software or firmware of a data processing system that features an embodiment of a graphics processor uses a version of the command sequence shown to set up, execute, and terminate a set of graphics operations. A sample command sequence is shown and described for purposes of example only as embodiments are not limited to these specific commands or to this command sequence. Moreover, the commands may be issued as batch of commands in a command sequence, such that the graphics processor will process the sequence of commands in at least partially concurrence.

[0114] In some embodiments, the graphics processor command sequence **910** may begin with a pipeline flush command **912** to cause any active graphics pipeline to complete the currently pending commands for the pipeline. In some embodiments, the 3D pipeline **922** and the media pipeline **924** do not operate concurrently. The pipeline flush is performed to cause the active graphics pipeline to complete any pending commands. In response to a pipeline flush, the command parser for the graphics processor will pause command processing until the active drawing engines complete pending operations and the relevant read caches are invalidated. Optionally, any data in the render cache that is marked 'dirty' can be flushed to memory. In some embodiments, pipeline flush command **912** can be used for pipeline synchronization or before placing the graphics processor into a low power state.

[0115] In some embodiments, a pipeline select command **913** is used when a command sequence requires the graphics processor to explicitly switch between pipelines. In some embodiments, a pipeline select command **913** is required only once within an execution context before issuing pipeline commands unless the context is to issue commands for both pipelines. In some embodiments, a pipeline flush command **912** is required immediately before a pipeline switch via the pipeline select command **913**.

[0116] In some embodiments, a pipeline control command **914** configures a graphics pipeline for operation and is used to program the 3D pipeline **922** and the media pipeline **924**. In some embodiments, pipeline control command **914** configures the pipeline state for the active pipeline. In one embodiment, the pipeline control command **914** is used for pipeline synchronization and to clear data from one or more cache memories within the active pipeline before processing a batch of commands.

[0117] In some embodiments, commands for the return buffer state **916** are used to configure a set of return buffers for the respective pipelines to write data. Some pipeline operations require the allocation, selection, or configuration of one or more return buffers into which the operations write intermediate data during processing. In some embodiments, the graphics processor also uses one or more return buffers to store output data and to perform cross thread communi-

cation. In some embodiments, configuring the return buffer state **916** includes selecting the size and number of return buffers to use for a set of pipeline operations.

[0118] The remaining commands in the command sequence differ based on the active pipeline for operations. Based on a pipeline determination **920**, the command sequence is tailored to the 3D pipeline **922** beginning with the 3D pipeline state **930** or the media pipeline **924** beginning at the media pipeline state **940**.

[0119] The commands to configure the 3D pipeline state **930** include 3D state setting commands for vertex buffer state, vertex element state, constant color state, depth buffer state, and other state variables that are to be configured before 3D primitive commands are processed. The values of these commands are determined at least in part based on the particular 3D API in use. In some embodiments, 3D pipeline state **930** commands are also able to selectively disable or bypass certain pipeline elements if those elements will not be used.

[0120] In some embodiments, 3D primitive **932** command is used to submit 3D primitives to be processed by the 3D pipeline. Commands and associated parameters that are passed to the graphics processor via the 3D primitive **932** command are forwarded to the vertex fetch function in the graphics pipeline. The vertex fetch function uses the 3D primitive **932** command data to generate vertex data structures. The vertex data structures are stored in one or more return buffers. In some embodiments, 3D primitive **932** command is used to perform vertex operations on 3D primitives via vertex shaders. To process vertex shaders, 3D pipeline **922** dispatches shader execution threads to graphics processor execution units.

[0121] In some embodiments, 3D pipeline **922** is triggered via an execute **934** command or event. In some embodiments, a register write triggers command execution. In some embodiments execution is triggered via a 'go' or 'kick' command in the command sequence. In one embodiment, command execution is triggered using a pipeline synchronization command to flush the command sequence through the graphics pipeline. The 3D pipeline will perform geometry processing for the 3D primitives. Once operations are complete, the resulting geometric objects are rasterized and the pixel engine colors the resulting pixels. Additional commands to control pixel shading and pixel back end operations may also be included for those operations.

[0122] In some embodiments, the graphics processor command sequence **910** follows the media pipeline **924** path when performing media operations. In general, the specific use and manner of programming for the media pipeline **924** depends on the media or compute operations to be performed. Specific media decode operations may be offloaded to the media pipeline during media decode. In some embodiments, the media pipeline can also be bypassed and media decode can be performed in whole or in part using resources provided by one or more general purpose processing cores. In one embodiment, the media pipeline also includes elements for general-purpose graphics processor unit (GPGPU) operations, where the graphics processor is used to perform SIMD vector operations using computational shader programs that are not explicitly related to the rendering of graphics primitives.

[0123] In some embodiments, media pipeline **924** is configured in a similar manner as the 3D pipeline **922**. A set of commands to configure the media pipeline state **940** are dispatched or placed into a command queue before the media object commands **942**. In some embodiments, commands for the media pipeline state **940** include data to configure the media pipeline elements that will be used to process the media objects. This includes data to configure the video decode and video encode logic within the media pipeline, such as encode or decode format. In some embodiments, commands for the media pipeline state **940** also support the use of one or more pointers to "indirect" state elements that contain a batch of state settings.

[0124] In some embodiments, media object commands **942** supply pointers to media objects for processing by the media pipeline. The media objects include memory buffers containing video data to be processed. In some embodiments, all media pipeline states must be valid before issuing a media object command **942**. Once the pipeline state is configured and media object commands **942** are queued, the media pipeline **924** is triggered via an execute command **944** or an equivalent execute event (e.g., register write). Output from media pipeline **924** may then be post processed by operations provided by the 3D pipeline **922** or the media pipeline **924**. In some embodiments, GPGPU operations are configured and executed in a similar manner as media operations.

[0125] FIG. **13** illustrates exemplary graphics software architecture for a data processing system **1000** according to some embodiments. In some embodiments, software architecture includes a 3D graphics application **1010**, an operating system **1020**, and at least one processor **1030**. In some embodiments, processor **1030** includes a graphics processor **1032** and one or more general-purpose processor core(s) **1034**. The graphics application **1010** and operating system **1020** each execute in the system memory **1050** of the data processing system.

[0126] In some embodiments, 3D graphics application **1010** contains one or more shader programs including shader instructions **1012**. The shader language instructions may be in a high-level shader language, such as the High Level Shader Language (HLSL) or the OpenGL Shader Language (GLSL). The application also includes executable instructions **1014** in a machine language suitable for execution by the general-purpose processor core **1034**. The application also includes graphics objects **1016** defined by vertex data.

[0127] In some embodiments, operating system **1020** is a Microsoft® Windows® operating system from the Microsoft Corporation, a proprietary UNIX-like operating system, or an open source UNIX-like operating system using a variant of the Linux kernel. The operating system **1020** can support a graphics API **1022** such as the Direct3D API, the OpenGL API, or the Vulkan API. When the Direct3D API is in use, the operating system **1020** uses a front-end shader compiler **1024** to compile any shader instructions **1012** in HLSL into a lower-level shader language. The compilation may be a just-in-time (JIT) compilation or the application can perform shader pre-compilation. In some embodiments, high-level shaders are compiled into low-level shaders during the compilation of the 3D graphics application **1010**. In some embodiments, the shader instructions **1012** are provided in an intermediate form, such as a version of the Standard Portable Intermediate Representation (SPIR) used by the Vulkan API.

[0128] In some embodiments, user mode graphics driver **1026** contains a back-end shader compiler **1027** to convert

the shader instructions **1012** into a hardware specific representation. When the OpenGL API is in use, shader instructions **1012** in the GLSL high-level language are passed to a user mode graphics driver **1026** for compilation. In some embodiments, user mode graphics driver **1026** uses operating system kernel mode functions **1028** to communicate with a kernel mode graphics driver **1029**. In some embodiments, kernel mode graphics driver **1029** communicates with graphics processor **1032** to dispatch commands and instructions.

[0129] One or more aspects of at least one embodiment may be implemented by representative code stored on a machine-readable medium which represents and/or defines logic within an integrated circuit such as a processor. For example, the machine-readable medium may include instructions which represent various logic within the processor. When read by a machine, the instructions may cause the machine to fabricate the logic to perform the techniques described herein. Such representations, known as "IP cores," are reusable units of logic for an integrated circuit that may be stored on a tangible, machine-readable medium as a hardware model that describes the structure of the integrated circuit. The hardware model may be supplied to various customers or manufacturing facilities, which load the hardware model on fabrication machines that manufacture the integrated circuit. The integrated circuit may be fabricated such that the circuit performs operations described in association with any of the embodiments described herein.

[0130] FIG. **14** is a block diagram illustrating an IP core development system **1100** that may be used to manufacture an integrated circuit to perform operations according to an embodiment. The IP core development system **1100** may be used to generate modular, re-usable designs that can be incorporated into a larger design or used to construct an entire integrated circuit (e.g., an SOC integrated circuit). A design facility **1130** can generate a software simulation **1110** of an IP core design in a high level programming language (e.g., C/C++). The software simulation **1110** can be used to design, test, and verify the behavior of the IP core using a simulation model **1112**. The simulation model **1112** may include functional, behavioral, and/or timing simulations. A register transfer level (RTL) design **1115** can then be created or synthesized from the simulation model **1112**. The RTL design **1115** is an abstraction of the behavior of the integrated circuit that models the flow of digital signals between hardware registers, including the associated logic performed using the modeled digital signals. In addition to an RTL design **1115**, lower-level designs at the logic level or transistor level may also be created, designed, or synthesized. Thus, the particular details of the initial design and simulation may vary.

[0131] The RTL design **1115** or equivalent may be further synthesized by the design facility into a hardware model **1120**, which may be in a hardware description language (HDL), or some other representation of physical design data. The HDL may be further simulated or tested to verify the IP core design. The IP core design can be stored for delivery to a 3rd party fabrication facility **1165** using non-volatile memory **1140** (e.g., hard disk, flash memory, or any non-volatile storage medium). Alternatively, the IP core design may be transmitted (e.g., via the Internet) over a wired connection **1150** or wireless connection **1160**. The fabrication facility **1165** may then fabricate an integrated circuit that is based at least in part on the IP core design. The fabricated integrated circuit can be configured to perform operations in accordance with at least one embodiment described herein.

[0132] FIGS. **15-17** illustrate exemplary integrated circuits and associated graphics processors that may be fabricated using one or more IP cores, according to various embodiments described herein. In addition to what is illustrated, other logic and circuits may be included, including additional graphics processors/cores, peripheral interface controllers, or general purpose processor cores.

[0133] FIG. **15** is a block diagram illustrating an exemplary system on a chip integrated circuit **1200** that may be fabricated using one or more IP cores, according to an embodiment. Exemplary integrated circuit **1200** includes one or more application processor(s) **1205** (e.g., CPUs), at least one graphics processor **1210**, and may additionally include an image processor **1215** and/or a video processor **1220**, any of which may be a modular IP core from the same or multiple different design facilities. Integrated circuit **1200** includes peripheral or bus logic including a USB controller **1225**, UART controller **1230**, an SPI/SDIO controller **1235**, and an I^2S/I^2C controller **1240**. Additionally, the integrated circuit can include a display device **1245** coupled to one or more of a high-definition multimedia interface (HDMI) controller **1250** and a mobile industry processor interface (MIPI) display interface **1255**. Storage may be provided by a flash memory subsystem **1260** including flash memory and a flash memory controller. Memory interface may be provided via a memory controller **1265** for access to SDRAM or SRAM memory devices. Some integrated circuits additionally include an embedded security engine **1270**.

[0134] FIG. **15** is a block diagram illustrating an exemplary graphics processor **1310** of a system on a chip integrated circuit that may be fabricated using one or more IP cores, according to an embodiment. Graphics processor **1310** can be a variant of the graphics processor **1210** of FIG. **15**. Graphics processor **1310** includes a vertex processor **1305** and one or more fragment processor(s) **1315**A**1315**N (e.g., **1315**A, **1315**B, **1315**C, **1315**D, through **1315**N-**1**, and **1315**N). Graphics processor **1310** can execute different shader programs via separate logic, such that the vertex processor **1305** is optimized to execute operations for vertex shader programs, while the one or more fragment processor (s) **1315**A-**1315**N execute fragment (e.g., pixel) shading operations for fragment or pixel shader programs. The vertex processor **1305** performs the vertex processing stage of the 3D graphics pipeline and generates primitives and vertex data. The fragment processor(s) **1315**A-**1315**N use the primitive and vertex data generated by the vertex processor **1305** to produce a framebuffer that is displayed on a display device. In one embodiment, the fragment processor (s) **1315**A-**1315**N are optimized to execute fragment shader programs as provided for in the OpenGL API, which may be used to perform similar operations as a pixel shader program as provided for in the Direct 3D API.

[0135] Graphics processor **1310** additionally includes one or more memory management units (MMUs) **1320**A-**1320**B, cache(s) **1325**A-**1325**B, and circuit interconnect(s) **1330**A-**1330**B. The one or more MMU(s) **1320**A-**1320**B provide for virtual to physical address mapping for graphics processor **1310**, including for the vertex processor **1305** and/or fragment processor(s) **1315**A-**1315**N, which may reference vertex or image/texture data stored in memory, in addition to vertex or image/texture data stored in the one or more

cache(s) **1325**A-**1325**B. In one embodiment the one or more MMU(s) **1320**A-**1320**B may be synchronized with other MMUs within the system, including one or more MMUs associated with the one or more application processor(s) **1205**, image processor **1215**, and/or video processor **1220** of FIG. **16**, such that each processor **1205**-**1220** can participate in a shared or unified virtual memory system. The one or more circuit interconnect(s) **1330**A-**1330**B enable graphics processor **1310** to interface with other IP cores within the SoC, either via an internal bus of the SoC or via a direct connection, according to embodiments.

[0136] FIG. **17** is a block diagram illustrating an additional exemplary graphics processor **1410** of a system on a chip integrated circuit that may be fabricated using one or more IP cores, according to an embodiment. Graphics processor **1410** can be a variant of the graphics processor **1210** of FIG. **15**. Graphics processor **1410** includes the one or more MMU(s) **1320**A-**1320**B, cache(s) **1325**A-**1325**B, and circuit interconnect(s) **1330**A-**1330**B of the integrated circuit **1300** of FIG. **16**.

[0137] Graphics processor **1410** includes one or more shader core(s) **1415**A-**1415**N (e.g., **1415**A, **1415**B, **1415**C, **1415**D, **1415**E, **1415**F, through **1315**N-**1**, and **1315**N), which provides for a unified shader core architecture in which a single core or type or core can execute all types of programmable shader code, including shader program code to implement vertex shaders, fragment shaders, and/or compute shaders. The exact number of shader cores present can vary among embodiments and implementations. Additionally, graphics processor **1410** includes an inter-core task manager **1405**, which acts as a thread dispatcher to dispatch execution threads to one or more shader core(s) **1415**A-**1415**N and a tiling unit **1418** to accelerate tiling operations for tile-based rendering, in which rendering operations for a scene are subdivided in image space, for example to exploit local spatial coherence within a scene or to optimize use of internal caches.

[0138] The following clauses and/or examples pertain to further embodiments:

[0139] One example embodiment may be a method comprising combining first and second kernels into a combined kernel by a compiler, receiving at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit, and offloading said combined kernel for execution on said graphics processing unit. The method may also include offloading execution of a first kernel using a ring task buffer with a fixed number of task slots, offloading execution of a second and all subsequent kernels using said ring task buffer, and offloading at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer. The method may also include resolving identification of said first and second and all subsequent kernels. The method may also include fetching parameters of said first and second and all subsequent kernels. The method may also include creating an object and writing said parameters and said identifications to said object. The method may also include blocking until a graphics processing unit completes a current task when no slots are available in the ring buffer. The method may also include starting a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit. The method may also include enabling users to decide when to engage offloading. The method may also include providing a mechanism to stop and start execution of the combined kernel.

[0140] In another example embodiment may be one or more non-transitory computer readable media storing instructions to perform a sequence comprising combining first and second kernels into a combined kernel by a compiler, receiving at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit, and offloading said combined kernel for execution on said graphics processing unit. The media may store instructions to perform a sequence including offloading execution of a first kernel using a ring task buffer with a fixed number of task slots, offloading execution of a second and all subsequent kernels using said ring task buffer, and offloading at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer. The media may store instructions to perform a sequence including resolving identification of said first and second and all subsequent kernels. The media may store instructions to perform a sequence including fetching parameters of said first and second and all subsequent kernels. The media may store instructions to perform a sequence including creating an object and writing said parameters and said identifications to said object. The media may store instructions to perform a sequence including blocking until a graphics processing unit completes a current task when no slots are available in the ring buffer. The media may store instructions to perform a sequence including starting a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit. The media may store instructions to perform a sequence including enabling users to decide when to engage offloading. The media may store instructions to perform a sequence including providing a mechanism to stop and start execution of the combined kernel.

[0141] Another example embodiment may be an apparatus a processor to combine first and second kernels into a combined kernel by a compiler, receive at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit, offload said combined kernel for execution on said graphics processing unit, and a memory coupled to said processor. The apparatus of said processor to offload execution of a first kernel using a ring task buffer with a fixed number of task slots, offload execution of a second and all subsequent kernels using said ring task buffer, and offload at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer. The apparatus of said processor to resolve identification of said first and second and all subsequent kernels. The apparatus of said processor to fetch parameters of said first and second and all subsequent kernels. The apparatus of said processor to create an object and writing said parameters and said identifications to said object. The apparatus of said processor to block until a graphics processing unit completes a current task when no slots are available in the ring buffer. The apparatus of said processor to start a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit. The apparatus of said processor to enable users to decide when to engage offloading. The apparatus of said processor to provide a mechanism to stop and start execution of the combined kernel.

[0142] The graphics processing techniques described herein may be implemented in various hardware architec-

tures. For example, graphics functionality may be integrated within a chipset. Alternatively, a discrete graphics processor may be used. As still another embodiment, the graphics functions may be implemented by a general purpose processor, including a multicore processor.

[0143] References throughout this specification to "one embodiment" or "an embodiment" mean that a particular feature, structure, or characteristic described in connection with the embodiment is included in at least one implementation encompassed within the present disclosure. Thus, appearances of the phrase "one embodiment" or "in an embodiment" are not necessarily referring to the same embodiment. Furthermore, the particular features, structures, or characteristics may be instituted in other suitable forms other than the particular embodiment illustrated and all such forms may be encompassed within the claims of the present application.

[0144] While a limited number of embodiments have been described, those skilled in the art will appreciate numerous modifications and variations therefrom. It is intended that the appended claims cover all such modifications and variations as fall within the true spirit and scope of this disclosure.

1. A method comprising:
combining first and second kernels into a combined kernel by a compiler;
receiving at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit; and
offloading said combined kernel for execution on said graphics processing unit.

2. The method of claim **1** further including:
offloading execution of a first kernel using a ring task buffer with a fixed number of task slots;
offloading execution of a second and all subsequent kernels using said ring task buffer; and
offloading at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer.

3. The method of claim **2** including resolving identification of said first and second and all subsequent kernels.

4. The method of claim **3** including fetching parameters of said first and second and all subsequent kernels.

5. The method of claim **4** including creating an object and writing said parameters and said identifications to said object.

6. The method of claim **5** including blocking until a graphics processing unit completes a current task when no slots are available in the ring buffer.

7. The method of claim **6** including starting a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit.

8. The method of claim **7** including enabling users to decide when to engage offloading.

9. The method of claim **8** including providing a mechanism to stop and start execution of the combined kernel.

10. One or more non-transitory computer readable media storing instructions to perform a sequence comprising:
combining first and second kernels into a combined kernel by a compiler;
receiving at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit; and
offloading said combined kernel for execution on said graphics processing unit.

11. The media of claim **10**, further storing instructions to perform a sequence including:
offloading execution of a first kernel using a ring task buffer with a fixed number of task slots;
offloading execution of a second and all subsequent kernels using said ring task buffer; and
offloading at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer.

12. The media of claim **11**, further storing instructions to perform a sequence including resolving identification of said first and second and all subsequent kernels.

13. The media of claim **12**, further storing instructions to perform a sequence including fetching parameters of said first and second and all subsequent kernels.

14. The media of claim **13**, further storing instructions to perform a sequence including creating an object and writing said parameters and said identifications to said object.

15. The media of claim **14**, further storing instructions to perform a sequence including blocking until a graphics processing unit completes a current task when no slots are available in the ring buffer.

16. The media of claim **15**, further storing instructions to perform a sequence including starting a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit.

17. The media of claim **16**, further storing instructions to perform a sequence including enabling users to decide when to engage offloading.

18. The media of claim **17**, further storing instructions to perform a sequence including providing a mechanism to stop and start execution of the combined kernel.

19. An apparatus comprising:
a processor to combine first and second kernels into a combined kernel by a compiler, receive at runtime, the combined kernel on a central processing unit for offloading to a graphics processing unit, offload said combined kernel for execution on said graphics processing unit; and
a memory coupled to said processor.

20. The apparatus of claim **19**, said processor to offload execution of a first kernel using a ring task buffer with a fixed number of task slots, offload execution of a second and all subsequent kernels using said ring task buffer, and offload at least two kernels from a central processing unit to a graphics processing unit via said ring task buffer.

21. The apparatus of claim **20**, said processor to resolve identification of said first and second and all subsequent kernels.

22. The apparatus of claim **21**, said processor to fetch parameters of said first and second and all subsequent kernels.

23. The apparatus of claim **22**, said processor to create an object and writing said parameters and said identifications to said object.

24. The apparatus of claim **23**, said processor to block until a graphics processing unit completes a current task when no slots are available in the ring buffer.

25. The apparatus of claim **24**, said processor to start a thread to periodically enqueue an exit task in said ring buffer to make the combined kernel finish and exit.

26. The apparatus of claim **25**, said processor to enable users to decide when to engage offloading.

27. The apparatus of claim **26**, said processor to provide a mechanism to stop and start execution of the combined kernel.

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