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#### ( 54 ) EFFICIENT SHARING AND COMPRESSION EXPANSION OF DATA ACROSS PROCESSING SYSTEMS

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#### ( 57 ) ABSTRACT

A mechanism is described for facilitating sharing of data and<br>compression expansion of models at autonomous machines.<br>A method of embodiments, as described herein, includes<br>detecting a first processor processing informatio a neural network at a first computing device , where the first computing device comprises a first autonomous machine. The method further includes facilitating the first processor to store one or more portions of the information in a library at a database, where the one or more portions are accessible to a second processor of a computing device .







FIG . 2A



## FIG . 2B



### FIG. 2C



FIG . 2D



## FIG. 3A



FIG. 3B



FIG. 4A









 $FIG. 4D$ 



FIG . 4E



#### $FIG.4F$



FIG. 5













FIG. 10







FIG. 13A



**FIG. 13B** 



FIG. 14





FIG .16







**C**<br>5







C N



E D.



Patent Application Publication

Figure



 $F1G.25$
# FIG. 26A GRAPHICS PROCESSOR COMMAND FORMAT 2600







## DATA PROCESSING SYSTEM - 2700

FIG. 27





FIG .28



FIG. 29



FIG 30



FIG. 31

### EFFICIENT SHARING AND COMPRESSION EXPANSION OF DATA ACROSS PROCESSING SYSTEMS

#### RELATED APPLICATION

[0001] This Application is a continuation of and claims the benefit of and priority to U.S. application Ser. No. 16/696, 852, entitled EFFICIENT SHARING AND COMPRES-SION EXPANSION OF DATA ACROSS PROCESSING<br>SYSTEMS, by Abhishek R. Appu, et al., filed Nov. 26, 2019, now allowed, which is a continuation of and claims the benefit of and priority to U.S. application Ser. No. 15/495, 081, entitled EFFICIENT SHARING AND COMPRES-SION EXPANSION OF DATA ACROSS PROCESSING<br>SYSTEMS, by Abhishek R. Appu, et al., filed Apr. 24, 2017, now issued as U.S. Pat. No. 10,497,084, the entire contents of which are incorporated herein by reference .

#### FIELD

data processing and more particularly to facilitate a tool for [0002] Embodiments described herein relate generally to data processing and more particularly to facilitate a tool for facilitating efficient sharing and compression expansion of data across processing systems.

#### BACKGROUND

[0003] Current parallel graphics data processing includes systems and methods developed to perform specific operations on graphics data such as, for example, linear interpolation, tessellation, rasterization, texture mapping, depth testing, etc. Traditionally, graphics processors used fixed function computational units to process graphics data; however, more recently, portions of graphics processors have been made programmable, enabling such processors to support a wider variety of operations for processing vertex and fragment data.<br>[0004] To further increase performance, graphics proces-

sors typically implement processing techniques such as pipelining that attempt to process, in parallel, as much graphics data as possible throughout the different parts of the graphics pipeline. Parallel graphics processors with single<br>instruction, multiple thread (SIMT) architectures are<br>designed to maximize the amount of parallel processing in the graphics pipeline. In an SIMT architecture, groups of parallel threads attempt to execute program instructions synchronously together as often as possible to increase processing efficiency. A general overview of software and hardware for SIMT architectures can be found in Shane Cook, CUDA Programming, Chapter 3, pages 37-51 (2013)<br>and/or Nicholas Wilt, CUDA Handbook, A Comprehensive<br>Guide to GPU Programming, Sections 2.6.2 to 3.1.2 (June<br>2013). instruction, multiple thread (SIMT) architectures are

[0005] Machine learning has been successful at solving many kinds of tasks. The computations that arise when training and using machine learning algorithms (e.g., neural networks) lend themselves naturally to efficient parallel implementations. Accordingly, parallel processors such as general-purpose graphic processing units (GPGPUs) have played a significant role in the practical implementation of deep neural networks. Parallel graphics processors with single instruction, multiple thread (SIMT) architectures are designed to maximize the amount of parallel processing in the graphics pipeline. In an SIMT architecture, groups of parallel threads attempt to execute program instructions synchronously together as often as possible to increase processing efficiency . The efficiency provided by parallel machine learning algorithm implementations allows the use of high capacity networks and enables those networks to be

[0006] Conventional techniques do not provide for data sharing across processing systems through a library, which is inefficient and cumbersome. Further, conventional techniques are limited to compression of data and thus do not anticipate expansion or re-expansion of compression models .

### BRIEF DESCRIPTION OF THE DRAWINGS

[0007] Embodiments are illustrated by way of example, and not by way of limitation, in the figures of the accompanying drawings in which like reference numerals refer to similar elements. So that the manner in which the above recited features can be understood in detail, a more particular description, briefly summarized above, may be had by reference to embodiments, some of which are illustrated in the appended drawings. It is to be noted, howe appended drawings illustrate only typical embodiments and are therefore not to be considered limiting of its scope, for the drawings may illustrate other equally effective embodi ments.

[0008] FIG. 1 is a block diagram illustrating a computer system configured to implement one or more aspects of the embodiments described herein.

embodiments described nerein.<br>
[0009] FIG. 2A-2D illustrate a parallel processor components, according to an embodiment.<br>
[0010] FIG. 3A-3B are block diagrams of graphics multiprocessors, according to embodiments.<br>
[0011]

smart resource distribution mechanism according to one

[0014] FIG. 7 illustrates smart resource distribution mechanism according to one embodiment.

[0015] FIG. 8A illustrates a tree-like communication structure for facilitating energy-efficient distribution of deep

learning according to one embodiment.<br>
[0016] FIG. 8B illustrates a process structure for facilitating effective communication in distributed deep learning

ing effective communication in distributed deep deficiency in energy, communication, and debugging at autonomous machines according to one embodiment.

[0018] FIG. 10 illustrates a machine learning software stack, according to an embodiment.

[0019] FIG. 11 illustrates a highly-parallel general-pur-<br>pose graphics processing unit, according to an embodiment. [0020] FIG. 12 illustrates a multi-GPU computing system.

according to an embodiment.<br>[0021] FIG. 13A-13B illustrate layers of exemplary deep<br>neural networks.

[0022] FIG. 14 illustrates training and deployment of a deep neural network.

[0023] FIG. 15 illustrates training and deployment of a deep neural network

[0024] FIG. 16 is a block diagram illustrating distributed<br>learning.<br>[0025] FIG. 17 illustrates an exemplary inferencing system on a chip (SOC) suitable for performing inferencing<br>using a trained model.

computer system with a processor having one or more processor cores and graphics processors. [0026] FIG. 18 is a block diagram of an embodiment of a

 $[0027]$  FIG. 19 is a block diagram of one embodiment of a processor having one or more processor cores, an integrated memory controller, and an integrated graphics processor.

[0028] FIG. 20 is a block diagram of one embodiment of a graphics processor which may be a discreet graphics a graphics processor which may be a discreet graphics<br>processing unit, or may be graphics processor integrated<br>with a plurality of processing cores.<br>[0029] FIG. 21 is a block diagram of an embodiment of a<br>graphics processi

unit instruction format according to an embodiment.<br>[0033] FIG. 25 is a block diagram of another embodiment of a graphics processor which includes a graphics pipeline, a media pipeline, a display engine, thread execution logic, and a render output pipeline. [0034] FIG. 26A is a block diagram illustrating a graphics processor com

[0035] FIG. 26B is a block diagram illustrating a graphics processor command sequence according to an embodiment. [0036] FIG. 27 illustrates exemplary graphics software architecture for a data processing system according to an embodiment.

[0037] FIG. 28 is a block diagram illustrating an IP core development system that may be used to manufacture an integrated circuit to perform operations according to an embodiment.

[0038] FIG. 29 is a block diagram illustrating an exemplary system on a chip integrated circuit that may be fabricated using one or more IP cores, according to an embodiment.

[0039] FIG. 30 is a block diagram illustrating an exemplary graphics processor of a system on a chip integrated circuit .

[ 0040 ] FIG . 31 is a block diagram illustrating an addi tional exemplary graphics processor of a system on a chip integrated circuit .

#### DETAILED DESCRIPTION

same convolution. Embodiments further provide for a novel [ 004 ] Embodiments provide for a novel technique for facilitating data sharing across processing systems using a surface library such that data produced on one graphics processor, application processor, etc., can be retrieved from the surface library and used by another if working on the same convolution. Embodiments further provide for a novel<br>technique for facilitating expansion or re-expansion of compressed models for performance and communication effi-<br>ciency.

 $[0042]$  it is to be noted that terms or acronyms like " convolutional neural network", "CNN", "neural network", "<br>"NN", "deep neural network", "DNN", "recurrent neural network", "RNN", and/or the like may be interchangeably referenced throughout this document. Further, terms like "autonomous machine" or simply "machine", "autonomous vehicle" or simply "vehicle", "autonomous agent" or simply "agent", "autonomous device" or "computing device", "robot", and/or the like, may be interchangeably referenced throughout this document. [0043] In some embodiment

 $(GPU)$  is communicatively coupled to host/processor cores to accelerate graphics operations, machine-learning operations, pattern analysis operations, and various general purpose GPU (GPGPU) functions. The GPU may be communicatively coupled to the host processor/cores over a bus or another interconnect (e.g., a high-speed interconnect such as PCIe or NVLink). In other embodiments, the GPU may be integrated on the same package or chip as the cores and communicatively coupled to the cores over an internal processor bus/interconnect (i.e., internal to the package or chip). Regardless of the manner in which the GPU is connected, the processor cores may allocate work to the GPU in the form of sequences of commands/instructions contained in a work descriptor. The GPU then uses dedicated circuitry/logic for efficiently processing these commands/ instructions .

[0044] In the following description, numerous specific details are set forth. However, embodiments, as described herein, may be practiced without these specific details. In other instances, well-known circuits, structures and techniques have not been shown in detail in order not to obscure the understanding of this description.

[0045] System Overview I

[0046] FIG. 1 is a block diagram illustrating a computing system 100 configured to implement one or more aspects of the embodiments described herein. The computing system 100 includes a processing subsystem 101 having one or more processor( $\overline{s}$ ) 102 and a system memory 104 communicating via an interconnection path that may include a memory hub 105. The memory hub 105 may be a separate component within a chipset component or may be integrated within the one or more processor(s) 102. The memory hub 105 couples with an I/O subsystem  $\frac{111}{111}$  via a communication link 106. The I/O subsystem 111 includes an I/O hub 107 that can enable the computing system 100 to receive input from one or more input device(s)  $108$ . Additionally, the I/O hub  $107$  can enable a display controller, which may be hub 107 can enable a display controller, which may be included in the one or more processor( $\frac{\sqrt{3}}{2}$ , to provide outputs to one or more display device( $\frac{\sqrt{3}}{2}$ ) 110A. In one embodiment, the one or more display device(s) 110A coupled with the I/O hub 107 can include a local, internal, .

or embedded display device.<br>  $[0047]$  In one embodiment, the processing subsystem 101 includes one or more parallel processor(s) 112 coupled to memory hub 105 via a bus or other communication link 113. The communication link 113 may be one of any number of standards based communication link technologies or protocols, such as, but not limited to PCI Express, or may be a vendor specific communications interface or communica tions fabric. In one embodiment, the one or more parallel  $processor(s)$  112 form a computationally focused parallel or vector processing system that an include a large number of processing cores and/or processing clusters, such as a many integrated core (MIC) processor. In one embodiment, the one or more parallel processor(s) 112 form a graphics processing subsystem that can output pixels to one of the one or more display device(s)  $110A$  coupled via the I/O Hub 107.

The one or more parallel processor( $s$ ) 112 can also include a display controller and display interface (not shown) to enable a direct connection to one or more display device(s) 110B .

[0048] Within the I/O subsystem 111, a system storage unit 114 can connect to the I/O hub 107 to provide a storage mechanism for the computing system 100. An I/O switch 116 can be used to provide an interface mechanism to enable connections between the I/O hub 107 and other components, such as a network adapter 118 and/or wireless network adapter 119 that may be integrated into the platform, and various other devices that can be added via one or more add-in device(s)  $120$ . The network adapter  $118$  can be an Ethernet adapter or another wired network adapter. The wireless network adapter 119 can include one or more of a Wi-Fi, Bluetooth, near field communication (NFC), or other network device that includes one or more wireless radios .

[0049] The computing system 100 can include other components not explicitly shown, including USB or other port connections, optical storage drives, video capture devices, and the like, may also be connected to the I/O hub 107.<br>Communication paths interconnecting the various components in FIG. 1 may be implemented using any suitable protocols, such as PCI (Peripheral Component Interconnect) based protocols (e.g., PCI-Express), or any other bus or point-to-point communication interfaces and/or protocol(s), such as the NV-Link high-speed interconnect, or interconnect protocols known in the art.

For the embounded in the embounded of processor (s) 112 incorporate circuitry optimized for graphics<br>and video processing, including, for example, video output<br>circuitry, and constitutes a graphics processing unit (GPU).<br>I parallel processor(s),  $112$  memory hub  $105$ , processor(s)  $102$ , and I/O hub  $107$  can be integrated into a system on chip (SoC) integrated circuit. Alternatively, the components of the computing system  $100$  can be int package to form a system in package (SIP) configuration. In one embodiment, at least a portion of the components of the computing system 100 can be integrated into a multi-chip<br>module (MCM), which can be interconnected with other<br>multi-chip modules into a modular computing system.<br>[0051] It will be appreciated that the computing system

100 shown herein is illustrative and that variations and modifications are possible. The connection topology, including the number and arrangement of bridges, the number of processor(s)  $102$ , and the number of parallel processor(s)  $112$ , may be modified as desired. For instance, in some embodiments, system memory 104 is connected to the  $processor(s)$  102 directly rather than through a bridge, while other devices communicate with system memory 104 via the memory hub  $105$  and the processor(s)  $102$ . In other alternative topologies, the parallel processor(s)  $112$  are connected to the I/O hub 107 or directly to one of the one or more processor(s)  $102$ , rather than to the memory hub 105. In other embodiments, the I/O hub 107 and memory hub 105 may be integrated into a single chip. Some embodiments may include two or more sets of processor( $s$ ) 102 attached via multiple sockets, which can couple with two or more instances of the parallel processor(s)  $112$ .

[0052] Some of the particular components shown herein are optional and may not be included in all implementations of the computing system 100. For example, any number of add in cards or peripherals may be supported, or some components may be eliminated. Furthermore, some architectures may use different terminology for components similar to those illustrated in FIG. 1. For example, the memory hub 105 may be referred to as a Northbridge in some architectures, while the I/O hub 107 may be referred to as a Southbridge.

[0053] FIG. 2A illustrates a parallel processor 200, according to an embodiment. The various components of the parallel processor 200 may be implemented using one or more integrated circuit devices, such as programmable processors, application specific integrated circuits (ASICs), or field programmable gate arrays (FPGA). The illustrated parallel processor 200 is a variant of the one or more parallel processor(s)  $112$  shown in FIG. 1, according to an embodiment.

[0054] In one embodiment, the parallel processor 200 includes a parallel processing unit includes an I/O unit 204 that enables communication with other devices, including other instances of the parallel processing unit 202. The I/O unit 204 may be directly connected to other devices. In one embodiment, the 1/0 unit 204 connects with other devices via the use of a hub or switch interface, such as memory hub 105. The connections between the memory hub  $105$  and the I/O unit  $204$  form a communication link 113. Within the parallel processing unit 202, the I/O unit 204 connects with a host interface 206 and a memory crossbar 216, where the host interface 206 receives commands directed to performing processing 216 receives commands directed to performing memory operations and the memory operations and the memory operations . [ 0055] When the host interface 206 receives a command

buffer via the I/O unit 204, the host interface 206 can direct work operations to perform those commands to a front end  $208$ . In one embodiment, the front end  $208$  couples with a scheduler 210, which is configured to distribute commands or other work items to a processing cluster array 212. In one embodiment, the scheduler 210 ensures that the processing cluster array 212 is properly configured and in a valid state before tasks are distributed to the processing clusters of the processing cluster array 212.

[ $0056$ ] The processing cluster array 212 can include up to "N" processing clusters (e.g., cluster 214A, cluster 214B, through cluster  $214N$ . Each cluster  $214A - 214N$  of the processing cluster array  $212$  can execute a large number of concurrent threads . The scheduler 210 can allocate work to the clusters  $214A - 214N$  of the processing cluster array  $212$  using various scheduling and/or work distribution algorithms, which may vary depending on the workload arising for each type of program or computation. The scheduling can be handled dynamically by the scheduler 210, or can be assisted in part by compiler logic during compilation of program logic configured for execution by the processing

program logic configured for execution by the program cluster array 212 can be allocated for process-<br>of processing cluster array 212 can be allocated for processing different types of programs or for performing different types of computations.

[0058] The processing cluster array 212 can be configured<br>to perform various types of parallel processing operations.<br>In one embodiment, the processing cluster array 212 is<br>configured to perform general-purpose parallel co filtering of video and/or audio data, performing modeling operations, including physics operations, and performing data transformations.

[0059] In one embodiment, the processing cluster array 212 is configured to perform parallel graphics processing operations. In embodiments in which the parallel processor 200 is configured to perform graphics processing o perform texture operations, as well as tessellation logic and other vertex processing logic. Additionally, the processing cluster array 212 can be configured to execute graphics processing related shader programs such as, but not limited to vertex shaders, tessellation shaders, geometry shaders, and pixel shaders. The parallel processing unit 202 can transfer data from system memory via the I/O unit 204 for processing. During processing the transferred data can be stored to on-chip memory (e.g., parallel processor memory 222) during processing, then written back to system memory.

[0060] In one embodiment, when the parallel processing unit 202 is used to perform graphics processing, the scheduler 210 can be configured to divide the processing workload into approximately equal sized tasks, to better enable distribution of the graphics processing operations to multiple clusters 214A-214N of the processing cluster array 212. In some embodiments, portions of the processing cluster array 212 can be configured to perform different types of processing. For example, a first portion may be configured to perform vertex shading and topology generation , a second etry shading, and a third portion may be configured to perform pixel shading or other screen space operations , to produce a rendered image for display. Intermediate data produced by one or more of the clusters 214A-214N may be stored in buffers to allow the intermediate data to be transmitted between clusters 214A-214N for further processing.  $[0061]$  During operation, the processing cluster array 212 can receive processing tasks to be executed via the scheduler 210, which receives commands defining processing tasks from front end 208. For graphics processing operations, processing tasks can include indices of data to be processed, e.g., surface (patch) data, primitive data, vertex data, and/or pixel data, as well as state parameters and commands defining how the data is to be processed (e.g., what program is to be executed). The scheduler  $210$  may be configured to fetch the indices corresponding to the tasks or may receive the indices from the front end 208. The front end 208 can be configured to ensure the processing cluster array 212 is configured to a valid state before the workload specified by incoming command buffers (e.g., batch-buffers, push buffers, etc.) is initiated.

[0062] Each of the one or more instances of the parallel processing unit 202 can couple with parallel processor memory 222. The parallel processor memory 222 can be accessed via the memory crossbar 216, which can receive memory requests from the processing cluster array 212 as well as the  $I/O$  unit 204. The memory crossbar 216 can access the parallel processor memory 222 via a memory interface 218. The memory interface 218 can include multiple partition units (e.g., partition unit 220A, partition unit 220B, through partition unit 220N) that can each couple to a portion  $(e.g., \text{ memory unit})$  of parallel processor memory 222. In one implementation, the number of partition units  $220A-220N$  is configured to be equal to the number of memory units, such that a first partition unit  $220A$  has a corresponding first memory unit 224A, a second partition unit 220B has a corresponding memory unit 224B, and an Nth partition unit 220N has a corresponding Nth memory unit 224N. In other embodiments, the number of partition units 220A-220N may not be equal to the number of memory devices.

[0063] In various embodiments, the memory units 224A-224N can include various types of memory devices, including dynamic random access memory (DRAM) or graphics random access memory, such as synchronous graphics random access memory (SGRAM), including graphics double data rate (GDDR) memory. In one embodiment, the memory units 224A-224N may also include 3D stacked memory, including but not limited to high bandwidth memory (HBM). Persons skilled in the art will appreciate that the specific implementation of the memory units 224A-224N can vary, and can be selected from one of various conventional designs. Render targets, such as frame buffers or texture maps may be stored across the memory units 224A-<br>224N, allowing partition units 220A-220N to write portions of each render target in parallel to efficiently use the available bandwidth of parallel processor memory 222. In some embodiments, a local instance of the parallel processor memory 222 may be excluded in favor of a unified memory design that utilizes system memory in conjunction with local cache memory.

[ $0064$ ] In one embodiment, any one of the clusters  $214A$ -214N of the processing cluster array 212 can process data that will be written to any of the memory units 224A-224N within parallel processor memory 222. The memory crossbar 216 can be configured to transfer the output of each cluster 214A-214N to any partition unit 220A-220N or to another cluster 214A-214N, which can perform additional processing operations on the output. Each cluster 214A-214N can communicate with the memory interface 218 through the memory crossbar 216 to read from or write to various external memory devices. In one embodiment, the memory crossbar 216 has a connection to the memory interface 218 to communicate with the I/O unit 204, as well as a connection to a local instance of the parallel processor memory 222, enabling the processing units within the different processing clusters 214A-214N to communicate with system memory or other memory that is not local to the parallel processing unit  $202$ . In one embodiment, the memory crossbar  $216$  can use virtual channels to separate traffic streams between the clusters  $214A - 214N$  and the partition units  $220A - 220N$ .

[0065] While a single instance of the parallel processing unit  $202$  is illustrated within the parallel processing unit  $202$  can be included. For example, multiple instances of the parallel processing unit 202 can be provided on a single add-in card,<br>or multiple add-in cards can be interconnected. The different<br>instances of the parallel processing unit 202 can be config-<br>ured to inter-operate even if the diff different numbers of processing cores, different amounts of<br>local parallel processor memory, and/or other configuration<br>differences. For example, and in one embodiment, some<br>instances of the parallel processing unit 202 ca instances . Systems incorporating one or more instances of the parallel processing unit 202 or the parallel processor 200 can be implemented in a variety of configurations and form factors, including but not limited to desktop, laptop, or handheld personal computers, servers, workstations, game

a mandneld personal computers, servers, workstations, game<br>consoles, and/or embedded systems.<br>[0066] FIG. 2B is a block diagram of a partition unit 220,<br>according to an embodiment. In one embodiment, the par-<br>tition unit 220 220A-220N of FIG. 2A. As illustrated, the partition unit 220 includes an L2 cache 221, a frame buffer interface 225, and a ROP 226 (raster operations unit). The L2 cache 221 is a read/write cache that is configured to perform load and store operations received from the memory crossbar 216 and ROP 226. Read misses and urgent write-back requests are output by L2 cache 221 to frame buffer interface  $\overline{225}$  for processing. Dirty updates can also be sent to the frame buffer via the frame buffer interface 225 for opportunistic processing. In one embodiment, the frame buffer interface 225 interfaces with one of the memory units in parallel processor memory,

such as the memory units  $224\text{\AA}-224\text{\AA}$  of FIG. 2A (e.g., within parallel processor memory 222).<br>[0067] In graphics applications, the ROP 226 is a processing unit that performs raster operations, such as stencil, z te processed graphics data that is stored in graphics memory. In some embodiments, the ROP 226 includes compression logic to compress z or color data that is written to memory and decompress z or color data that is read from memory. In some embodiments, the ROP 226 is included within each processing cluster (e.g., cluster 214A-214N of FIG. 2A) instead of within the partition unit 220. In such embodiment, read and write requests for pixel data are transmitted over the memory crossbar 216 instead of pixel fra

[0068] The processed graphics data may be displayed on a display device, such as one of the one or more display device  $(s)$  110 of FIG. 1, routed for further processing by the  $processor(s)$  102, or routed for further processing by one of the processing entities within the parallel processor 200 of FIG. 2A.

 $[0009]$  FIG. 2C is a block diagram of a processing cluster 214 within a parallel processing unit, according to an embodiment. In one embodiment, the processing cluster is an instance of one of the processing clusters 214A-214N of FIG. 2A. The processing cluster 214 can be configured to execute many threads in parallel, where the term "thread" refers to an instance of a particular program executing on a particular set of input data. In some embodiments, single-instruction, multiple-data (SIMD) instruction issue techniques are used to support parallel execution of a large number of threads without providing multiple independent<br>instruction units. In other embodiments, single-instruction,<br>multiple-thread (SIMT) techniques are used to support par-<br>allel execution of a large number of generall where all processing engines typically execute identical instructions , SIMT execution allows different threads to

subset of a SIMT processing regime. more readily follow divergent execution paths through a given thread program . Persons skilled in the art will under stand that a SIMD processing regime represents a functional

[0070] Operation of the processing cluster 214 can be controlled via a pipeline manager 232 that distributes processing tasks to SIMT parallel processors. The pipeline manager 232 receives instructions from the scheduler 210 of FIG. 2A and manages execution of those instructions via a graphics multiprocessor 234 and/or a texture unit 236. The illustrated graphics multiprocessor 234 is an exemplary instance of an SIMT parallel processor. However, various types of SIMT parallel processors of differing architectures may be included within the processing cluster 214. One or more instances of the graphics multiprocessor 234 can be included within a processing cluster 214. The graphics multiprocessor 234 can process data and a data crossbar 240 can be used to distribute the processed data to one of multiple possible destinations, including other shader units.<br>The pipeline manager 232 can facilitate the distribution of processed data by specifying destinations for processed data to be distributed vis the data crossbar 240.

[ $0071$ ] Each graphics multiprocessor 234 within the processing cluster 214 can include an identical set of functional execution logic (e.g., arithmetic logic units, load-store units, etc.). The functional execution logic can be configured in a pipelined manner in which new instructions can be issued execution logic may be provided. The functional logic supports a variety of operations including integer and floating point arithmetic comparison operations, Boolean operations bit-shifting, and computation of various algebraic functions. In one embodiment, the same functional - unit hardware can be leveraged to perform different operations hard any combination of functional units may be present.<br> **[0072]** The instructions transmitted to the processing clus-

ter 214 constitutes a thread. A set of threads executing across the set of parallel processing engines is a thread group. A thread group executes the same program on different input data. Each thread within a thread group can be assigned to a different processing engine within a graphics multiprocessor 234. A thread group may include fewer threads than the number of processing engines within the graphics multiprothe number of processing engines, one or more of the processing engines may be idle during cycles in which that thread group is being processed. A thread group may also include more threads than the number of processing engines within the graphics multiprocessor 234. When the thread group includes more threads than the number of processing engines within the graphics multiprocessor 234, processing can be performed over consecutive clock cycles. In one embodiment, multiple thread groups can be executed concessor 234. When a thread group includes fewer threads than

embodiment , multiprocessor **234** . [ 0073] In one embodiment , the graphics multiprocessor 234 includes an internal cache memory to perform load and store operations. In one embodiment, the graphics multiprocessor 234 can forego an internal cache and use a cache memory (e.g.,  $L1$  cache 308) within the processing cluster 214. Each graphics multiprocessor 234 also has access to L2 caches within the partition units ( $e.g.,$  partition units  $220A - 220N$  of FIG. 2A) that are shared among all processing clusters 214 and may be used to transfer data between threads . The graphics multiprocessor 234 may also access

off-chip global memory, which can include one or more of local parallel processor memory and/or system memory.<br>Any memory external to the parallel processing unit 202 may be used as global memory. Embodiments in which the processing cluster 214 includes multiple instances of the graphics multiprocessor  $234$  can share common instructions and data, which may be stored in the L1 cache  $308$ .

[0074] Each processing cluster 214 may include an MMU 245 (memory management unit) that is configured to map virtual addresses into physical addresses. In other embodiments, one or more instances of the MMU 245 may reside within the memory interface 218 of FIG. 2A. The MMU 245 includes a set of page table entries (PTEs) used to map a virtual address to a physical address of a tile (talk more about tiling) and optionally a cache line index. The MMU 245 may include address translation lookaside buffers (TLB) or caches that may reside within the graphics multiprocessor 234 or the L1 cache or processing cluster 214. The physical address is processed to distribute surface data access locality to allow efficient request interleaving among partition units. The cache line index may be used to request for a cache line is a hit or miss .

[0075] In graphics and computing applications, a processing cluster 214 may be configured such that each graphics multiprocessor 234 is coupled to a texture unit 236 for performing texture mapping operations, e.g., determi the texture data. Texture data is read from an internal texture L1 cache (not shown) or in some embodiments from the L1 cache within graphics multiprocessor 234 and is fetched from an L2 cache, local parallel processor memory, or system memory, as needed. Each graphics multiprocessor 234 outputs processed tasks to the data crossbar 240 to provide the processed task to another processing cluster 214 for further processing or to store the processed task in an L2 cache, local parallel processor memory, or system memory via the memory crossbar 216. A preROP 242 (pre-raster operations unit) is configured to receive data from graphics multiprocessor  $234$ , direct data to ROP units, which may be located with partition units as described her tion units 220A-220N of FIG. 2A). The preROP 242 unit can perform optimizations for color blending, organize pixel color data, and perform address translations.

[0076] It will be appreciated that the core architecture described herein is illustrative and that variations and modi-<br>fications are possible. Any number of processing units, e.g., graphics multiprocessor  $234$ , texture units  $236$ , preROPs  $242$ , etc., may be included within a processing cluster  $214$ .<br>Further, while only one processing cluster  $214$  is shown, a parallel processing unit as describe using separate and distinct processing units, L1 caches, etc.

[0077] FIG. 2D shows a graphics multiprocessor 234, according to one embodiment. In such embodiment, the graphics multiprocessor 234 couples with the pipeline manager 232 of the processing cluster 214. The graphics mul tiprocessor 234 has an execution pipeline including but not limited to an instruction cache 252, an instruction unit 254, an address mapping unit 256, a register file 258, one or more general purpose graphics processing unit (GPGPU) cores 262, and one or more load/store units 266. The GPGPU cores 262 and load/store units 266 are coupled with cache memory 272 and shared memory 270 via a memory and cache interconnect 268 .

[0078] In one embodiment, the instruction cache 252 receives a stream of instructions to execute from the pipeline manager 232. The instructions are cached in the instruction cache 252 and dispatched for execution by the instruction unit 254. The instruction unit 254 can dispatch instructions as thread groups (e.g., warps), with each thread of the thread group assigned to a different execution unit within GPGPU core 262. An instruction can access any of a local, shared, or global address space by specifying an address within a unified address space . The address mapping unit 256 can be used to translate addresses in the unified address space into a distinct memory address that can be accessed by the load/store units 266.

each functional unit is allocated a dedicated portion of the [0079] The register file 258 provides a set of registers for the functional units of the graphics multiprocessor 324. The register file 258 provides temporary storage for operands<br>connected to the data paths of the functional units (e.g.,<br>GPGPU cores 262, load/store units 266) of the graphics<br>multiprocessor 324. In one embodiment, the registe is divided between each of the functional units such that register file 258. In one embodiment, the register file 258 is divided between the different warps being executed by the graphics multiprocessor 324.

[0080] The GPGPU cores 262 can each include floating point units (FPUs) and/or integer arithmetic logic units (ALUs) that are used to execute instructions of the graphics multiprocessor 324. The GPGPU cores 262 can be similar in architecture or can differ in architecture, according to embodiments. For example, and in one embodiment, a first portion of the GPGPU cores 262 include a single precision FPU and an integer ALU while a second portion of the GPGPU cores include a double precision FPU. In one embodiment, the FPUs can implement the IEEE 754-2008 standard for floating point arithmetic or enable variable precision floating point arithmetic. The graphics multiprocessor 324 can additionally include one or more fixed function or special function units to perform specific funcfunction or such as copy rectangle or pixel blending operations. In one embodiment one or more of the GPGPU cores can also include fixed or special function logic.

[0081] The memory and cache interconnect 268 is an interconnect network that connects each of the functional units of the graphics multiprocessor 324 to the register file 258 and to the shared memory 270. In one embodiment, the memory and cache interconnect 268 is a crossbar interconnect that allows the load/store unit 266 to implement load and store operations between the shared memory 270 and the register file 258. The register file 258 can operate at the same frequency as the GPGPU cores 262, thus data transfer between the GPGPU cores 262 and the register file 258 is very low latency . The shared memory 270 can be used to enable communication between threads that execute on the functional units within the graphics multiprocessor 234. The cache memory 272 can be used as a data cache for example, to cache texture data communicated between the functional units and the texture unit 236. The shared memory 270 can also be used as a program managed cached. Threads executing on the GPGPU cores 262 can programmatically store

data within the shared memory in addition to the automati cally cached data that is stored within the cache memory 272.

[0082] FIGS. 3A-3B illustrate additional graphics multiprocessors, according to embodiments. The illustrated graphics multiprocessors 325, 350 are variants of the graphics multiprocessor 234 of FIG. 2C. The illustrated graphics<br>multiprocessors 325, 350 can be configured as a streaming<br>multiprocessor (SM) capable of simultaneous execution of a<br>large number of execution threads.

[0083] FIG. 3A shows a graphics multiprocessor 325 according to an additional embodiment. The graphics multiprocessor 325 includes multiple additional instances of execution resource units relative to the graphics multiprocessor 234 of FIG. 2D. For example, the graphics multiprocessor 325 can include multiple instances of the instruction<br>unit 332A-332B, register file 334A-334B, and texture unit(s) 344A-344B. The graphics multiprocessor 325 also includes<br>multiple sets of graphics or compute execution units (e.g.,<br>GPGPU core 336A-336B, GPGPU core 337A-337B,<br>GPGPU core 338A-338B) and multiple sets of load/store units 340A-340B. In one embodiment, the execution resource units have a common instruction cache 330, texture and/or data cache memory 342, and shared memory 346. The various components can communicate via an interconnect fabric 327. In one embodiment, the interconnect fabric 327 includes one or more crossbar switches to enable examples equality components of the graphics multiprocessor 325.<br> **[0084]** FIG. 3B shows a graphics multiprocessor 350

according to an additional embodiment. The graphics processor includes multiple sets of execution resources 356A 356D, where each set of execution resource includes multiple instruction units, register files, GPGPU cores, and load store units, as illustrated in FIG. 2D and FIG. 3A. The execution resources 356A-356D can work in concert with texture unit(s)  $360A-360D$  for texture operations, while sharing an instruction cache 354 , and shared memory 362. In one embodiment, the execution resources 356A-356D can share an instruction cache 354 and shared memory 362, as well as multiple instances of a texture and/or data cache memory 358A-358B. The various components can communicate via an interconnect fabric 352 similar to the intercon nect fabric 327 of FIG. 3A.

[0085] Persons skilled in the art will understand that the architecture described in FIGS. 1, 2A-2D, and 3A-3B are descriptive and not limiting as to the scope of the present embodiments. Thus, the techniques described herein may be implemented on any properly configured processing unit, including, without limitation, one or more mobile application processors, one or more desktop or server central processing units (CPUs) including multi-core CPUs, one or more parallel processing units, such as the parallel ing unit 202 of FIG. 2A, as well as one or more graphics processors or special purpose processing units, without departure from the scope of the embodiments described herein.

[0086] In some embodiments, a parallel processor or GPGPU as described herein is communicatively coupled to host/processor cores to accelerate graphics operations, machine-learning operations, pattern analysis operations, and various general purpose GPU (GPGPU) functions. The GPU may be communicatively coupled to the host proces sor/cores over a bus or other interconnect (e.g., a high-speed interconnect such as PCIe or NVLink). In other embodiments , the GPU may be integrated on the same package or chip as the cores and communicatively coupled to the cores over an internal processor bus/interconnect (i.e., internal to the package or chip). Regardless of the manner in which the GPU is connected, the processor cores may allocate work to the GPU in the form of sequences of commands/instructions contained in a work descriptor. The GPU then uses dedicated circuitry/logic for efficiently processing these commands/ instructions .

[0087] Techniques for GPU to Host Processor Interconnection

[0088] FIG. 4A illustrates an exemplary architecture in which a plurality of GPUs  $410-413$  are communicatively coupled to a plurality of multi-core processors  $405-406$  over high-speed links  $440-443$  (e.g., buses, point connects, etc.). In one embodiment, the high-speed links 440-443 support a communication throughput of  $4$  GB/s,  $30$  GB/s or higher, depending on the implementation. Various interconnect protocols may be used including, but not limited to, PCIe 4.0 or 5.0 and NVLink 2.0. However, the underlying principles of the invention are not limited to any particular communication protocol or throughput.

[0089] In addition, in one embodiment, two or more of the GPUs 410-413 are interconnected over high-speed links 444-445, which may be implemented using the same or different protocols/links than those used for high-speed links 440-443. Similarly, two or more of the multi-core processors 405-406 may be connected over high speed link 433 which may be symmetric multi-processor (SMP) buses operating at 20 GB/s, 30 GB/s, 120 GB/s or higher. Alternatively, all communication between the various system components shown in FIG. 4A may be accomplished using the same protocols/links (e.g., over a common interconnection fabric). As mentioned, however, the underlying principles of the invention are not limited to any particular type of intercon nect technology.

**405-406** is communicatively coupled to a processor memory<br>**405-406** is communicatively coupled to a processor memory<br>**401-402**, via memory interconnects **430-431**, respectively, and each GPU **410-413** is communicatively c memory 420-423 over GPU memory interconnects 450-453, respectively. The memory interconnects 430-431 and 450-453 may utilize the same or different memory access technologies. By way of example, and not limitation, the processor memories  $401-402$  and GPU memories  $420-423$  may be volatile memories such as dynamic random access<br>memories (DRAMs) (including stacked DRAMs), Graphics DDR SDRAM (GDDR) (e.g., GDDR5, GDDR6), or High Bandwidth Memory (HBM) and/or may be non-volatile memories such as 3D XPoint or Nano-Ram. In one embodiment, some portion of the memories may be volatile memory and another portion may be non-volatile memory 9

(e.g., using a two-level memory (2LM) hierarchy).<br>[0091] As described below, although the various processors 405-406 and GPUs 410-413 may be physically coupled to a particular memory  $401-402$ ,  $420-423$ , respectively, a unified memory architecture may be implemented in which the same virtual system address space (also referred to as the "effective address" space) is distributed among all of the various physical memories. For example, processor memories 401-402 may each comprise 64 GB of the system memory address space and GPU memories 420-423 may

each comprise 32 GB of the system memory address space (resulting in a total of 256 GB addressable memory in this example).

[0092] FIG. 4B illustrates additional details for an inter-<br>connection between a multi-core processor 407 and a graphics acceleration module 446 in accordance with one embodiment. The graphics acceleration module 446 may include one or more GPU chips integrated on a line card which is coupled to the processor 407 via the high-speed link 440. Alternatively, the graphics acceleration module 446 may be integrated on the same package or chip as the processor 407.<br>[0093] The illustrated processor 407 includes a plurality of

cores 460A-460D, each with a translation lookaside buffer 461A-461D and one or more caches 462A-462D. The cores may include various other components for executing instructions and processing data which are not illustrated to avoid obscuring the underlying principles of the invention (e.g., instruction fetch units, branch prediction units, decoders, execution units, reorder buffers, etc.). The caches 462A- $462D$  may comprise level 1 (L1) and level 2 (L2) caches. In addition, one or more shared caches 426 may be included in the caching hierarchy and shared by sets of the cores 460A-460D. For example, one embodiment of the processor 407 includes 24 cores, each with its own L1 cache, twelve shared L2 caches, and twelve shared L3 caches. In this embodiment, one of the L2 and L3 caches are shared by two adjacent cores. The processor 407 and the graphics accelerator integration module 446 connect with system memory

441, which may include processor memories 401-402.<br>[ 0094] Coherency is maintained for data and instructions stored in the various caches 462A-462D, 456 and system memory 441 via inter-core communication over a coherence bus 464. For example, each cache may have cache coherency logic/circuitry associated therewith to communicate to over the coherence bus 464 in response to detected reads or writes to particular cache lines. In one implementation, a cache snooping protocol is implemented over the coherence<br>bus 464 to snoop cache accesses. Cache snooping/coherency techniques are well understood by those of skill in the art and will not be described in detail here to avoid obscuring the underlying principles of the invention.

[0095] In one embodiment, a proxy circuit 425 communicatively couples the graphics acceleration module 446 to the coherence bus 464, allowing the graphics acceleration module 446 to participate in the cache coherence protocol as a peer of the cores. In particular, an interface 435 provides connectivity to the proxy circuit 425 over high-speed link 440 (e.g., a PCIe bus, NVLink, etc.) and an interface 437 connects the graphics acceleration module 446 to the link 440.

[0096] In one implementation, an accelerator integration circuit 436 provides cache management, memory access, context management, and interrupt management services on behalf of a plurality of graphics processing engines 4 432, N of the graphics acceleration module 446. The graphics processing engines 431, 432, N may each comprise a separate graphics processing unit (GPU). Alternatively, the graphics processing engines 431, 432, N may comprise different types of graphics processing engines within a GPU such as graphics execution units, media processing engines such as graphics execution units , media processing engines ( e.g. , video encoders / decoders ) , samplers , and blit engines . In other words , the graphics acceleration module may be a GPU with a plurality of graphics processing engines 431 432, N or the graphics processing engines 431-432, N may be individual GPUs integrated on a common package, line card, or chip.

[0097] In one embodiment, the accelerator integration circuit 436 includes a memory management unit (MMU) 439 for performing various memory management functions such as virtual-to-physical memory translations (also referred to as effective-to-real memory translations) and memory access protocols for accessing system memory 441. The MMU 439 may also include a translation lookaside buffer (TLB) (not shown) for caching the virtual/effective to physical/real address translations. In one implementation, a cache 438 stores commands and data for efficient access by the graphics processing engines 431-432, N. In one embodiment, the data stored in cache 438 and graphics memories 433-434, N is kept coherent with the core caches 462A-462D, 456 and system memory 411. As mentioned, this may be accomplished via proxy circuit 425 which takes part in the cache coherency mechanism on behalf of cache 438 and memories 433-434, N (e.g., sending updates to the cache 438 related to modifications/accesses of cache lines on processor caches 462A-462D, 456 and receiving updates from the cache 438).<br>[ $0098$ ] A set of registers 445 store context data for threads

executed by the graphics processing engines 431-432 , N and a context management circuit 448 manages the thread contexts. For example, the context management circuit 448 may perform save and restore operations to save and restore contexts of the various threads during contexts switches (e.g., where a first thread is saved and a second thread is stored so that the second thread can be execute by a graphics processing engine). For example, on a context switch, the context management circuit 448 may store current register values to a designated region in memory (e.g., identified by a context pointer). It may then restore the register values when returning to the context. In one embodiment, an interrupt management circuit 447 receives and processes<br>interrupts received from system devices.<br>[0099] In one implementation, virtual/effective addresses

from a graphics processing engine 431 are translated to real/physical addresses in system memory 411 by the MMU 439. One embodiment of the accelerator integration circuit 436 supports multiple (e.g., 4, 8, 16) graphics accelerator modules  $446$  and/or other accelerator devices. The graphics accelerator module 446 may be dedicated to a single application executed on the processor 407 or may be shared between multiple applications. In one embodiment, a virtualized graphics execution environment is presented in which the resources of the graphics processing engines 431-432, N are shared with multiple applications or virtual machines (VMs). The resources may be subdivided into "slices" which are allocated to different VMs and/or applications based on the processing requirements and priorities associated with the VMs and/or applications.

 $[0100]$  Thus, the accelerator integration circuit acts as a bridge to the system for the graphics acceleration module 446 and provides address translation and system memory cache services. In addition, the accelerator integration circuit 436 may provide virtualization facilities for the host processor to manage virtualization of the graphics processing

engines, interrupts, and memory management.<br>
[0101] Because hardware resources of the graphics processing engines 431-432, N are mapped explicitly to the real address space seen by the host processor 407, any host

processor can address these resources directly using an effective address value. One function of the accelerator integration circuit  $436$ , in one embodiment, is the physical separation of the graphics processing engines 431-432, N so that they appear to the system as independent units.

may be volatile memories such as DRAMs (including stacked DRAMs), GDDR memory (e.g., GDDR5, GDDR6), or HBM, and/or may be non-volatile memories such as 3D [0102] As mentioned, in the illustrated embodiment, one or more graphics memories  $433-434$ , M are coupled to each of the graphics processing engines 431-432, N, respectively.<br>The graphics memories 433-434, M store instructions and<br>data being processed by each of the graphics processing<br>engines 431-432, N. The graphics memories 433-434 XPoint or Nano-Ram.

[0103] In one embodiment, to reduce data traffic over link 440, biasing techniques are used to ensure that the data stored in graphics memories 433-434. M is data which will be used most frequently by the graphics processing engines 431-432, N and preferably not used by the cores 460A-460D (at least not frequently). Similarly, the biasing mechanism<br>attempts to keep data needed by the cores (and preferably not the graphics processing engines  $431-432$ , N) within the caches 462A-462D, 456 of the cores and system memory 411 .

[0104] FIG. 4C illustrates another embodiment in which the accelerator integration circuit 436 is integrated within<br>the processor 407. In this embodiment, the graphics prothe processing engines 431-432, N communicate directly over the high-speed link 440 to the accelerator integration circuit 436 via interface 437 and interface 435 (which, again, may be utilize any form of bus or interface protocol). The accelerator integration circuit 436 may perform the same operations as those described with respect to FIG. 4B, but potentially at a higher throughput given its close proximity to the coherency bus 462 and caches 462A-462D, 426.

[0105] One embodiment supports different programming models including a dedicated-process programming model (no graphics acceleration module virtualization) and shared programming models (with virtualization). The latter may include programming models which are controlled by the accelerator integration circuit 436 and programming models which are controlled by the graphics acceleration module 446 .

[ $0106$ ] In one embodiment of the dedicated process model, graphics processing engines  $431-432$ , N are dedicated to a single application or process under a single operating system. The single application can funnel other

requests to the graphics engines  $431-432$ , N, providing<br>virtualization within a VM/partition.<br>[0107] In the dedicated-process programming models, the<br>graphics processing engines  $431-432$ , N, may be shared by<br>multiple VM cessing engines 431-432, N to allow access by each operating system. For single-partition systems without a hypervisor, the graphics processing engines 431-432, N are owned by the operating system. In both cases, the operating system can virtualize the graphics processing engines 431-432, N to

provide access to each process or application.<br>[0108] For the shared programming model, the graphics<br>acceleration module 446 or an individual graphics process-<br>ing engine 431-432, N selects a process element using a process handle. In one embodiment, process elements are stored in system memory 411 and are addressable using the effective address to real address translation techniques described herein. The process handle may be an implementation-specific value provided to the host process when registering its context with the graphics processing engine 431-432, N (that is, calling system software to add process element to the process element linked list). The lower 16-bits of the process handle may be the offset of the process element within the process element linked list.

[0109] FIG. 4D illustrates an exemplary accelerator integration slice 490. As used herein, a "slice" comprises a specified portion of the processing resources of the accelerator integration circuit 436. Application effective address space 482 within system memory 411 stores process elements 483. In one embodiment, the process elements 483 are stored in response to GPU invocations  $481$  from applications  $480$  executed on the processor  $407$ . A process element  $483$  contains the process state for the corresponding application  $480$ . A work descriptor (WD)  $484$ the WD 484 is a pointer to the job request queue in the application's address space 482 .

[0110] The graphics acceleration module 446 and/or the individual graphics processing engines 431-432, N can be shared by all or a subset of the processes in the system. Embodiments of the invention include an infrastructure for setting up the process state and sending a WD 484 to a graphics acceleration module 446 to start a job in a virtu alized environment.

[0111] In one implementation, the dedicated-process programming model is implementation-specific. In this model, a single process owns the graphics acceleration module 446 or an individual graphics processing engine 431. Because the graphics acceleration module  $446$  is owned by a single process, the hypervisor initializes the accelerator integration process initializes the accelerator integration and the operating system initializes the accelerator integration circuit  $436$  for the owning process at the time when the graphics acceleration module 446 is assigned.

[0112] In operation, a WD fetch unit 491 in the accelerator integration slice 490 fetches the next WD 484 which includes an indication of the work to be done by one of the graphics processing engines of the graphics acceleration module 446. Data from the WD 484 may be stored in registers 445 and used by the MMU 439, interrupt management circuit 447 and/or context management circuit 446 as illustrated. For example, one embodiment of the MMU 439 includes segment/page walk circuitry for accessing segment/<br>page tables 486 within the OS virtual address space 485. The interrupt management circuit 447 may process interrupt events 492 received from the graphics acceleration module 446. When performing graphics operations, an effective address 493 generated by a graphics processing engine 431-432 , N is translated to a real address by the MMU 439 .

[0113] In one embodiment, the same set of registers 445 are duplicated for each graphics processing engine 431-432, N and/or graphics acceleration module 446 and may be initialized by the hypervisor or operating system. Each of these duplicated registers may be included in an accelerator integration slice 490. Exemplary registers that may be initialized by the hypervisor are shown in Table 1.

TABLE 1

#### Hypervisor Initialized Registers

- Slice Control Register
- Real Address (RA) Scheduled Processes Area Pointer<br>Authority Mask Override Register
- 
- 4 Interrupt Vector Table Entry Offset
- Interrupt Vector Table Entry Limit<br>State Register
- 
- Logical Partition ID
- Real address (RA) Hypervisor Accelerator Utilization Record Pointer Storage Description Register
- 

[0114] Exemplary registers that may be initialized by the operating system are shown in Table 2.

TABLE 2

Operating System Initialized Registers	
Process and Thread Identification	

- 1 Process and Thread Identification Effective Address ( EA ) Context Save / Restore Pointer
- 2 Virtual Address (VA) Accelerator Utilization Record Pointer
- 3 Virtual Address (VA) Storage Segment Table Pointer
- 4 5 Authority Mask
- 6 Work descriptor

[0115] In one embodiment, each WD 484 is specific to a particular graphics acceleration module 446 and/or graphics processing engines 431-432, N. It contains all the information a graphics processing engine 431-432, N requires to do its work or it can be a pointer to a memory location where the application has set up a command queue of work to be

a embodiment of a shared model . This embodiment includes the completed.<br> **(0116]** FIG. 4E illustrates additional details for one a hypervisor real address space 498 in which a process<br>element list 499 is stored. The hypervisor real address space<br>498 is accessible via a hypervisor 496 which virtualizes the<br>graphics acceleration module engines for the system 495.

[ 0117 ] The shared programming models allow for all or a subset of processes from all or a subset of partitions in the system to use a graphics acceleration module 446. There are two programming models where the graphics acceleration module 446 is shared by multiple processes and partitions:<br>time-sliced shared and graphics directed shared.

[0118] In this model, the system hypervisor  $496$  owns the graphics acceleration module  $446$  and makes its function available to all operating systems 495. For a graphics acceleration module 446 to support virtualization by the system hypervisor 496, the graphics acceleration module 446 may adhere to the following requirements: 1) An application's job request must be autonomous (that is, the state does not need to be maintained between jobs), or the graphics acceleration module 446 must provide a context save and restore mechanism. 2) An application's job request is guaranteed by the graphics acceleration module 446 to complete in a specified amount of time, including any translation faults, or the graphics acceleration module 446 provides the ability to preempt the processing of the job. 3) The graphics acceleration module 446 must be guaranteed fairness between processes when operating in

[ $0119$ ] In one embodiment, for the shared model, the application  $480$  is required to make an operating system  $495$ 

system call with a graphics acceleration module 446 type, a work descriptor ( $\overline{WD}$ ), an authority mask register ( $\overline{AMP}$ ) value, and a context save/restore area pointer (CSRP). The graphics acceleration module 446 type describes the targeted<br>acceleration function for the system call. The graphics<br>acceleration module 446 type may be a system-specific<br>value. The WD is formatted specifically for the gra pointer to a queue of commands, or any other data structure to describe the work to be done by the graphics acceleration module 446. In one embodiment, the AMR value is the AMR<br>state to use for the current process. The value passed to the operating system is similar to an application setting the AMR. If the accelerator integration circuit 436 and graphics acceleration module 446 implementations do not support a User Authority Mask Override Register (UAMOR), the operating system may apply the current UAMOR value to the AMR value before passing the AMR in the hypervisor call. The hypervisor 496 may optionally apply the current Authority Mask Override Register (AMOR) value before placing the AMR into the process element 483. In one embodiment, the CSRP is one of the registers 445 containing the effective address of an area in the application's address space 482 for the graphics acceleration module 446 to save and restore the context state . This pointer is optional if no state is required to be saved between jobs or when a job is preempted. The context save/restore area may be pinned

system memory.<br>[0120] Upon receiving the system call, the operating system 495 may verify that the application 480 has registered<br>and been given the authority to use the graphics acceleration<br>module 446. The operating syst visor 496 with the information shown in Table 3.

TABLE 3

#### OS to Hypervisor Call Parameters

- 1 A work descriptor (WD)<br>2 An Authority Mask Register (AMR) value (potentially masked).
- 2 An effective address (EA) Context Save/Restore Area Pointer (CSRP) 4 A process ID (PID) and optional thread ID (TID) 5 A virtual address (VA) accelerator utilization record pointer (AURP)
- 
- 
- 5 The virtual address of the storage segment table pointer (SSTP) 7 A logical interrupt service number (LISN)
- 

[0121] Upon receiving the hypervisor call, the hypervisor 496 verifies that the operating system 495 has registered and been given the authority to use the graphics acceleration module 446. The hypervisor 496 then puts the process element 483 into the process element linked list for the corresponding graphics acceleration module 446 type. The process element may include the information shown in Table 4

TABLE 4

#### Process Element Information

- 
- 1 A work descriptor (WD)<br>
2 An Authority Mask Register (AMR) value (potentially masked).<br>
3 An effective address (EA) Context Save/Restore Area Pointer (CSRP)<br>
4 A process ID (PID) and optional thread ID (TID)<br>
5 A virtual
- 

6 The virtual address of the storage segment table pointer (SSTP)

TABLE 4-continued Process Element Information

# 7 A logical interrupt service number (LISN)<br>8 Interrupt vector table, derived from the hypervisor call parameters.<br>9 A state register (SR) value

10 A logical partition ID (LPID)

11 A real address ( RA ) hypervisor accelerator utilization record pointer 12 The Storage Descriptor Register ( SDR )

[0122] In one embodiment, the hypervisor initializes a plurality of accelerator integration slice 490 registers 445.

[0123] As illustrated in FIG. 4F, one embodiment of the invention employs a unified memory addressable via a common virtual memory address space used to access the physical processor memories 401-402 and GPU memories 420-423. In this implementation, operations executed on the GPUs 410-413 utilize the same virtual/effective memory address space to access the processors memories 401-402 and vice versa, thereby simplifying programmability. In one embodiment, a first portion of the virtual/effective address space is allocated to the processor memory 401, a second portion to the second processor memory 402, a third portion to the GPU memory 420, and so on. The entire virtual/ effective memory space (sometimes referred to as the effective address space) is thereby distributed across each of the processor memories  $401-402$  and GPU memories  $420-423$ , allowing any processor or GPU to access any physical memory with a virtual address mapped to that memory.

[0124] In one embodiment, bias/coherence management circuitry 494A-494E within one or more of the MMUs 439A-439E ensures cache coherence between the caches of the host processors (e.g.,  $405$ ) and the GPUs 410-413 and implements biasing techniques indicating the physical memories in which certain types of data should be stored. While multiple instances of bias/coherence management<br>circuitry 494A-494E are illustrated in FIG. 4F, the bias/<br>coherence circuitry may be implemented within the MMU of one or more host processors 405 and/or within the accelerator integration circuit 436 .

[0125] One embodiment allows GPU-attached memory 420-423 to be mapped as part of system memory, and accessed using shared virtual memory (SVM) technology. but without suffering the typical performance drawbacks associated with full system cache coherence . The ability to GPU-attached memory 420-423 to be accessed as system<br>memory without onerous cache coherence overhead provides a beneficial operating environment for GPU offload. This arrangement allows the host processor 405 software to setup operands and access computation results, without the overhead of tradition I/O DMA data copies. Such traditional copies involve driver calls, interrupts and memory mapped I/O (MMIO) accesses that are all inefficient relative to simple memory accesses. At the same time, the ability to access GPU attached memory 420-423 without cache coher ence overheads can be critical to the execution time of an offloaded computation. In cases with substantial streaming write memory traffic, for example, cache coherence overhead can significantly reduce the effective write bandwidth seen by a GPU 410-413. The efficiency of operand setup, the efficiency of results access, and the efficiency of GPU computation all play a role in determining the effectiveness of GPU offload.

[0126] In one implementation, the selection of between GPU bias and host processor bias is driven by a bias tracker data structure. A bias table may be used, for example, which may be a page-granular structure (i.e., controlled at the granularity of a memory page) that includes 1 or 2 bits per<br>GPU-attached memory page. The bias table may be imple-<br>mented in a stolen memory range of one or more GPUattached memories 420-423 , with or without a bias cache in the GPU  $410-413$  (e.g., to cache frequently/recently used entries of the bias table). Alternatively, the entire bias table may be maintained within the GPU.

processor bias complete the request like a normal memory  $[0127]$  In one implementation, the bias table entry associated with each access to the GPU-attached memory 420-423 is accessed prior the actual access to the GPU memory, causing the following operations. First, local requests from<br>the GPU 410-413 that find their page in GPU bias are forwarded directly to a corresponding GPU memory 420-423. Local requests from the GPU that find their page in host bias are forwarded to the processor  $405$  (e.g., over a high-<br>speed link as discussed above). In one embodiment, requests<br>from the processor  $405$  that find the requested page in host read. Alternatively, requests directed to a GPU-biased page may be forwarded to the GPU 410-413. The GPU may then transition the page to a host processor bias if it is not

the currently using the page.<br> **[0128]** The bias state of a page can be changed either by a software-based mechanism, a hardware-assisted softwarebased mechanism, or, for a limited set of cases, a purely hardware-based mechanism.

[0129] One mechanism for changing the bias state employs an API call (e.g. OpenCL), which, in turn, calls the GPU's device driver which, in turn, sends a message (or enqueues a command descriptor) to the GPU directing it to change the bias state and, for some transitions, perform a cache flushing operation in the host. The cache flushing operation is required for a transition from host processor 405 bias to GPU bias, but is not required for the opposite transition.

[0130] In one embodiment, cache coherency is maintained<br>by temporarily rendering GPU-biased pages uncacheable by<br>the host processor 405. To access these pages, the processor<br>405 may request access from the GPU 410 which ma mentation. Thus, to reduce communication between the processor 405 and GPU 410 it is beneficial to ensure that GPU-biased pages are those which are required by the GPU but not the host processor 405 and vice versa.

[0131] Graphics Processing Pipeline

[0132] FIG. 5 illustrates a graphics processing pipeline 500, according to an embodiment. In one embodiment, a graphics processor can implement the illustrated graphics processing pipeline 500. The graphics processor can b 2A, which, in one embodiment, is a variant of the parallel processor(s)  $112$  of FIG. 1. The various parallel processing systems can implement the graphics processing pipeline 500 via one or more instances of the parallel processing unit  $(e.g.,$  parallel processing unit 202 of FIG. 2A) as described herein. For example, a shader unit (e.g., graphics multiprocessor 234 of FIG. 2D) may be configured to perform the functions of one or more of a vertex processing unit 504, a tessellation control processing unit 508, a tessellation evaluation processing unit 512, a geometry processing unit 516, and a fragment/pixel processing unit 524. The functions of data assembler 502, primitive assemblers 506, 514, 518, tessellation unit 510, rasterizer 522, and raster operations<br>unit 526 may also be performed by other processing engines<br>within a processing cluster (e.g., processing cluster 214 of<br>FIG. 3A) and a corresponding partition u unit  $220A - 220N$  of FIG. 2C). The graphics processing pipeline 500 may also be implemented using dedicated processing units for one or more functions. In one embodi-<br>ment, one or more portions of the graphics processing pipeline 500 can be performed by parallel processing logic within a general-purpose processor (e.g., CPU). In one embodiment, one or more portions of the graphics processing pipeline 500 can access on-chip memory (e.g., parallel processor memory 222 as in FIG. 2A) via a memory interface 528, which may be an instance of the memory interface 218 of FIG. 2A.

[0133] In one embodiment, the data assembler 502 is a processing unit that collects vertex data for surfaces and primitives. The data assembler 502 then outputs the vertex data, including the vertex attributes, to the vertex processing unit 504. The vertex processing unit 504 is a programmable execution unit that executes vertex shader programs, lighting and transforming vertex data as specified by the vertex shader programs. The vertex processing unit 504 reads data that is stored in cache, local or system memory for use in processing the vertex data and may be programmed to representation to a world space coordinate space or a nor-malized device coordinates space.

[ $0134$ ] A first instance of a primitive assembler  $506$ receives vertex attributes from the vertex processing unit 504. The primitive assembler 506 readings stored vertex attributes as needed and constructs graphics primitives for

attributes as heeded and constructs graphics primitives for<br>processing by tessellation control processing unit 508. The<br>graphics primitives include triangles, line segments, points,<br>patches, and so forth, as supported by v tessellation evaluation processing unit 512. The tessellation control processing unit 508 can also compute tessellation factors for edges of geometric patches. A tessellation factor applies to a single edge and quantifies a view-dependent level of detail associated with the edge. A tessellation unit 510 is configured to receive the tessellation factors for edges<br>of a patch and to tessellate the patch into multiple geometric<br>primitives such as line, triangle, or quadrilateral primitives,<br>which are transmitted to a tes operates on parameterized coordinates of the subdivided patch to generate a surface representation and vertex attri butes for each vertex associated with the geometric primitives.

[0136] A second instance of a primitive assembler 514 receives vertex attributes from the tessellation evaluation processing unit 512, reading stored vertex attributes as needed, and constructs graphics primitives for processing by the geometry processing unit 516. The geometry processing unit 516 is a programmable execution unit that executes geometry shader programs to transform graphics primitives received from primitive assembler 514 as specified by the geometry shader programs. In one embodiment, the geometry processing unit 516 is programmed to subdivide the graphics primitives into one or more new graphics primi-

tives and calculate parameters used to rasterize the new<br>graphics primitives.<br>[0137] In some embodiments, the geometry processing<br>unit 516 can add or delete elements in the geometry stream.<br>The geometry processing unit 516 parameters and vertices from the geometry processing unit 516 and constructs graphics primitives for processing by a viewport scale, cull, and clip unit 520. The geometry processing unit 516 reads data that is stored in parallel processor memory or system memory for use in processing the geometry data. The viewport scale, cull, and clip unit 520 performs clipping, culling, and viewport scaling and outputs<br>processed graphics primitives to a rasterizer 522.<br>[0138] The rasterizer 522 can perform depth culling and

other depth-based optimizations. The rasterizer 522 also performs scan conversion on the new graphics primitives to generate fragments and outputs those fragments and asso ciated coverage data to the fragment/pixel processing unit 524 .

[0139] The fragment/pixel processing unit 524 is a programmable execution unit that is configured to execute fragment shader programs or pixel shader programs. The fragment/pixel processing unit 524 transforming fragments or pixels received from rasterizer 522, as specified by the fragment or pixel shader programs. For example, the fragment/pixel processing unit  $\overline{524}$  may be programmed to perform operations included but not limited to texture mapping, shading, blending, texture correction and perspective correction to produce shaded fragments or pixels that are output to a raster operations unit 526. The fragment/pixel processing unit 524 can read data that is stored in either the parallel processor memory or the system memory for use shader programs may be configured to shade at sample, pixel, tile, or other granularities, depending on the sampling rate configured for the processing units.

[0140] The raster operations unit  $526$  is a processing unit that performs raster operations including, but not limited to stencil, z test, blending, and the like, and outputs pixel data as processed graphics data to be storage in graphics memory, e.g., parallel processor memory 222 as in FIG. 2A, and/or system memory 104 as in FIG. 1, to be displayed on the one or more display device( $s$ ) 110 or for further processing by one of the one or more processor( $s$ ) 102 or parallel processor( $s$ ) 112. In some embodiments, the raster operations unit 526 is configured to compress z or color data that is written to memory and decompress z or color data that is read from memory.

[0141] FIG. 6 illustrates a computing device 600 hosting a data sharing and compression expansion mechanism ("sharing and expansion mechanism")  $610$  according to one embodiment. Computing device 600 represents a communication and data processing device including (but not limited to) smart wearable devices, smartphones, virtual reality (VR) devices, head-mounted display (HMDs), mobile computers, Internet of Things (IoT) devices, laptop computers, desktop computers, server computers, etc., and be similar to or the same as computing device 100 of FIG.<br>1; accordingly, for brevity, clarity, and ease of understanding, many of the details stated above with reference to FIGS.

1-5 are not further discussed or repeated hereafter.<br>
[0142] Computing device 600 may further include (with-<br>
out limitations) an autonomous machine or an artificially intelligent agent, such as a mechanical agent or machine, an electronics agent or machine, a virtual agent or machine, an electro-mechanical agent or machine, etc. Examples of autonomous machines or artificially intelligent agents may include (without limitation) robots, autonomous vehicles<br>(e.g., self-driving cars, self-flying planes, self-sailing boats,<br>etc.), autonomous equipment (self-operating construction<br>vehicles, self-operating medical equipment like. Throughout this document, "computing device" may be interchangeably referred to as "autonomous machine" or " artificially intelligent agent" or simply " robot".<br>
[0143] It contemplated that although " autonomous

vehicle" and "autonomous driving" are referenced throughout this document, embodiments are not limited as such. For example, "autonomous vehicle" is not limed to an automobile but that it may include any number and type of autonomous machines, such as robots, autonomous equipment, household autonomous devices, and/or the like, and any one<br>or more tasks or operations relating to such autonomous machines may be interchangeably referenced with autonomous driving.

[0144] Computing device 600 may further include (without limitations) large computing systems, such as server computers, desktop computers, etc., and may further include set-top boxes (e.g., Internet-based cable television set-top boxes, etc.), global positioning system (GPS)-based devices, etc. Computing device 600 may include mobile computing devices serving as communication devices, such as cellular phones including smartphones, personal digital smart televisions, television platforms, wearable devices (e.g., glasses, watches, bracelets, smartcards, jewelry, clothing items, etc.), media players, etc. For example, in one embodiment, computing device 600 may include a mobile computing device employing a computer platform hosting an integrated circuit ("IC"), such as system on a chip ("SoC" or "SOC"), integrating various hardware and/or software components of computing device 600 on a single chip.

[0145] As illustrated, in one embodiment, computing device 600 may include any number and type of hardware and/or software components, such as (without limitation) graphics processing unit ("GPU" or simply "graphics processor") 614, graphics driver (also referred to as "GPU driver", "graphics driver logic", "driver logic", user-mode driver (UMD), UMD, user-mode driver framework<br>(UMDF), UMDF, or simply "driver") 616, central processing unit ("CPU" or simply "application processor") 612, memory 608, network devices, drivers, or the like, as well as input/output  $(I/O)$  sources 604, such as touchscreens, touch panels, touch pads, virtual or regular keyboards, virtual or regular mice, ports, connectors, etc. Computing device 600 may include operating system (OS) 606 serving as an interface between hardware and/or physical resources of the computer device  $600$  and a user. It is contemplated that graphics processor  $614$  and application processor  $612$  may<br>be one or more of processor(s)  $102$  of FIG. 1.<br>[0146] It is to be appreciated that a lesser or more<br>equipped system than the example described above may be

preferred for certain implementations. Therefore, the configuration of computing device 600 may vary from implementation to implementation depending upon numerous factors, such as price constraints, performance requirement factors technological improvements, or other circumstances.<br> **[0147]** Embodiments may be implemented as any or a

combination of: one or more microchips or integrated circuits interconnected using a parentboard, hardwired logic, software stored by a memory device and executed by a microprocessor, firmware, an application specific integrated circuit (ASIC), and/or a field programmable gate array (FPGA). The terms "logic", "module", "component", "en example, software or hardware and/or combinations of software and hardware.

[0148] In one embodiment, sharing and expansion mechanism 610 may be hosted or facilitated by operating system 606 of computing device 600. In another embodiment, sharing and expansion mechanism 610 may be hosted by or part of graphics processing unit ("GPU" or simply "graphics processor") 614 or firmware of graphics processor 614. For example, sharing and expansion mechanism 610 may be embedded in or implemented as part of the processing hardware of graphics processor 614. Similarly, in yet another embodiment, sharing and expansion mechanism 610 may be hosted by or part of central processing unit ("CPU" or simply "application processor")  $\dot{\textbf{6}}$ 12. For example, sharing and expansion mechanism  $\dot{\textbf{6}}$ 10 may be embedded in or implemented as part of the processing hardware of application processor 612. In yet another embodiment, sharing and expansion mechanism 610 may be hosted by or part of any number and type of components of computing device portion may be hosted by or part of graphics processor 614, another portion may be hosted by or part of application processor 612, while one or more portions of sharing and expansion mechanism 610 may be hosted by or part of operating system 606 and/or any number and type of devices<br>of computing device 600. It is contemplated that one or more portions or components of sharing and expansion mechanism 610 may be employed as hardware, software, and/or firmware.

[ 0149 ] It is contemplated that embodiments are not limited to any particular implementation or hosting of sharing and expansion mechanism 610 and that sharing and expansion mechanism 610 and one or more of its components may be implemented as hardware, software, firmware, or any combination thereof.

[0150] Computing device 600 may host network interface (s) to provide access to a network, such as a LAN, a wide area network (WAN), a metropolitan area network (MAN), a personal area network (PAN), Bluetooth, a cloud network, a mobile network (e.g.,  $3^{rd}$  Generation (3G),  $4^{th}$  Generation  $(4G)$ , etc.), an intranet, the Internet, etc. Network interface (s) may include, for example, a wireless network interface having antenna, which may represent one or more antenna (e). Network interface $(s)$  may also include, for example, a wired network interface to communicate with remote devices via network cable, which may be, for example, an Ethernet cable, a coaxial cable, a fiber optic cable, a serial cable, or a parallel cable.<br>
[0151] Embodiments may be provided, for example, as a

computer program product which may include one or more

machine-readable media having stored thereon machine-executable instructions that, when executed by one or more machines such as a computer, network of computers, or other electronic devices, may result in the one or more machines carrying out operations in accordance with embodiments described herein. A machine-readable medium<br>may include, but is not limited to, floppy diskettes, optical disks, CD-ROMs (Compact Disc-Read Only Memories), and magneto-optical disks, ROMs, RAMs, EPROMs (Erasable Programmable Read Only Memories), EEPROMs (Electrically Erasable Programmable Read Only Memories), magnetic or optical cards, flash memory, or other type of media/machine-readable medium suitable for storing machine-executable instructions.<br>[0152] Moreover, embodiments may be downloaded as a

computer program product, wherein the program may be transferred from a remote computer  $(e.g., a server)$  to a requesting computer (e.g., a client) by way of one or more data signals embodied in and/or modulated by a carrier wave or other propagation medium via a communication link (e.g., a modem and/or network connection).

noted that throughout this document, terms like graphics<br>domain" may be referenced interchangeably with "graphics<br>processing unit", "graphics processor", or simply "GPU" and similarly, "CPU domain" or "host domain" may be<br>referenced interchangeably with "computer processing<br>unit", "application processor", or simply "CPU".<br>[0154] It is to be noted that terms like "node", "computing<br>node", " [0153] Throughout the document, term "user" may be interchangeably referred to as "viewer", "observer", "person", "individual", "end-user", and/or the like. It is to be

server , "cloud server computer , "machine", "host machine " , " device " , " computing device " , " computer " , " computing system " , and the like , may be used interchange ably throughout this document . It is to be further noted that " software program", " package", " software package", and the like, may be used interchangeably throughout this document. Also, terms like "job", "input", "request", "message", and the like, may be used interchangeably throughout this document.

[0155] FIG. 7 illustrates sharing and expansion mechanism 610 of FIG. 6 according to one embodiment. For brevity, many of the details already discussed with reference to FIGS. 1-6 are not repeated or discussed hereafter. In one embodiment, sharing and expansion mechanism 610 may include any number and type of components, such as (with-<br>out limitations): detection/observation logic 701; library generation/mapping logic 703; data sharing/retrieval logic 705; communication/compatibility logic 707; and compression/expansion logic 709.<br>[0156] As aforementioned, data sharing is an efficient manner for having access to

going through long processes and procedures or re-inventing<br>the wheel. However, conventional data sharing techniques<br>are limited in their user as they do not provide for data sharing across processing systems, where one processing system can retrieve any portion of shared data if that portion is relevant to the work being performed by or at the<br>processing system.<br>[0157] Embodiments provide for a novel technique for<br>offering to share data produced on any number and type of

processing systems or devices, such as graphics processor

614 , application processor 612 , field programmable gate array ( $FPGA$ ), application-specific integrated circuit ( $ASIC$ ), and/or the like, using a surface library. For example, in one embodiment, any processor working on the same convolution can retrieve data from the surface library, where this data had been produced by another processor. In another embodiment, this data may be stored persistently across runs.

[0158] In one embodiment, detection/observation logic 701 may be used to detect and observe processors, such as graphics processor 614, as they work on, for example, convolution, where this information may be shared with s example, if graphics processor 614 is working on convolution neural networks (CNNs), graphics processor 614 may be facilitated by sharing/retrieval logic 703 to store the intermediate neural network (NN) data as a data surface in surface library 731 located at one or more cloud databases or datacenters, such as database( $s$ ) 730, in communication with computing devices 600, 740 over one or more communication medium(s) 725, such as a cloud network.<br>[0159] In one embodiment, library logic 703 may be used

to generate one or more surface libraries, such as surface library 731, as desired or necessitated, as graphics processor 614 is detected by detection/observation logic 701 as working on the convolution and having data that is capable of being stored at surface library 731 and subsequently used by one or more other processing devices , such as application processor 612, application processor 742, graphics processor 744, etc. In another embodiment, if there are available sufficient number of or amount of space in surface libraries, additional or new surface libraries may not be generated and data sharing/retrieval logic 703 may simply be used to trigger graphics processor  $614$  to store its data at one of the libraries, such as surface library  $731$ , and establish mapping of the data stored with the corresponding processing device, such as graphics processor 614, as facilitated by library logic 705 .

[ $0160$ ] Once the data has been stored, in one embodiment, it is now available to other processing devices, whether it be application processor  $612$  at computing device  $600$ , or one or more processing devices at another such as application and graphics processors  $742$ ,  $744$  at computing device 740, over communication medium 730. For example, if graphics processor 744 is also working to the same or similar convolution or problem as graphics proces sor 614, graphics processor 744 may access the intermediate NN data at surface library 731 as facilitated by sharing/ ntrieval logic 703.<br>[0161] For example, if computing devices 600, 740 are

two autonomous machines (e.g., vehicles, robots, etc.) that are side-by-side, having a similar view, experiencing the same environmental conditions, etc., the two graphics processors 612, 742, respectively, can shared the NN data using surface library  $731$  at database(s)  $730$  over communication medium( $s$ ) 725.

[0162] Further, in one embodiment, any surfaces produced in this matter may be optionally compressed by compression/expansion logic 709 for an even faster transmission time. Moreover, shared surfaces, as facilitated by sha library 731, may be used for cross-checking results by multiple deep learning systems, such as autonomous machines 600, 740 as further illustrated in FIG. 8A.

[0163] Embodiments further provide for a novel technique<br>for re-expansion of compressed models for achieving high<br>performance and communication efficiency. For example,<br>certain models can be too large to send over the air layers and one or more of the models may be expended back to their original size, such as in high performance computing (HPC).

approach to communicating data models between a number [0164] Conventional techniques are merely limited in their of autonomous machines, such as vehicles, etc., which is inefficient in most cases, such as when there are millions of autonomous vehicles, such as autonomous machines 600 and 740, involved on the road, requiring a large amount of bandwidth to deliver data models over one or more communication medium(s)  $725$ .

[0165] Embodiments provide for a novel technique for compressing data models, while expanding them with an artefact to allow for smooth communication over commu nication medium(s)  $725$  without being costly, such as in terms of system or network resources, bandwidth, etc.

[0166] For example, in case of a large number of autonomous vehicles, such as autonomous machines 600, 740, being on the road, it would be desirable to facilitate communication of information between such vehicles, such as traffic data, whether information, emergency alerts, etc., as<br>frequency and quickly as desired or necessitated. However, conventional systems are not capable of expanding compressed models.

[0167] In one embodiment, as further illustrated with respect to FIG. 8B, compression/expansion logic 709 may be used to compress a data model and assign an artefact to the compressed model such that the artefact serves as both an extension to the compressed model and a form of identification if the compressed model is communicated from one machine 600 to another machine 740 over communication medium(s) 725.

 $[0168]$  It is contemplated that "artefact" or the use of artefact is merely according to one embodiment and that there may be other several techniques by which a compressed model may be re-combined to get the original model, such additional techniques may include (without limitation) using a "light" retraining within the vehicle, "hints" from peer vehicles/drivers, and/or the like. [0169] For example, in one embodiment, an original or r

artefact by compression/expansion logic 709, where this compressed model and the corresponding artefact are com municated from one autonomous machine 600 to another autonomous machine 740 over one or more communication medium $(s)$  725 (e.g., cloud, Internet, etc.). The artefact may then be received at autonomous machine 740 (in autonomous vehicle, for example), followed by the reception of a combo of compressed model and artifact. The two are then separated and autonomous machine 740 can now use the model in its original and uncompressed model.

[0170] Further, communication/compatibility logic 707 may be used to facilitate the needed communication and compatibility between any number of devices of computing<br>device 600 and various components of sharing and expan-<br>sion mechanism 610.<br>[0171] Communication/compatibility logic 707 may be<br>used to facilitate dynamic communicat

between computing device 600 and any number and type of

other computing devices (such as mobile computing device, desktop computer, server computing device, etc.); processing devices or components (such as CPUs, GPUs, etc.); capturing/sensing/detecting devices (such as capturin camera sensors, red green blue ("RGB" or "rgb") sensors, microphones, etc.); display devices (such as output components including display screens, display areas, display projectors, etc.); user/context-awareness components identification/verification sensors/devices (such as biometric sensors/detectors, scanners, etc.); database $(s)$  730, such as memory or storage devices, databases, and/or data sources (such as data storage devices, hard drives, solid-state drives, hard disks, memory cards or devices, memory circuits, etc.); communication medium( $s$ ) 725, such as one or more communication channels or networks (e.g., cloud networks, the Internet, intranets, cellular networks, proximity networks, such as Bluetooth, Bluetooth low energy (BLE), Bluetooth Smart, Wi-Fi proximity, Radio Frequency Identification (RFID), Near Field Communication (NFC), Body Area<br>Network (BAN), etc.); wireless or wired communications and relevant protocols (e.g., Wi-Fi®, WiMAX, Ethernet, etc.); connectivity and location management techniques; software applications/websites (e.g., social and/or business networking websites, etc., business applications, games and other entertainment applications, etc.); and programming languages, etc., while ensuring compatibility with changing technologies, parameters, protocols, standards, etc.

prizy Further, any use of a particular brand, word, term,<br>phrase, name, and/or acronym, such as "detecting", "observ-<br>ing", "training", "selecting", "compressing", "associating",<br>"applying", "sharing", "storing", "retrievi library " , " compressed model " , " expanded model " , " expand ing " , " training set " , " agent " , " machine " , " vehicle " , " robot " , " driving " , " CNN " , " DNN " , " NN " , " execution unit " , " EU " , " shared local memory " , " SLM " , " graphics streams " , " cache " , " graphics cache " , " GPU " , " graphics processor " , « GPU domain " , " GPGPU " , " CPU " , " application proces sor", "CPU domain", "graphics driver", "workload", "application", "graphics pipeline", "pipeline processes", "API", "3D API", "OpenGL®", "DirectX®", "hardware", "software", "agent", "graphics driver", "kernel mode graphics<br>driver", "user-mode driver", "kernel mode graphics<br>driver", "user-mode driver", "task", "process", "operation",<br>"software application", "game", etc., should not be r limit embodiments to software or devices that carry that label in products or in literature external to this document. 2

[0173] It is contemplated that any number and type of components may be added to and/or removed from compression and expansion mechanism 610 to facilitate various<br>embodiments including adding, removing, and/or enhancing<br>certain features. For brevity, clarity, and ease of understand-<br>ing of compression and expansion mechanism computing device, are not shown or discussed here. It is contemplated that embodiments, as described herein, are not limited to any particular technology, topology, system,<br>architecture, and/or standard and are dynamic enough to<br>adopt and adapt to any future changes.<br>[0174] FIG. 8A illustrates a network setup 800 for data<br>sharing and ret

discussed with reference to FIGS. 1-7 may not be discussed or repeated hereafter. Any processes relating to setup 800 may be performed by processing logic that may comprise<br>hardware (e.g., circuitry, dedicated logic, programmable<br>logic, etc.), software (such as instructions run on a process-<br>ing device), or a combination thereof, as facil recited in linear sequences for brevity and clarity in presentation; however, it is contemplated that any number of them can be performed in parallel, asynchronously, or in different orders. Further, embodiments are not limited to any particular architectural placement, framework, setup, or structure of processes and/or components, such as setup 800.

[0175] As illustrated here and described with reference to FIG. 7, the novel technique for data sharing and retrieval as facilitated by sharing and expansion mechanism 610 may hosted and used by any number and type of computing<br>devices, such as autonomous machines 600, 740, 810 (e.g.,<br>vehicles, robots, etc.) over one or more communication<br>mediums 725, such as a cloud network.

[0176] In the illustrated embodiment, database(s) 730, such as cloud databases(s) or datacenter(s), may host one or more surface libraries, such as surface library 731, where data, such as intermediate NN data, may be stored, retrieved, and shared by any number and type of processing devices, such as processors  $612$ ,  $614$ ,  $742$ ,  $744$  and  $812$ ,  $814$  of autonomous machines  $600$ ,  $740$ , and  $810$ , respectively.

[0177] FIG. 8B illustrates a transaction sequence 850 for data sharing and retrieval across processing systems accord ing to one embodiment. For brevity, many of the details previously discussed with reference to FIGS. 1-7 may not be discussed or repeated hereafter. Any processes relating to<br>transaction sequence 850 may be performed by processing<br>logic that may comprise hardware (e.g., circuitry, dedicated logic, programmable logic, etc.), software (such as instructions run on a processing device), or a combination thereof, as facilitated by sharing and expansion mechanism 610 of FIG. 6. The processes associated with transaction sequence **850** may be illustrated or recited in linear sequences for brevity and clarity in presentation; however, it is contemplated that any number of them can be performed in parallel, asynchronously, or in different orders. Further, embodiments are not limited to any particular architectural placement, framework, setup, or structure of processes and/or components, such as transaction sequence 850.

[0178] As illustrated, in one embodiment, original model 851A is selected and the compressed at block 853 and expanded with artefact 855A resulting in compressed model 857A. This compressed model 857A is then transmitted along with artefact from one autonomous machine or any computing device to another autonomous machine or any other computing device at block 859 over communication medium 725 (e.g., cloud network, Internet, proximity network, Bluetooth, etc.).

 $[0179]$  The transmitted compressed model is received at the receiving autonomous machine at 861, where it is received as a combined package 863 having a combination<br>of compressed model 857B (that is the same as compressed model 857A) and artefact 855B (which is the same as artefact 855A). At the autonomous machine, artefact 855B is removed and compressed model 857B is expanded back into uncompressed model 851B (which is the same as uncompressed original model 851).

[0180] FIG. 9 illustrates a method 900 for facilitating data sharing across processing devices using surface library

according to one embodiment. For brevity, many of the details previously discussed with reference to FIGS. 1-8B may not be discussed or repeated hereafter. Any processes relating to method 900 may be performed by processing<br>logic that may comprise hardware (e.g., circuitry, dedicated<br>logic, programmable logic, etc.), software (such as instruc-<br>tions run on a processing device), or a combin as facilitated by sharing and expansion mechanism 610 of FIG. 6. The processes associated with method 900 may be illustrated or recited in linear sequences for brevity and clarity in presentation; however, it is contemplated that any<br>number of them can be performed in parallel, asynchronously, or in different orders.<br>[0181] Method 900 begins at block 901 with detection of

a processor, such as a graphics processor at an autonomous machine, working on a CNN. At block 903, the processor is facilitated to store any intermediate NN data relating to the CNN as a data surface in a surface library at a database, such as a cloud database, over a communication medium, such as a cloud network. In one embodiment, this surface library and other such surface libraries may be created at various database or datacenters and make available to any number<br>and type of processors at various autonomous machines.

block 907, this second processor, such as graphics processor, [0182] At bock 905, another processor of this or another autonomous is working on the same CNN and access the data surface stored at the surface library by the graphics processor. This processor may be another processor may be an application processor at the same or another autonomous machine or another graphics processor or other such processors at the same or another autonomous machine . At at another autonomous machine then accesses the data surface and proceeds with retrieving this stored data to be used for working on the CNN. Further, surface produced this may be optionally compressed for faster transmission time.<br>[0183] Machine Learning Overview

[0184] A machine learning algorithm is an algorithm that can learn based on a set of data. Embodiments of machine learning algorithms can be designed to model high-level tion algorithms can be used to determine which of several categories to which a given input belong; regression algorithms can output a numerical value given an input; and pattern recognition algorithms can be used to gener abstractions within a data set. For example, image recogni-

translated text or perform text to speech and/or speech<br>recognition.<br>[0185] An exemplary type of machine learning algorithm<br>is a neural network. There are many types of neural networks; a simple type of neural network is a feedforward network. A feedforward network may be implemented as an acyclic graph in which the nodes are arranged in layers.<br>Typically, a feedforward network topology includes an input<br>layer and an output layer that are separated by at least one<br>hidden layer. The hidden layer transforms in the input layer into a representation that is useful for generating output in the output layer. The network nodes are fully connected via edges to the nodes in adjacent layers, but there are no edges between nodes within each layer . Data received at the nodes of an input layer of a feedforward network are propagated (i.e., "fed forward") to the nodes of the output layer via an activation function that calculates the states of the nodes of each successive layer in the network based on coefficients ("weights") respectively associated with each of the edges connecting the layers. Depending on

the specific model being represented by the algorithm being executed , the output from the neural network algorithm can take various forms.

[0186] Before a machine learning algorithm can be used to model a particular problem, the algorithm is trained using a training data set. Training a neural network involves selecting a network topology, using a set of training data representing a problem being modeled by the network, and adjusting the weights until the network model performs with a minimal error for all instances of the training data set. For example, during a supervised learning training process for a neural network, the output produced by the network in response to the input representing an instance in a training data set is compared to the "correct" labeled output for that instance, an error signal representing the difference between the output and the labeled output is calculated, and the weights associated with the connections are adjusted to minimize that error as the error signal is backward propagated through the layers of the network. The network is considered "trained" when the errors for each of the outputs generated from the instances of the training data set are minimized .

[0187] The accuracy of a machine learning algorithm can be affected significantly by the quality of the data set used<br>to train the algorithm. The training process can be computationally intensive and may require a significant amount of time on a conventional general-purpose processor. Accordingly, parallel processing hardware is used to train many types of machine learning algorithms. This is particularly<br>useful for optimizing the training of neural networks, as the<br>computations performed in adjusting the coefficients in<br>neural networks lend themselves naturally to mentations. Specifically, many machine learning algorithms<br>and software applications have been adapted to make use of<br>the parallel processing hardware within general-purpose

graphics processing devices.<br>[0188] FIG. 10 is a generalized diagram of a machine<br>learning software stack  $1000$ . A machine learning applica-<br>tion  $1002$  can be configured to train a neural network using<br>a training datase application 1002 can include training and inference functionality for a neural network and/or specialized software that can be used to train a neural network before deployment. The machine learning application 1002 can implement any type of machine intelligence including but not limited to image recognition, mapping and localization, autonomous navigation, speech synthesis, medical imaging, or language translation.

[0189] Hardware acceleration for the machine learning application 1002 can be enabled via a machine learning framework 1004 can provide a library of machine learning primitives. can provide a norary of machine learning primitives.<br>Machine learning primitives are basic operations that are<br>commonly performed by machine learning algorithms.<br>Without the machine learning framework 1004, developers<br>of m plary primitives include tensor convolutions, activation

functions, and pooling, which are computational operations that are performed while training a convolutional neural network (CNN). The machine learning framework 1004 can also provide primitives to implement basic linear algebra subprograms performed by many machine-learning algorithms, such as matrix and vector operations.

[ 0190 ] The machine learning framework 1004 can process input data received from the machine learning application 1002 and generate the appropriate input to a compute framework 1006. The compute framework 1006 can abstract the underlying instructions provided to the GPGPU driver 1008 to enable the machine learning framework 1004 to take advantage of hardware acceleration via the GPGPU hard ware 1010 without requiring the machine learning framework 1004 to have intimate knowledge of the architecture of the GPGPU hardware 1010. Additionally, the compute framework 1006 can enable hardware acceleration for the machine learning framework 1004 across a variety of types and generations of the GPGPU hardware 1010.

[0191] GPGPU Machine Learning Acceleration

[0192] FIG. 11 illustrates a highly-parallel general-purpose graphics processing unit 1100, according to an embodi-<br>ment. In one embodiment, the general-purpose processing unit (GPGPU) 1100 can be configured to be particularly efficient in processing the type of computational workloads associated with training deep neural networks. Additionally, the GPGPU 1100 can be linked directly to other instances of the GPGPU to create a multi-GPU cluster to improve training speed for particularly deep neural networks.

[0193] The GPGPU 1100 includes a host interface 1102 to enable a connection with a host processor. In one embodi-<br>ment, the host interface 1102 is a PCI Express interface. However, the host interface can also be a vendor specific communications interface or communications fabric. The GPGPU 1100 receives commands from the host processor and uses a global scheduler 1104 to distribute execution threads associated with those commands to a set of compute clusters 1106A-H. The compute clusters 1106A-H share a cache memory 1108. The cache memory 1108 can serve as a higher-level cache for cache memories within the compute clusters 1106A-H.

 $[0194]$  The GPGPU 1100 includes memory 1114A-B coupled with the compute clusters 1106A-H via a set of memory controllers 1112A-B. In various embodiments, the memory 1114A-B can include various types of memory memory 1114A-B can include various types of memory<br>devices including dynamic random access memory (DRAM) or graphics random access memory, such as synchronous graphics random access memory (SGRAM), including graphics double data rate (GDDR) memory. In one embodiment, the memory units 224A-N may also include 3D stacked memory , including but not limited to high bandwidth memory ( HBM ) .

 $[0195]$  In one embodiment, each compute cluster GPLAB06A-H includes a set of graphics multiprocessors, such as the graphics multiprocessor 400 of FIG. 4A. The graphics multiprocessors of the compute cluster multiple types of integer and floating point logic units that can perform computational operations at a range of precisions including suited for machine learning computations . For floating-point units in each of the compute clusters 1106A-H can be configured to perform 16-bit or 32-bit floating point including suited for machine learning computations. For

operations , while a different subset of the floating - point units can be configured to perform 64-bit floating point operations .

a couples the GPGPU 1100 with a GPU link 1110 that enables [0196] Multiple instances of the GPGPU 1100 can be configured to operate as a compute cluster. The communication mechanism used by the compute cluster for synchronization and data exchange varies across embodiments. In one embodiment, the multiple instances of the GPGPU 1100 communicate over the host interface 1102. In one embodi ment. the GPGPU 1100 includes an I/O hub 1108 that a direct connection to other instances of the GPGPU. In one embodiment, the GPU link 1110 is coupled to a dedicated GPU-to-GPU bridge that enables communication and synchronization between multiple instances of the GPGPU 1100. In one embodiment, the GPU link 1110 couples with a high-speed interconnect to transmit and receive data to other GPGPUs or parallel processors. In one embodiment, the multiple instances of the GPGPU 1100 are located in separate data processing systems and communicate via a network device that is accessible via the host interface 1102 . In one embodiment, the GPU link 1110 can be configured to enable a connection to a host processor in addition to or as an alternative to the host interface 1102.

[0197] While the illustrated configuration of the GPGPU 1100 can be configured to train neural networks, one embodiment provides alternate configuration of the GPGPU 1100 that can be configured for deployment within a high performance or low power inferencing platform. In an inferencing configuration, the GPGPU 1100 includes fewer<br>of the compute clusters 1106A-H relative to the training configuration. Additionally, memory technology associated with the memory 1114A-B may differ between inferencing and training configurations. In one embodiment, the inferencing configuration of the GPGPU 1100 can support inferencing specific instructions. For example, an inferencing configuration can provide support for one or more 8 during inferencing operations for deployed neural networks.

[0198] FIG. 12 illustrates a multi-GPU computing system 1200, according to an embodiment. The multi-GPU computing system 1200 can include a processor 1202 coupled to multiple GPGPUs 1206A-D via a host interface switch 1204. The host interface switch 1204, in one embodiment, is a PCI express switch device that couples the processor 1202 to a PCI express bus over which the processor 1202 can communicate with the set of GPGPUs 1206A-D. Each of the multiple GPGPUs 1206A-D can be an instance of the GPGPU 1100 of FIG. 11. The GPGPUs 1206A-D can interconnect via a set of high-speed point to point GPU to GPU links 1216. The high-speed GPU to GPU links can connect to each of the GPGPUs 1206A-D via a dedicated GPU link , such as the GPU link 1110 as in FIG . 11. The P2P GPU links 1216 enable direct communication between each of the GPGPUs 1206A-D without requiring communication over the host interface bus to which the processor 1202 is connected. With GPU-to-GPU traffic directed to the P2P GPU links, the host interface bus remains available for system memory access or to communicate with other instances of the multi-GPU computing system 1200, for example, via one or more network devices. While in the illustrated embodiment the GPGPUs 1206A-D connect to the processor 1202 via the host interface switch 1204, in one embodiment the processor 1202 includes direct support for the P2P GPU links 1216 and can connect directly to the GPGPUs 1206A-D.

Machine Learning Neural Network Implementations

[0199] The computing architecture provided by embodiments described herein can be configured to perform the types of parallel processing that is particularly suited for training and deploying neural networks for machine learning. A neural network can be generalized as a network of functions having a graph relationship. As is well-known in the art, there are a variety of types of neural network implementations used in machine learning. One exemplary type of neural network is the feedforward network, as

of the network. The computations for a CNN include applyproduce the output of that filter. Convolution is a specialized type of neural network is the feeding network is the feed of neural network is the Convolutional Neural Network (CNN). A CNN is a specialized feedforward neural network for processing data having a known, grid-like topology, such as image data. Accordingly, CNNs are commonly used for compute vision and image recognition applications, but they also may be used for other types of pattern recognition such as speech and language processing. The nodes in the CNN input layer are organized into a set of "filters" (feature detectors inspired by the receptive fields found in the retina), and the output of each set of filters is propagated to nodes in successive layers ing the convolution mathematical operation to each filter to kind of mathematical operation performed by two functions to produce a third function that is a modified version of one<br>of the two original functions. In convolutional network terminology, the first function to the convolution can be referred to as the input, while the second function can be referred to as the convolution kernel. The output may be referred to as the feature map. For example, the input to a convolution layer can be a multidimensional array of data that defines the various color components of an input image.<br>The convolution kernel can be a multidimensional array of parameters , where the parameters are adapted by the training

[0201] Recurrent neural networks (RNNs) are a family of feedforward neural networks that include feedback connections between layers. RNNs enable modeling of sequential data by sharing parameter data across different parts of the The cycles represent the influence of a present value of a variable on its own value at a future time, as at least a portion of the output data from the RNN is used as feedback for processing subsequent input in a sequence. This feature makes RNNs particularly useful for language processing due to the variable nature in which language data can be com neural network. The architecture for a RNN includes cycles.

[0202] The figures described below present exemplary feedforward, CNN, and RNN networks, as well as describe a general process for respectively training and deploying these descriptions are exemplary and non-limiting as to any<br>specific embodiment described herein and the concepts<br>illustrated can be applied generally to deep neural networks<br>and machine learning techniques in general.

[0203] The exemplary neural networks described above can be used to perform deep learning . Deep learning is machine learning using deep neural networks . The deep

neural networks used in deep learning are artificial neural networks composed of multiple hidden layers , as opposed to layer. Deeper neural networks are generally more computationally intensive to train. However, the additional hidden layers of the network enable multistep pattern recognition that results in reduced output error relative t

mat results in reduced output error relative to shanow<br>machine learning techniques.<br>[0204] Deep neural networks used in deep learning typi-<br>cally include a front-end network to perform feature recog-<br>inition coupled to a b model. Instead, deep neural networks can learn features based on statistical structure or correlation within the input data. The learned features can be provided to a mathematical model that can map detected features to an output. The mathematical model used by the network is generally specialized for the specific task to be performed, and different models will be used to perform different task.

[0205] Once the neural network is structured, a learning model can be applied to the network to train the network to perform specific tasks . The learning model describes how to adjust the weights within the model to reduce the output error of the network . Backpropagation of errors is a common method used to train neural networks . An input vector is presented to the network for processing . The output of the network is compared to the desired output using a loss function and an error value is calculated for each of the neurons in the output layer . The error values are then propagated backwards until each neuron has an associated original output. The network can then learn from those errors using an algorithm, such as the stochastic gradient descent algorithm, to update the weights of the of the neural network.

[0206] FIG. 13A-B illustrate an exemplary convolutional neural network. FIG. 13A illustrates various layers within a CNN. As shown in FIG. 13A, an exemplary CNN used to model image processing can receive input 1302 describing<br>the red, green, and blue (RGB) components of an input<br>image. The input 1302 can be processed by multiple convolutional layers (e.g., convolutional layer  $1304$ , convolutional layer  $1306$ ). The output from the multiple convolutional layers may optionally be processed by a set of fully connected layers 1308. Neurons in a fully connected layer have full connections to all activations in the previous layer, as previously described for a feedforward network. The output from the fully connected layers 1308 can be used to generate an output result from the network. The activations within the fully connected layers 1308 can be computed using matrix multiplication instead of convolution. Not all CNN implementations are make use of fully connected layers DPLA08. For example, in some implementations the convolutional layer 1306 can generate output for the CNN.

[0207] The convolutional layers are sparsely connected, which differs from traditional neural network configuration found in the fully connected layers 1308. Traditional neural network layers are fully connected, such that every output unit interacts with every input unit. However, the convolutional layers are sparsely connected because the output of the convolution of a field is input (instead of the respective state value of each of the nodes in the field) to the nodes of the subsequent layer, as illustrated. The kernels associated<br>with the convolutional layers perform convolution opera-<br>tions, the output of which is sent to the next layer. The<br>dimensionality reduction performed within the

a process large images.<br>
[0208] FIG. 13B illustrates exemplary computation stages<br>
within a convolutional layer of a CNN. Input to a convolu-<br>
tional layer 1312 of a CNN can be processed in three stages<br>
of a convolutional l a convolution stage 1316 , a detector stage 1318 , and a pooling stage 1320. The convolution layer 1314 can then output data to a successive convolutional layer. The final convolutional layer of the network can generate output feature map data or provide input to a fully connected layer, for example, to generate a classification value for the input to the CNN.

[0209] In the convolution stage 1316 performs several convolutions in parallel to produce a set of linear activations. The convolution stage 1316 can include an affine transformation, which is any transformation that can be specified as a linear transformation plus a translation . Affine transfor mations include rotations, translations, scaling, and combinations of these transformations. The convolution stage computes the output of functions (e.g., neurons) that are connected to specific regions in the input, which can be determined as the local region associated with the neuron. The neurons compute a dot product between the weights of the neurons and the region in the local input to which the neurons are connected. The output from the convolution stage 1316 defines a set of linear activations that are processed by successive stages of the convolutional layer 1314. [0210] The linear activations can be processed by a detector stage 1318. In the detector stage 1318, each linear activation is processed by a non-linear activation function.<br>The non-linear activation function increases the nonlinear properties of the overall network without affecting the receptive fields of the convolution layer . Several types of non-linear activation functions may be used. One particular type is the rectified linear unit (ReLU), which uses an activation function defined as  $f(x) = max(0, x)$ , such that the activation is thresholded at zero.

[0211] The pooling stage 1320 uses a pooling function that replaces the output of the convolutional layer 1306 with a summary statistic of the nearby outputs. The pooling function can be used to introduce translation invariance into the neural network, such that small translations to the input do not change the pooled outputs. Invariance to local translation can be useful in scenarios where the presence of a feature in the input data is more important than the precise location of the feature. Various types of pooling functions can be used during the pooling stage 1320, including max pooling, average pooling, and 12-norm pooling. Additionally, some CNN implementations do not include a pooling stage. Instead, such implementations substitute and additional convolution stage having an increased stride relative

[0212] The output from the convolutional layer  $1314$  can then be processed by the next layer 1322. The next layer 1322 can be an additional convolutional layer or one of the fully connected layers 1308. For example, the first convolutional layer 1304 of FIG. 13A can output to the second convolutional layer 1306 , while the second convolutional layer can output to a first layer of the fully connected layers 1308.

[0213] FIG. 14 illustrates an exemplary recurrent neural network 1400. In a recurrent neural network (RNN), the previous state of the network influences the output of the current state of the network. RNNs can be built in a variety of ways using a variety of functions. The use of RNNs generally revolves around using mathematical models to predict the future based on a prior sequence of inputs. For example, an RNN may be used to perform statistical language modeling to predict an upcoming word given a previous sequence of words . The illustrated RNN 1400 can be described has having an input layer 1402 that receives an input vector, hidden layers 1404 to implement a recurrent function, a feedback mechanism 1405 to enable a 'memory' of previous states, and an output layer 1406 to output a result. The RNN 1400 operates based on time-steps. The state of the RNN at a given time step is influenced based on the previous time step via the feedback mechanism 1405. For a given time step, the state of the hidden layers 1404 is defined by the previous state and the input at the current time<br>step. An initial input  $(x_1)$  at a first-time step can be processed by the hidden layer 1404. A second input  $(x_2)$  can be processed by the hidden layer 1404 using state information that is determined during the processing of the initial input  $(x_1)$ . A given state can be computed as  $s = f(Ux_t + Ws_{t-1})$ , where U and W are parameter matrices. The function f is generally a nonlinearity, such as the hyperbolic tangent function (Tan h) or a variant of the rectifier function  $f(x)$  $=$ max $(0, x)$ . However, the specific mathematical function used in the hidden layers 1404 can vary depending on the specific implementation details of the RNN 1400.

[0214] In addition to the basic CNN and RNN networks described, variations on those networks may be enabled. One example RNN variant is the long short term memory (LSTM) RNN. LSTM RNNs are capable of learning long-<br>term dependencies that may be necessary for processing<br>longer sequences of language. A variant on the CNN is a convolutional deep belief network, which has a structure similar to a CNN and is trained in a manner similar to a deep belief network. A deep belief network (DBN) is a generative neural network that is composed of multiple layers of stochastic (random) variables. DBNs can be trained layerby-layer using greedy unsupervised learning. The learned weights of the DBN can then be used to provide pre-train neural networks by determining an optimal initial set of weights for the neural network.

[0215] FIG. 15 illustrates training and deployment of a deep neural network . Once a given network has been struc tured for a task the neural network is trained using a training t dataset 1502. Various training frameworks 1504 have been developed to enable hardware acceleration of the training process. For example, the machine learning framework 1004 of FIG. 10 may be configured as a training framework 1004. The training framework 1004 can hook into an untrained neural network 1506 and enable the untrained neural net to be trained using the parallel processing resources described herein to generate a trained neural net 1508.

[0216] To start the training process the initial weights may be chosen randomly or by pre-training using a deep belief network. The training cycle then be performed in either a supervised or unsupervised manner.

is converging towards a model suitable to generating correct [0217] Supervised learning is a learning method in which training is performed as a mediated operation, such as when the training dataset 1502 includes input paired with the desired output for the input, or where the training dataset includes input having known output and the output of the neural network is manually graded. The network processes the inputs and compares the resulting outputs against a set of expected or desired outputs. Errors are then propagated back through the system. The training framework 1504 can adjust to adjust the weights that control the untrained neural network 1506. The training framework 1504 can provide tools to monitor how well the untrained neural network 1506 answers based on known input data. The training process occurs repeatedly as the weights of the network are adjusted to refine the output generated by the neural network. The training process can continue until the neural network reaches a statistically desired accuracy associated with a trained neural net 1508. The trained neural network 1508 can then be deployed to implement any number of machine

learning operations.<br>
[0218] Unsupervised learning is a learning method in<br>
which the network attempts to train itself using unlabeled data. Thus, for unsupervised learning the training dataset 1502 will include input data without any associated output data. The untrained neural network 1506 can learn groupings<br>within the unlabeled input and can determine how individual<br>inputs are related to the overall dataset. Unsupervised training can be used to generate a self-organizing map, which is a type of trained neural network 1507 capable of performing operations useful in reducing the dimensionality of data.<br>Unsupervised training can also be used to perform anomaly<br>detection, which allows the identification of data points in an input dataset that deviate from the normal patterns of the data.

[0219] Variations on supervised and unsupervised training may also be employed. Semi-supervised learning is a technique in which in the training dataset 1502 includes a a mix of labeled and unlabeled data of the same distribution. Incremental learning is a variant of supervised learning in which<br>input data is continuously used to further train the model. Incremental learning enables the trained neural network 1508 to adapt to the new data 1512 without forgetting the knowledge instilled within the network during initial train-

ing.<br>[ 0220] Whether supervised or unsupervised, the training . process for particularly deep neural networks may be too Instead of using a single compute node, a distributed network of computational nodes can be used to accelerate the

training process.<br>
[0221] FIG. 16 is a block diagram illustrating distributed learning. Distributed learning is a training model that uses multiple distributed computing nodes to perform supervised or unsupervised training of a neural network . The distributed computational nodes can each include one or more host processors and one or more of the general-purpose processing nodes, such as the highly-parallel general-purpose<br>graphics processing unit 1100 as in FIG. 1100. As illustrated,<br>distributed learning can be performed model parallelism<br>1602, data parallelism 1604, or a combination of data parallelism 1604.<br>[ 0222] In model parallelism 1602, different computational

nodes in a distributed system can perform training compu-

tations for different parts of a single network. For example, each layer of a neural network can be trained by a different processing node of the distributed system. The benefits of model parallelism include the ability to large models. Splitting the computations associated with different layers of the neural network enables the training of very large neural networks in which the weights of all layers would not fit into the memory of a single computational<br>node. In some instances, model parallelism can be particularly<br>useful in performing unsupervised training of large neural networks .

[0223] In data parallelism 1604, the different nodes of the distributed network have a complete instance of the model and each node receives a different portion of the data . The results from the different nodes are then combined. While different approaches to data parallelism are possible, data parallel training approaches all require a technique of combining results and synchronizing the model parameters between each node. Exemplary approaches to combining data include parameter averaging and update based data parallelism. Parameter averaging trains each node on a subset of the training data and sets the global parameters (e.g., weights, biases) to the average of the parameters from each node. Parameter averaging uses a central parameter server that maintains the parameter data. Update based data parallelism is similar to parameter averaging except that instead of transferring parameters from the nodes to the Additionally, update based data parallelism can be per-<br>formed in a decentralized manner, where the updates are<br>compressed and transferred between nodes.<br>[0224] Combined model and data parallelism 1606 can be<br>implemented,

each computational node includes multiple GPUs. Each node can have a complete instance of the model with separate GPUs within each node are used to train different

[0225] Distributed training has increased overhead relative to training on a single machine. However, the parallel processors and GPGPUs described herein can each implement various techniques to reduce the overhead of distributed training, including techniques to enable high bandwidth GPU-to-GPU data transfer and accelerated remote data

Synchronization.<br>
[0226] Exemplary Machine Learning Applications<br>
[0227] Machine learning can be applied to solve a variety of technological problems, including but not limited to computer vision, autonomous driving and navigation, speech recognition, and language processing. Computer vision has traditionally been one of the most active research areas for machine learning applications . Applications of computer vision range from reproducing human visual abilities, such as recognizing faces, to creating new categories of visual abilities. For example, computer vision applications can be configured to recognize sound waves from the vibrations induced in objects visible in a video . Parallel vision applications to be trained using significantly larger training dataset than previously feasible and enables inferencing systems to be deployed using low power parallel processors .

[ 0228 ] Parallel processor accelerated machine learning has autonomous driving applications including lane and road sign recognition, obstacle avoidance, navigation, and driving control. Accelerated machine learning techniques can be used to train driving models based on datasets that define the appropriate responses to specific training input. The parallel processors described herein can enable rapid training of the increasingly complex neural networks used for autonomous<br>driving solutions and enables the deployment of low power<br>inferencing processors in a mobile platform suitable for<br>integration into autonomous vehicles.

[0229] Parallel processor accelerated deep neural networks have enabled machine learning approaches to automatic speech recognition (ASR). ASR includes the creation of a function that computes the most probable linguistic sequence given an input acoustic sequence . Accelerated machine learning using deep neural networks have enabled<br>the replacement of the hidden Markov models (HMMs) and Gaussian mixture models (GMMs) previously used for ASR. [0230] Parallel processor accelerated machine learning can also be used to accelerate natural language processing. Automatic learning procedures can make use of statistical inference algorithms to produce models that are robust to erroneous or unfamiliar input. Exemplary natural language processor applications include automatic machine transla

[0231] The parallel processing platforms used for machine<br>learning can be divided into training platforms and deploy-<br>ment platforms. Training platforms are generally highly<br>parallel and include optimizations to accelerate highly-parallel general-purpose graphics processing unit 1100 of FIG. 1100 and the multi-GPU computing system 1200 of FIG. 1200. On the contrary, deployed machine learning platforms generally include lower power parallel processors suitable for use in products such as cameras, processors suitable for use in autonomous vehicles and autonomous vehicles . [0232] FIG . 17 illustrates an exemplary inferencing sys-

tem on a chip (SOC) 1700 suitable for performing inferencing using a trained model. The SOC 1700 can integrate processing components including a media processor 1702, a vision processor 1704, a GPGPU 1706 and a multi-core processor 1708. The SOC 1700 can additionally include on-chip memory 1705 that can enable a shared on-chip data pool that is accessible by each of the processing components. The processing components can be optimized for low power operation to enable deployment to a variety of machine learning platforms, including autonomous vehicles and autonomous robots. For example, one implementation of the SOC 1700 can be used as a portion of the main control system for an autonomous vehicle. Where the SOC 1700 is configured for use in autonomous vehicles the SOC is designed and configured for compliance with the relevant functional safety standards of the deployment jurisdiction.

[0233] During operation, the media processor  $1702$  and vision processor 1704 can work in concert to accelerate computer vision operations. The media processor 1702 can enable low latency decode of multiple high-resolution (e.g., 4K, 8K) video streams. The decoded video streams can be written to a buffer in the on-chip-memory 1705. The vision processor 1704 can then parse the decoded video and perform preliminary processing operations on the frames of the decoded video in preparation of processing the frames using a trained image recognition model. For example, the vision processor 1704 can accelerate convolution operations for a CNN that is used to perform image recognition on the high-resolution video data, while back end model computations are performed by the GPGPU 1706. [0234] The multi-core processor 1708 can include control logic to assist with sequencing and synchronization of data

transfers and shared memory operations performed by the media processor 1702 and the vision processor 1704. The multi-core processor 1708 can also function as an application processor to execute software applications that can make use of the inferencing compute capability of the GPGPU 1706. For example, at least a portion of the navigation and driving logic can be implemented in software executing on the multi-core processor 1708. Such software can directly issue computational workloads to the GPGPU 1706 or the computational workloads can be issued to the multi-core processor 1708, which can offload at least a multion of those operations to the GPGPU 1706.<br>[0235] The GPGPU 1706 can include compute clusters

such as a low power configuration of the compute clusters 1106A-1106H within the highly-parallel general-purpose graphics processing unit 1100. The compute clusters within the GPGPU 1706 can support instruction that are specifically optimized to perform inferencing computations on a trained neural network. For example, the GPGPU 1706 can support instructions to perform low precision computations such as 8-bit and 4-bit integer vector operations.<br>[ 0236] System Overview II

[0237] FIG. 18 is a block diagram of a processing system 1800, according to an embodiment. In various embodiments, the system 1800 includes one or more processors 1802 and one or more graphics processors 1808 , and may be a single processor desktop system, a multiprocessor workstation system, or a server system having a large number of processors 1802 or processor cores 1807. In on embodiment, the system  $1800$  is a processing platform incorporated within a system - on-a-chip (SoC) integrated circuit for use in mobile, handheld, or embedded devices.

[0238] An embodiment of system 1800 can include, or be incorporated within a server-based gaming platform, a game console, including a game and media console, a mobile gaming console, a handheld game console, or an online game console. In some embodiments system 1800 is a mobile phone, smart phone, tablet computing device or mobile Internet device . Data processing system 1800 can also include, couple with, or be integrated within a wearable device, such as a smart watch wearable device, smart eyewear device, augmented reality device, or virtual reality device. In some embodiments, data processing system 1800 is a television or set top box device having one or more processors 1802 and a graphical interface generated by one or more graphics processors 1808. incorporated within a server-based gaming platform, a game

[0239] In some embodiments, the one or more processors 1802 each include one or more processor cores 1807 to process instructions which, when executed, perform operations for system and user software. In some embodiments, each of the one or more processor cores 1807 is configured to process a specific instruction set 1809. In some embodiments, instruction set 1809 may facilitate Complex Instruction Set Computing (CISC), Reduced Instruction Set Computing (RISC), or computing via a Very Long Instruction Word (VLIW). Multiple processor cores 1807 may each process a different instruction set 1809 , which may include instructions to facilitate the emulation of other instruction sets. Processor core 1807 may also include other processing devices, such a Digital Signal Processor (DSP).

[0240] In some embodiments, the processor 1802 includes cache memory 1804. Depending on the architecture, the processor 1802 can have a single internal cache or multiple levels of internal cache. In some embodiments, the cache memory is shared among various components of the pro cessor 1802. In some embodiments , the processor 1802 also uses an external cache (e.g., a Level-3  $(L3)$  cache or Last Level Cache (LLC)) (not shown), which may be shared among processor cores 1807 using known cache coherency techniques. A register file 1806 is additionally included in processor 1802 which may include different types of registers for storing different types of data (e.g., integer registers, floating point registers, status registers, and an instruction pointer register). Some registers may be general-purpose registers, while other registers may be specific to the design of the processor 1802.

[0241] In some embodiments, processor  $1802$  is coupled<br>to a processor bus  $1810$  to transmit communication signals<br>such as address, data, or control signals between processor<br> $1802$  and other components in system  $1800$ . architecture, including a memory controller hub 1816 and an Input Output ( $I/O$ ) controller hub 1830. A memory controller hub 1816 facilitates communication between a memory device and other components of system 1800, while an I/O<br>Controller Hub (ICH) 1830 provides connections to I/O devices via a local I/O bus. In one embodiment, the logic of the memory controller hub 1816 is integrated within the processor

[0242] Memory device 1820 can be a dynamic random access memory (DRAM) device, a static random access<br>memory (SRAM) device, flash memory device, phasechange memory device, or some other memory device having suitable performance to serve as process memory. In one embodiment, the memory device 1820 can operate as system memory for the system 1800, to store data 1822 and instructions 1821 for use when the one or more processors 1802 executes an application or process. Memory controller hub 1816 also couples with an optional external graphics processor 1812, which may communicate with the one or more graphics processors 1808 in processors 1802 to perform graphics and media operations.

[0243] In some embodiments, ICH 1830 enables peripherals to connect to memory device 1820 and processor 1802 via a high-speed I/O bus. The I/O peripherals include, but are not limited to, an audio controller 1846, a firmware interface 1828, a wireless transceiver 1826 (e.g., Wi-Fi, Bluetooth), a data storage device 1824 (e.g., hard disk drive, flash memory, etc.), and a legacy I/O controller 1840 for coupling legacy (e.g., Personal System 2 (PS/2)) devices to the system. One or more Universal Serial Bus (USB) controllers 1842 connect input devices, such as keyboard and mouse 1844 combinations. A network controller 1834 may also couple to ICH 1830. In some embodiments, a high-<br>performance network controller (not shown) couples to<br>processor bus 1810. It will be appreciated that the system 1800 shown is exemplary and not limiting, as other types of data processing systems that are differently configured may also be used. For example, the I/O controller hub 1830 may be integrated within the one or more proces memory controller hub  $1816$  and  $I/O$  controller hub  $1830$ may be integrated into a discreet external graphics proces sor, such as the external graphics processor 1812.

[0244] FIG. 19 is a block diagram of an embodiment of a processor 1900 having one or more processor cores 1902A 1902N, an integrated memory controller 1914, and an integrated graphics processor 1908. Those elements of FIG. 19 having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such. Processor 1900 can include additional cores up to and including additional core 1902N represented by the dashed lined boxes . Each of processor cores 1902A-1902N includes one or more internal cache units 1904A-1904N. In some embodiments, each processor core also has access to one or more shared cached units 1906 .

[0245] The internal cache units 1904A-1904N and shared cache units 1906 represent a cache memory hierarchy within the processor 1900. The cache memory hierarchy may include at least one level of instruction and data cache within each processor core and one or more levels of shared mid-level cache, such as a Level 2 (L2), Level 3 (L3), Level 4 (L4), or other levels of cache, where the highest level of cache before external memory is classified as the LLC . In some embodiments, cache coherency logic maintains coherency between the various cache units 1906 and 1904A 1904N.

[0246] In some embodiments, processor 1900 may also include a set of one or more bus controller units 1916 and a system agent core 1910. The one or more bus controller units 1916 manage a set of peripheral buses, such as one or more Peripheral Component Interconnect buses (e.g., PCI, PCI Express). System agent core 1910 provides management functionality for the various processor components. In some embodiments, system agent core 1910 includes one or more integrated memory controllers 1914 to manage access to various external memory devices (not shown).

[0247] In some embodiments, one or more of the processor cores 1902A-1902N include support for simultaneous multi-threading. In such embodiment, the system agent core 1910 includes components for coordinating and operating cores 1902A-1902N during multi-threaded processing. System agent core 1910 may additionally include a power control unit (PCU), which includes logic and components to regulate the power state of processor cores 1902A-1902N and graphics processor 1908.

[ 0248 ] In some embodiments , processor 1900 additionally includes graphics processor 1908 to execute graphics pro cessing operations . In some embodiments , the graphics processor 1908 couples with the set of shared cache units 1906 , and the system agent core 1910 , including the one or . more integrated memory controllers 1914. In some embodi ments, a display controller 1911 is coupled with the graphics processor 1908 to drive graphics processor output to one or more coupled displays. In some embodiments, display controller 1911 may be a separate module coupled with the graphics processor via at least one interconnect, or may be integrated within the graphics processor 1908 or system agent core 1910.

[ 0249 ] In some embodiments , a ring based interconnect unit 1912 is used to couple the internal components of the processor 1900. However, an alternative interconnect unit may be used, such as a point-to-point interconnect, a switched interconnect, or other techniques, including techniques well known in the art. In some embodiments, graphics processor 1908 couples with the ring interconnect 1912 via an I/O link 1913.

[0250] The exemplary I/O link 1913 represents at least one of multiple varieties of I/O interconnects, including an on-package I/O interconnect which facilitates communication between various processor components and a high-<br>performance embedded memory module 1918, such as an eDRAM module. In some embodiments, each of the processor cores 1902-1902N and graphics processor 1908 use embedded memory modules 1918 as a shared Last Level Cache.

[0251] In some embodiments, processor cores 1902A-1902N are homogenous cores executing the same instruction set architecture. In another embodiment, processor cores 1902A-1902N are heterogeneous in terms of instruction set architecture (ISA), where one or more of processor cores 1902A-N execute a first instruction set, while at least one of the other cores executes a subset of the first instruction set or a different instruction set . In one embodiment processor cores 1902A-1902N are heterogeneous in terms of micro-<br>architecture, where one or more cores having a relatively higher power consumption couple with one or more power cores having a lower power consumption. Additionally, processor 1900 can be implemented on one or more chips or as an SoC integrated circuit having the illustrated compo nents, in addition to other components.

[ 0252 ] FIG . 20 is a block diagram of a graphics processor 2000 , which may be a discrete graphics processing unit , or may be a graphics processor integrated with a plurality of processing cores. In some embodiments, the graphics processor communicates via a memory mapped I/O interface to registers on the graphics processor and with commands placed into the processor memory. In some embodiments, graphics processor 2000 includes a memory interface 2014 to access memory. Memory interface 2014 can be an interface to local memory, one or more internal caches, one or more shared external caches, and/or to system memory.

[0253] In some embodiments, graphics processor 2000 also includes a display controller 2002 to drive display output data to a display device 2020. Display controller 2002 includes hardware for one or more overlay planes for the interface elements. In some embodiments, graphics processor 2000 includes a video codec engine 2006 to encode. decode, or transcode media to, from, or between one or more media encoding formats, including, but not limited to Moving Picture Experts Group (MPEG) formats such as MPEG-2, Advanced Video Coding (AVC) formats such as H.264/ MPEG-4 AVC, as well as the Society of Motion Picture & Television Engineers (SMPTE)  $421M/VC-1$ , and Joint Photographic Experts Group (JPEG) formats such as JPEG, and Motion JPEG (MJPEG) formats.

[0254] In some embodiments, graphics processor 2000 includes a block image transfer (BLIT) engine 2004 to perform two-dimensional (2D) rasterizer operations including, for example, bit-boundary block transfers. However, in one embodiment, 2D graphics operations are performed using one or more components of graphics processing<br>engine (GPE) 2010. In some embodiments, graphics pro-<br>cessing engine 2010 is a compute engine for performing<br>graphics operations, including three-dimensional (3D)<br>graphic [0255] In some embodiments, GPE 2010 includes a 3D pipeline 2012 for performing 3D operations, such as rendering three-dimensional images and scenes using processing functions that act upon 3D primitive shapes (e.g., rectangle, triangle, etc.). The 3D pipeline 2012 includes programmable and fixed function elements that perform various tasks within the element and/or spawn execution threads to a 3D/Media sub-system 2015. While 3D pipeline 2012 can be used to perform media operations, an embodi-<br>ment of GPE  $2010$  also includes a media pipeline  $2016$  that is specifically used to perform media operations, such as video post-processing and image enhancement.

[0256] In some embodiments, media pipeline 2016 includes fixed function or programmable logic units to perform one or more specialized media operations, such as video decode acceleration, video de-interlacing, and video encode acceleration in place of, or on behalf of video codec engine 2006. In some embodiments, media pipeline 2016 additionally includes a thread spawning unit to spawn threads for execution on 3D/Media sub-system 2015. The spawned threads perform computations for the media operations on one or more graphics execution units included in 3D/Media sub-system 2015.

[0257] In some embodiments, 3D/Media subsystem 2015 includes logic for executing threads spawned by 3D pipeline 2012 and media pipeline 2016. In one embodiment, the pipelines send thread execution requests to 3D/Media subsystem 2015, which includes thread dispatch logic for arbitrating and dispatching the various requests to available thread execution resources . The execution resources include an array of graphics execution units to process the 3D and 2015 includes one or more internal caches for thread instructions and data. In some embodiments, the subsystem also includes shared memory, including registers and addressable memory , to share data between threads and to store output data .

[0258] 3D/Media Processing<br>[0259] FIG. 21 is a block diagram of a graphics processing<br>engine 2110 of a graphics processor in accordance with<br>some embodiments. In one embodiment, the graphics processing engine ( GPE  $\,$  2110 is a version of the GPE 2010 shown in FIG. 20. Elements of FIG. 21 having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such. For example, the 3D pipeline 2012 and media pipeline  $2016$  of FIG. 20 are illustrated. The media pipeline 2016 is optional in some embodiments of the GPE 2110 and may not be explicitly included within the GPE 2110. For example, be explicited with the GPE 2110.<br>For explicit in a separate media and/or image processor is coupled to the GPE 2110.<br>For all  $(0.260)$  In some embodiments, GPE 2110 couples with or

includes a command streamer 2103, which provides a command stream to the 3D pipeline 2012 and/or media pipelines 2016. In some embodiments, command streamer 2103 is coupled with memory, which can be system memory, or one or more of internal cache memory and shared cache memory. In some embodiments, command streamer 2103 receives commands from the memory and sends the com mands to 3D pipeline 2012 and/or media pipeline 2016. The commands are directives fetched from a ring buffer, which stores commands for the 3D pipeline 2012 and media pipeline 2016. In one embodiment, the ring buffer can additionally include batch command buffers storing batches<br>of multiple commands. The commands for the 3D pipeline<br>2012 can also include references to data stored in memory,<br>such as but not limited to vertex and geometry da pipeline 2016 process the commands and data by performing<br>operations via logic within the respective pipelines or by<br>dispatching one or more execution threads to a graphics core array 2114.

[0261] In various embodiments, the 3D pipeline 2012 can execute one or more shader programs, such as vertex shaders, geometry shaders, pixel shaders, fragment shaders, compute shaders, or other shader programs, by processing<br>the instructions and dispatching execution threads to the<br>graphics core array 2114. The graphics core array 2114 provides a unified block of execution resources . Multi purpose execution logic (e.g., execution units) within the graphic core array 2114 includes support for various 3D API shader languages and can execute multiple simultaneous execution threads associated with multiple shaders.

[0262] In some embodiments, the graphics core array  $2114$  also includes execution logic to perform media functions, such as video and/or image processing. In one embodiment,<br>the execution units additionally include general-purpose<br>logic that is programmable to perform parallel general<br>purpose computational operations, in addition to gra processing operations. The general-purpose logic can perform processing operations in parallel or in conjunction with general purpose logic within the processor core(s)  $1807$  of FIG. 18 or core  $1902A-1902N$  as in FIG. 19.

[0263] Output data generated by threads executing on the graphics core array 2114 can output data to memory in a unified return buffer (URB) 2118. The URB 2118 can store data for multiple threads. In some embodiments, the URB 2118 may be used to send data between different threads executing on the graphics core array 2114. In some embodi-<br>ments, the URB 2118 may additionally be used for synchromization between threads on the graphics core array and fixed function logic within the shared function logic 2120. [0264] In some embodiments, graphics core array 2114 is scalable, such that the array includes a variable number of graphics cores, each having a variable number of execution units based on the target power and performance level of GPE 2110. In one embodiment, the execution resources are dynamically scalable, such that execution resources may be enabled or disabled as needed.

[0265] The graphics core array 2114 couples with shared function logic 2120 that includes multiple resources that are shared between the graphics cores in the graphics core array.<br>The shared functions within the shared function logic 2120 are hardware logic units that provide specialized supplemental functionality to the graphics core array 2114. In various embodiments, shared function logic 2120 includes but is not limited to sampler 2121, math 2122, and interthread communication ( $ITC$ ) 2123 logic. Additionally, some embodiments implement one or more cache( $s$ ) 2125 within the shared function logic 2120. A shared function is implemented where the demand for a given specialized function is insufficient for inclusion within the graphics core array 2114. Instead a single instantiation of that specialized function is implemented as a stand-alone entity in the shared function logic 2120 and shared among the execution resources within the graphics core array 2114. The precise set of functions that are shared between the graphics core array 2114 and included within the graphics core array 2114

a varies between embodiments.<br>[0266] FIG. 22 is a block diagram of another embodiment<br>of a graphics processor 2200. Elements of FIG. 22 having<br>the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

[0267] In some embodiments, graphics processor  $2200$  includes a ring interconnect  $2202$ , a pipeline front-end  $2204$ , a media engine 2237, and graphics cores 2280A-2280N. In some embodiments, ring interconnect 2202 couples the graphics processor to other processing units, including other graphics processors or one or more general-purpose processor cores. In some embodiments, the graphics processor is one of many processors integrated within a multi-core

ments, media engine 2237 includes a Video Quality Engine processing system.<br>
[0268] In some embodiments, graphics processor 2200<br>
receives batches of commands via ring interconnect 2202. The incoming commands are interpreted by a command streamer  $2203$  in the pipeline front-end  $2204$ . In some embodiments, graphics processor  $2200$  includes scalable execution logic to perform 3D geometry processing and media processing via the graphics core(s) 2280A-2280N. For 3D geometry processing commands, command streamer 2203 supplies commands to geometry pipeline 2236. For at least some media processing commands, command streamer 2203 supplies the commands to a video front end 2234, which couples with a media engine 2237. In some embodi-(VQE)  $2230$  for video and image post-processing and a multi-format encode/decode (MFX)  $2233$  engine to provide hardware-accelerated media data encode and decode. In some embodiments, geometry pipeline 2236 and media engine 2237 each generate execution threads for the thread execution resources provided by at least one graphics core 2280A.

[0269] In some embodiments, graphics processor 2200 includes scalable thread execution resources featuring modular cores 2280A-2280N (sometimes referred to as core slices), each having multiple sub-cores 2250A-2250N, 2260A-2260N (sometimes referred to as core sub-slices). In some embodiments, graphics processor 2200 can have an number of graphics cores 2280A through 2280N. In some embodiments, graphics processor 2200 includes a graphics core 2280A having at least a first sub-core 2250A and a second core sub-core 2260A. In other embodiments, the graphics processor is a low power processor with a single sub-core (e.g.,  $2250$ A). In some embodiments, graphics processor 2200 includes multiple graphics cores 2280A-2280N, each including a set of first sub-cores 2250A-2250N and a set of second sub-cores 2260A-2260N. Each sub-core in the set of first sub-cores 2250A-2250N includes at least a first set of execution units 2252A-2252N and media/texture samplers 2254A-2254N. Each sub-core in the set of second sub-cores 2260A-2260N includes at least a second set of execution units 2262A-2262N and samplers 2264A-2264N. In some embodiments, each sub-core 2250A-2250N, 2260A-2260N shares a set of shared resources 2270A-2270N. In some embodiments, the shared resources include shared cache memory and pixel operation logic. Other shared resources may also be included in the various embodiments of the graphics processor.

[0270] Execution Logic<br>[0271] FIG. 23 illustrates thread execution logic 2300 including an array of processing elements employed in some embodiments of a GPE. Elements of FIG. 23 having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

instruction cache  $2300$ , data port  $2314$ , sampler  $2310$ , and [0272] In some embodiments, thread execution logic 2300 includes a pixel shader 2302, a thread dispatcher 2304, instruction cache 2306, a scalable execution unit array including a plurality of execution units 2308A-2308N, sampler 2310, a data cache 2312, and a data port 2314. In one embodiment, the included components are interconnected via an interconnect fabric that links to each of the 2300 includes one or more connections to memory, such as system memory or cache memory , through one or more of execution unit array 2308A-2308N. In some embodiments, each execution unit (e.g. 2308A) is an individual vector processor capable of executing multiple simultaneous threads and processing multiple data elements in parallel for each thread. In some embodiments, execution unit array 2308A-2308N includes any number individual execution units .

[0273] In some embodiments, execution unit array 2308A-2308N is primarily used to execute "shader" programs. In some embodiments, the execution units in array 2308A-<br>2308N execute an instruction set that includes native support for many standard 3D graphics shader instructions, such that shader programs from graphics libraries (e.g., Direct 3D and OpenGL) are executed with a minimal translation. The execution units support vertex and geometry processing (e.g., vertex programs, geometry programs, vertex shaders), pixel processing (e.g., pixel shaders, fragment shaders) and general-purpose processing (e.g., compute and media shaders ) .

[0274] Each execution unit in execution unit array 2308A-2308N operates on arrays of data elements. The number of data elements is the "execution size," or the number of channels for the instruction. An execution channel is a logical unit of execution for data element access, masking, and flow control within instructions. The number of channels may be independent of the number of physical Arithmetic Logic Units (ALUs) or Floating Point Units (FPUs) for a particular graphics processor. In some embodiments,<br>execution units 2308A-2308N support integer and floating-<br>point data types.<br>[0275] The execution unit instruction set includes single<br>instruction multiple data (SIM

tiple thread (SIMT) instructions. The various data elements can be stored as a packed data type in a register and the execution unit will process the various elements based on the data size of the elements. For example, when operating on a 256-bit wide vector, the 256 bits of the vector are stored in a register and the execution unit operates on the vector as four separate 64-bit packed data elements (Quad-Word (QW) size data elements), eight separate 32-bit packed data elements (Double Word (DW) size data elements), sixteen separate 16-bit packed data elements (Word (W) size data elements), or thirty-two separate 8-bit data elements (byte (B) size data elements). However, different vector widths and register sizes are possible .

 $[0276]$  One or more internal instruction caches  $(e.g., 2306)$ are included in the thread execution logic 2300 to cache thread instructions for the execution units. In some embodiments, one or more data caches (e.g., 2312) are included to cache thread data during thread execution. In some embodi-<br>ments, sampler 2310 is included to provide texture sampling for 3D operations and media sampling for media operations.<br>In some embodiments, sampler 2310 includes specialized<br>texture or media sampling functionality to process texture or media data during the sampling process before providing the

[0277] During execution, the graphics and media pipelines send thread initiation requests to thread execution logic 2300 via thread spawning and dispatch logic. In some embodiments, thread execution logic 2300 includes a local thread dispatcher 2304 that arbitrates thread initiation requests from the graphics and media pipelines and instantiates the requested threads on one or more execution units 2308A 2308N. For example, the geometry pipeline (e.g., 2236 of FIG. 22) dispatches vertex processing, tessellation, or geometry processing threads to thread execution logic 2300 (FIG. 23). In some embodiments, thread dispatcher 2304 can also process runtime thread spawning requests from the executing shader programs.

[0278] Once a group of geometric objects has been processed and rasterized into pixel data, pixel shader  $2302$  is invoked to further compute output information and cause results to be written to output surfaces (e.g., co depth buffers, stencil buffers, etc.). In some embodiments, pixel shader 2302 calculates the values of the various vertex attributes that are to be interpolated across the rasterized object. In some embodiments, pixel shader 2302 then executes an application programming interface (API)-supplied pixel shader program. To execute the pixel shader program, pixel shader 2302 dispatches threads to an execution unit (e.g., 2308A) via thread dispatcher 2304. In some embodiments, pixel shader 2302 uses texture sampling logic in sampler  $2310$  to access texture data in texture maps stored<br>in memory. Arithmetic operations on the texture data and the input geometry data compute pixel color data for each geometric fragment, or discards one or more pixels from further processing.

[ $0279$ ] In some embodiments, the data port  $2314$  provides a memory access mechanism for the thread execution logic 2300 output processed data to memory for processing on a graphics processor output pipeline. In some embodiments, the data port  $2314$  includes or couples to one or more cache memories (e.g., data cache  $2312$ ) to cache data for memory access via the data port.

[ $0280$ ] FIG. 24 is a block diagram illustrating a graphics processor instruction formats  $2400$  according to some embodiments. In one or more embodiment, the graphics processor execution units support an instruction set having instructions in multiple formats. The solid lined boxes illustrate the components that are generally included in an execution unit instruction, while the dashed lines include components that are optional or that are only included in a sub-set of the instructions. In some embodiments, instruction format 2400 described and illustrated are macro-instructions, in that they are instructions supplied to the execution unit, as opposed to micro-operations resulting from instruction decode once the instruction is processed.

[0281] In some embodiments, the graphics processor execution units natively support instructions in a 128-bit

instruction format 2410. A 64-bit compacted instruction format 2430 is available for some instructions based on the selected instruction, instruction options, and number of operands. The native 128-bit instruction format 2410 provides access to all instruction options, while some options and operations are restricted in the 64-bit instruction format 2430. The native instructions available in the 64-bit instruction format 2430 vary by embodiment. In some embodiments, the instruction is compacted in part using a set of index values in an index field 2413. The execution unit hardware references a set of compaction tables based on the index values and uses the compaction table outputs to reconstruct a native instruction in the 128-bit instruction format 2410.

[0282] For each format, instruction opcode 2412 defines the operation that the execution unit is to perform. The execution units execute each instruction in parallel across the multiple data elements of each operand. For example, in response to an add instruction the execution unit performs a resenting a texture element or picture element. By default, the execution unit performs each instruction across all data channels of the operands. In some embodiments, instruction control field 2414 enables control over certain execution options, such as channels selection (e.g., predication) and data channel order (e.g., swizzle). For 128-bit instructions 2410 an exec-size field 2416 limits the number of data channels that will be executed in parallel. In some embodiments, exec-size field 2416 is not available for use in the 64-bit compact instruction format 2430.

[0283] Some execution unit instructions have up to three operands including two source operands, src0 2420, src1 2422 , and one destination 2418. In some embodiments , the execution units support dual destination instructions, where<br>one of the destinations is implied. Data manipulation instructions can have a third source operand (e.g., SRC2 2424), where the instruction opcode 2412 determines the number of source operands. An instruction's last source

operand can be an immediate (e.g., hard-coded) value passed<br>with the instruction.<br>[0284] In some embodiments, the 128-bit instruction for-<br>mat 2410 includes an access/address mode information 2426 specifying, for example, whether direct register addressing<br>mode or indirect register addressing mode is used. When<br>direct register addressing mode is used, the register address<br>of one or more operands is directly provided instruction 2410.

[0285] In some embodiments, the 128-bit instruction format 2410 includes an access/address mode field 2426, which specifies an address mode and/or an access mode for the instruction. In one embodiment, the access mode to define a data access alignment for the instruction. Some embodiments support access modes including a 16-byte aligned access mode and a 1-byte aligned access mode, where the byte alignment of the access mode determines the access alignment of the instruction operands. For example, when in<br>a first mode, the instruction 2410 may use byte-aligned addressing for source and destination operands and when in<br>a second mode, the instruction  $2410$  may use 16-bytealigned addressing for all source and destination operands. [0286] In one embodiment, the address mode portion of the access/address mode field 2426 determines whether the instruction is to use direct or indirect addressing. When direct register addressing mode is used bits in the instruction 2410 directly provide the register address of one or more operands. When indirect register addressing mode is used, the register address of one or more operands may be computed based on an address register value and an address immediate field in the instruction.

[0287] In some embodiments, instructions are grouped based on opcode 2412 bit-fields to simplify Opcode decode 2440. For an 8-bit opcode, bits 4, 5, and 6 allow the execution unit to determine the type of opcode. The precise opcode grouping shown is merely an example. In some embodiments, a move and logic opcode group  $2442$  includes data movement and logic instructions (e.g., move (mov), compare (cmp)). In some embodiments, move and logic group 2442 shares the five most significant bits (MSB), where move (mov) instructions are in the form of 0000xxxxb and logic instructions are in the form of 0001xxxxb. A flow control instruction group 2444 (e.g., call, jump (jmp)) includes instructions in the form of 0010xxxxb (e.g., 0x20). A miscellaneous instruction group  $2446$ includes a mix of instructions, including synchronization<br>instructions (e.g., wait, send) in the form of 0011xxxxb (e.g.,<br>0x30). A parallel math instruction group 2448 includes<br>component-wise arithmetic instructions (e.g.  $(mul)$ ) in the form of 0100xxxxb (e.g., 0x40). The parallel math group 2448 performs the arithmetic operations in parallel across data channels . The vector math group 2450 includes arithmetic instructions (e.g.,  $dp4$ ) in the form of 0101xxxxb (e.g.,  $0x50$ ). The vector math group performs out the vector such as dot product calculations on vector operands.<br>
[0288] Graphics Pipeline

[0289] FIG. 25 is a block diagram of another embodiment of a graphics processor 2500. Elements of FIG. 25 having the same reference numbers (or names) as the elements of any other figure herein can operate or function in any manner similar to that described elsewhere herein, but are not limited to such.

[ $0290$ ] In some embodiments, graphics processor  $2500$  includes a graphics pipeline  $2520$ , a media pipeline  $2530$ , a display engine  $2540$ , thread execution logic  $2550$ , and a render output pipeline 2570. In some embodiments, graphics processor 2500 is a graphics processor within a multi-core processing system that includes one or more general-purpose processing cores . The graphics processor is controlled by register writes to one or more control registers ( not shown) or via commands issued to graphics processor 2500 via a ring interconnect 2502. In some embodiments, ring interconnect 2502 couples graphics processor 2500 to other processing components, such as other graphics processors or general-purpose processors. Commands from ring interconnect 2502 are interpreted by a command streamer 2503,

which supplies instructions to individual components of<br>graphics pipeline 2520 or media pipeline 2530.<br>[0291] In some embodiments, command streamer 2503<br>directs the operation of a vertex fetcher 2505 that reads<br>vertex data commands provided by command streamer 2503. In some embodiments, vertex fetcher 2505 provides vertex data to a vertex shader 2507, which performs coordinate space transformation and lighting operations to each vertex. In some embodiments, vertex fetcher 2505 and vertex shader 2507 execute vertex-processing instructions by dispatching execution threads to execution units 2552A, 2552B via a thread dispatcher 2531 .

[0292] In some embodiments, execution units 2552A, 2552B are an array of vector processors having an instruction set for performing graphics and media operations . In some embodiments, execution units 2552A, 2552B have an attached L1 cache 2551 that is specific for each array or shared between the arrays. The cache can be configured as a data cache, an instruction cache, or a single cache that is

partitioned to contain data and instructions in different<br>partitions.<br>[0293] In some embodiments, graphics pipeline 2520<br>includes tessellation components to perform hardware-ac-<br>celerated tessellation of 3D objects. In som a programmable hull shader 2511 configures the tessellation operations. A programmable domain shader 2517 provides back-end evaluation of tessellation output. A tessellator 2513 operates at the direction of hull shader 2511 and contains special purpose logic to generate a set of detailed provided as input to graphics pipeline 2520. In some embodiments, if tessellation is not used, tessellation components 2511, 2513, 2517 can be bypassed.

[0294] In some embodiments, complete geometric objects can be processed by a geometry shader 2519 via one or more threads dispatched to execution units 2552A, 2552B, or can proceed directly to the clipper 2529. In some embodiments,<br>the geometry shader operates on entire geometric objects,<br>rather than vertices or patches of vertices as in previous<br>stages of the graphics pipeline. If the tessel the geometry shader 2519 receives input from the vertex shader 2507. In some embodiments, geometry shader 2519 is programmable by a geometry shader program to perform<br>geometry tessellation if the tessellation units are disabled.<br>[0295] Before rasterization, a clipper 2529 processes vertex data. The clipper  $2529$  may be a fixed function clipper or a programmable clipper having clipping and geometry depth test component 2573 in the render output pipeline<br>2570 dispatches pixel shaders to convert the geometric<br>objects into their per pixel representations. In some embodiments, pixel shader logic is included in thread execution logic 2550. In some embodiments, an application can bypass rasterization and access un-rasterized vertex data via a stream out unit 2523 .

[0296] The graphics processor 2500 has an interconnect bus, interconnect fabric, or some other interconnect mechanism that allows data and message passing amongst the major components of the processor. In some embodiments, execution units 2552A, 2552B and associated cache(s) 2551, texture and media sampler 2554, and texture/sampler cache 2558 interconnect via a data port 2556 to perform memory access and communicate with render output pipeline components of the processor. In some embodiments, sampler 2554, caches 2551, 2558 and execution units 2552A, 2552B each have separate memory access paths.

[ $0297$ ] In some embodiments, render output pipeline 2570 contains a rasterizer and depth test component 2573 that converts vertex-based objects into an associated pixel-based<br>representation. In some embodiments, the render output presentation . In some embody embody includes a windower/masker unit to perform fixed function triangle and line rasterization. An associated render cache 2578 and depth cache 2579 are also available in some embodiments. A pixel operations component 2577 performs pixel-based operations on the data, though in some instances, pixel operations associated with 2D operations

( e.g. bit block image transfers with blending ) are performed by the 2D engine 2541 , or substituted at display time by the display controller 2543 using overlay display planes . In some embodiments, a shared  $\overline{L}3$  cache 2575 is available to all graphics components, allowing the sharing of data without the use of main system memory.

[0298] In some embodiments, graphics processor media pipeline  $2530$  includes a media engine 2537 and a video front end 2534. In some embodiments, video front end 2534 receives pipeline commands from the command streamer 2503. In some embodiments, media pipeline 2530 includes a separate command streamer. In some embodiments, video front-end 2534 processes media commands before sending the command to the media engine 2537. In some embodi-<br>ments, media engine 2537 includes thread spawning functionality to spawn threads for dispatch to thread execution logic 2550 via thread dispatcher 2531.

[0299] In some embodiments, graphics processor  $2500$  includes a display engine  $2540$ . In some embodiments, display engine  $2540$  is external to processor  $2500$  and couples with the graphics processor via the ring interconnect 2502, or some other interconnect bus or fabric. In some embodiments, display engine 2540 includes a 2D engine 2541 and a display controller 2543. In some embodiments, display engine 2540 contains special purpose logic capable of operating independently of the 3D pipeline. In some embodiments, display controller 2543 couples with a device (not shown), which may be a system integrated display device, as in a laptop computer, or an external display device attached via a display device connector.

[0300] In some embodiments, graphics pipeline 2520 and<br>media pipeline 2530 are configurable to perform operations<br>based on multiple graphics and media programming inter-<br>faces and are not specific to any one application pr ming interface (API). In some embodiments, driver software<br>for the graphics processor translates API calls that are<br>specific to a particular graphics or media library into commands that can be processed by the graphics processor. In some embodiments, support is provided for the Open Graphics Library (OpenGL) and Open Computing Language (OpenCL) from the Khronos Group, the Direct3D library from the Microsoft Corporation, or support may be provided to both OpenGL and D3D. Support may also be provided for the Open Source Computer Vision Library (OpenCV). A future API with a compatible 3D pipeline would also be supported if a mapping can be made from the pipeline of the future API to the pipeline of the graphics processor.

[0301] Graphics Pipeline Programming

[ $0302$ ] FIG.  $26A$  is a block diagram illustrating a graphics processor command format 2600 according to some embodi ments. FIG. 26B is a block diagram illustrating a graphics processor command sequence 2610 according to an embodi ment. The solid lined boxes in FIG. 26A illustrate the components that are generally included in a graphics com mand while the dashed lines include components that are optional or that are only included in a sub-set of the graphics commands. The exemplary graphics processor command format 2600 of FIG. 26A includes data fields to identify a target client 2602 of the command, a command operation target client 2602 of the command , a command operation code ( opcode ) 2604 , and the relevant data 2606 for the 2 command. A sub-opcode 2605 and a command size 2608 are also included in some commands.

[0303] In some embodiments, client 2602 specifies the client unit of the graphics device that processes the com

mand data . In some embodiments , a graphics processor command parser examines the client field of each command to condition the further processing of the command and route the command data to the appropriate client unit. In some embodiments, the graphics processor client units include a memory interface unit, a render unit, a 2D unit, a<br>3D unit, and a media unit. Each client unit has a corresponding processing pipeline that processes the commands. Once the command is received by the client unit , the client unit reads the opcode 2604 and, if present, sub-opcode 2605 to determine the operation to perform. The client unit performs the command using information in data field 2606. For some commands an explicit command size 2608 is expected to specify the size of the command. In some embodiments, the command parser automatically determines the size of at least some of the commands based on the command opcode . In some embodiments, commands are aligned via multiples of a double word.

[0304] The flow diagram in FIG. 26B shows an exemplary graphics processor command sequence 2610. In some embodiments, software or firmware of a data processing system that features an embodiment of a graphics processor uses a version of the command sequence shown to set up, execute, and terminate a set of graphics operations. A sample command sequence is shown and described for purposes of example only as embodiments are not limited to these<br>specific commands or to this command sequence. Moreover, the commands may be issued as batch of commands in a command sequence, such that the graphics processor will process the sequence of commands in at least partially concurrence .

[0305] In some embodiments, the graphics processor command sequence 2610 may begin with a pipeline flush command 2612 to cause any active graphics pipeline to complete the currently pending commands for the pipeline.<br>In some embodiments, the 3D pipeline 2622 and the media<br>pipeline 2624 do not operate concurrently. The pipeline flush is performed to cause the active graphics pipeline to complete any pending commands. In response to a pipeline flush, the command parser for the graphics processor will pause<br>command processing until the active drawing engines com-<br>plete pending operations and the relevant read caches are invalidated. Optionally, any data in the render cache that is marked 'dirty' can be flushed to memory. In some embodi-<br>ments, pipeline flush command 2612 can be used for pipeline synchronization or before placing the graphics processor into a low power state.

[0306] In some embodiments, a pipeline select command 2613 is used when a command sequence requires the graphics processor to explicitly switch between pipelines. In some embodiments, a pipeline select command 2613 is required only once within an execution context before issuing pipeline commands unless the context is to issue commands for both pipelines. In some embodiments, a pipeline flush command is 2612 is required immediately before a pipeline switch via the pipeline select command 2613.

[0307] In some embodiments, a pipeline control command 2614 configures a graphics pipeline for operation and is used to program the 3D pipeline 2622 and the media pipeline 2624. In some embodiments, pipeline control command 2614 configures the pipeline state for the active pipeline . In one embodiment, the pipeline control command 2614 is used for pipeline synchronization and to clear data from one or more cache memories within the active pipeline before processing a batch of commands. [0308] In some embodiments, commands for the return

buffer state 2616 are used to configure a set of return buffers for the respective pipelines to write data. Some pipeline operations require the allocation, selection, or configuration<br>of one or more return buffers into which the operations write<br>intermediate data during processing. In some embodiments,<br>the graphics processor also uses one or to store output data and to perform cross thread communi cation. In some embodiments, configuring the return buffer state 2616 includes selecting the size and number of return buffers to use for a set of pipeline operations.

[0309] The remaining commands in the command sequence differ based on the active pipeline for operations. Based on a pipeline determination 2620, the command sequence is tailored to the 3D pipeline 2622 beginning with the 3D pipeline state 2630, or the media pipeline 2624

beginning at the media pipeline state  $2640$ .<br>[0310] The commands for the 3D pipeline state  $2630$  include 3D state setting commands for vertex buffer state, vertex element state, constant color state, depth buffer state, and other state variables that are to be configured before 3D primitive commands are processed. The values of these commands are determined at least in part b lar 3D API in use. In some embodiments, 3D pipeline state 2630 commands are also able to selectively disable or bypass certain pipeline elements if those elements will not be used .

[0311] In some embodiments, 3D primitive 2632 command is used to submit 3D primitives to be processed by the 3D pipeline. Commands and associated parameters that are passed to the graphics processor via the 3D primitive 2632 command are forwarded to the vertex fetch function in the graphics pipeline. The vertex fetch function uses the 3D primitive 2632 command data to generate vertex data struc tures. The vertex data structures are stored in one or more return buffers. In some embodiments, 3D primitive 2632 command is used to perform vertex operations on 3D primitives via vertex shaders. To process vertex shaders, 3D

pipeline 2622 dispatches shader execution threads to graphics processor execution units.<br>[0312] In some embodiments, 3D pipeline 2622 is triggered via an execute 2634 command or event. In some embodiments, a register write In some embodiments execution is triggered via a 'go' or 'kick' command in the command sequence. In one embodi-<br>ment command execution is triggered using a pipeline synchronization command to flush the command sequence<br>through the graphics pipeline. The 3D pipeline will perform<br>geometry processing for the 3D primitives. Once operations are complete, the resulting geometric objects are rasterized and the pixel engine colors the resulting pixels . Additional commands to control pixel shading and pixel back end

[0313] In some embodiments, the graphics processor com-<br>mand sequence 2610 follows the media pipeline 2624 path when performing media operations. In general, the specific use and manner of programming for the media pipeline 2624 depends on the media or compute operations to be performed. Specific media decode operations may be offloaded to the media pipeline during media decode . In some embodi ments, the media pipeline can also be bypassed and media decode can be performed in whole or in part using resources provided by one or more general-purpose processing cores.<br>In one embodiment, the media pipeline also includes elements for general-purpose graphics processor unit (GPGPU) operations, where the graphics processor is used to perform SIMD vector operations using computational shader programs that are not explicitly related to the rendering of

graphics primitives.<br>[ $0314$ ] In some embodiments, media pipeline  $2624$  is configured in a similar manner as the 3D pipeline 2622. A set of commands to configure the media pipeline state 2640 are dispatched or placed into a command queue before the media object commands 2642. In some embodiments, commands for the media pipeline state 2640 include data to configure the media pipeline elements that will be used to process the media objects. This includes data to configure the video decode and video encode logic within the media pipeline, such as encode or decode format. In some embodiments, commands for the media pipeline state 2640 also support the use of one or more pointers to "indirect" state elements that contain a batch of state settings.

[0315] In some embodiments, media object commands 2642 supply pointers to media objects for processing by the media pipeline. The media objects include memory buffers containing video data to be processed. In some embodiments, all media pipeline states must be valid before issuing a media object command 2642. Once the pipeline state is configured and media object commands  $2642$  are queued, the media pipeline  $2624$  is triggered via an execute command  $2644$  or an equivalent execute event (e.g., register write). Output from media pipeline 2624 may then be post processed by operations provided by the 3D pipeline 2622 or the media pipeline 2624. In some embodiments, GPGPU operations are configured and executed in a similar manner a as media operations.<br>
[0316] Graphics Software Architecture

[0317] FIG. 27 illustrates exemplary graphics software architecture for a data processing system 2700 according to a some embodiments. In some embodiments, software architecture includes a 3D graphics application 2710, an operating system 2720, and at least one processor 2730. In some embodiments, processor 2730 includes a graphics processor

2732 and one or more general-purpose processor core(s)  $2734$ . The graphics application 2710 and operating system  $2720$  each execute in the system memory  $2750$  of the data processing system. [0318] In some embodiments, Level Shader Language (HLSL) or the OpenGL Shader Language (GLSL). The application also includes executable instructions 2714 in a machine language suitable for execution by the general purpose processor core (s)  $2734$ . The application also includes graphics objects 2716 defined by vertex data.

[0319] In some embodiments, operating system 2720 is a Microsoft® Windows® operating system from the Microsoft Corporation, a proprietary UNIX-like operating<br>system, or an open source UNIX-like operating system<br>using a variant of the Linux kernel. The operating system<br>2720 can support a graphics API 2722 such as the Direc the operating system 2720 uses a front-end shader compiler 2724 to compile any shader instructions 2712 in HLSL into a lower-level shader language. The compilation may be a just-in-time (JIT) compilation or the application can perform shader pre-compilation. In some embodiments, high-level shaders are compiled into low-level shaders during the

with a kernel mode graphics driver  $2729$ . In some embodicompilation of the 3D graphics application 2710.<br>[0320] In some embodiments, user mode graphics driver 2726 contains a back-end shader compiler 2727 to convert the shader instructions 2712 into a hardware specific representation. When the OpenGL API is in use, shader instructions 2712 in the GLSL high-level language are passed to a user mode graphics driver 2726 for compilation. In some embodiments, user mode graphics driver 2726 uses operating system kernel mode functions 2728 to communicate ments, kernel mode graphics driver 2729 communicates with graphics processor 2732 to dispatch commands and instructions.

## [0321] IP Core Implementations

[0322] One or more aspects of at least one embodiment may be implemented by representative code stored on a machine-readable medium which represents and/or defines logic within an integrated circuit such as a processor. For example, the machine-readable medium may include instructions which represent various logic within the processor. When read by a machine, the instructions may cause the machine to fabricate the logic to perform the techniques described herein. Such representations, known as "IP cores," are reusable units of logic for an integrated circuit that may be stored on a tangible, machine-readable medium as a hardware model that describes the structure of the integrated circuit. The hardware model may be supplied to various customers or manufacturing facilities, which load the hardware model on fabrication machines that manufacture the integrated circuit . The integrated circuit may be fabricated such that the circuit performs operations described in association with any of the embodiments described herein.

ciation with any of the embodiments described herein.<br>[0323] FIG. 28 is a block diagram illustrating an IP core development system 2800 that may be used to manufacture an integrated circuit to perform operations according to an embodiment . The IP core development system 2800 may be used to generate modular, re-usable designs that can be incorporated into a larger design or used to construct an entire integrated circuit (e.g., an SOC integrated circuit). A design facility 2830 can generate a software simulation 2810 of an IP core design in a high-level programming language (e.g.,  $C/C_{++}$ ). The software simulation 2810 can be used to design, test, and verify the behavior of the IP core using a simulation model 2812. The simulation model 2812 may include functional, behavioral, and/or timing simulations. A register transfer level (RTL) design 2815 can then be created or synthesized from the simulation model 2812. The RTL design 2815 is an abstraction of the behavior of the inte grated circuit that models the flow of digital signals between using the modeled digital signals. In addition to an RTL design 2815, lower-level designs at the logic level or transistor level may also be created, designed, or synthesized. Thus, the particular details of the initial design and simulation may vary.

[0324] The RTL design 2815 or equivalent may be further synthesized by the design facility into a hardware model 2820, which may be in a hardware description language (HDL), or some other representation of physical design data. The HDL may be further simulated or tested to verify the IP core design . The IP core design can be stored for delivery to a  $3^{rd}$  party fabrication facility 2865 using non-volatile memory 2840 (e.g., hard disk, flash memory, or any nonvolatile storage medium). Alternatively, the IP core design may be transmitted (e.g., via the Internet) over a wired connection 2850 or wireless connection 2860. The fabrica tion facility 2865 may then fabricate an integrated circuit that is based at least in part on the IP core design. The fabricated integrated circuit can be configured to perform<br>operations in accordance with at least one embodiment<br>described herein.<br>[0325] Exemplary System on a Chip Integrated Circuit

[0326] FIGS. 29-31 illustrate exemplary integrated circuits and associated graphics processors that may be fabri cated using one or more IP cores, according to various embodiments described herein. In addition to what is illus-<br>trated, other logic and circuits may be included, including additional graphics processors/cores, peripheral interface controllers, or general purpose processor cores.

[0327] FIG. 29 is a block diagram illustrating an exemplary system on a chip integrated circuit 2900 that may be fabricated using one or more IP cores, according to an embodiment. Exemplary integrated circuit 2900 includes one or more application processor(s) 2905 (e.g., CPUs), at least one graphics processor 2910, and may additionally<br>include an image processor 2915 and/or a video processor<br>2920, any of which may be a modular IP core from the same or multiple different design facilities . Integrated circuit 2900 includes peripheral or bus logic including a USB controller 2925, UART controller 2930, an SPI/SDIO controller 2935, and an  $I^2S/I^2C$  controller **2940**. Additionally, the integrated circuit can include a display device **2945** coupled to one or more of a high-definition multimedia interface (HDMI) controller 2950 and a mobile industry processor interface (MIPI) display interface 2955. Storage may be provided by a flash memory subsystem 2960 including flash memory and a flash memory controller. Memory interface may be provided via a memory controller 2965 for access to SDRAM or SRAM memory devices. Some integrated circuits additionally include an embedded security engine 2970.

[0328] FIG. 30 is a block diagram illustrating an exemplary graphics processor 3010 of a system on a chip integrated circuit that may be fabricated using one or more IP cores, according to an embodiment. Graphics processor 29. Graphics processor 3010 includes a vertex processor 3005 and one or more fragment processor(s) 3015A-3015N (e.g., 3015A, 3015B, 3015C, 3015D, through 3015N-1, and 3015N). Graphics processor 3010 can execute different shader programs via separate logic, such that the vertex processor 3005 is optimized to execute operations for vertex shader programs, while the one or more fragment processor (s)  $3015A-3015N$  execute fragment (e.g., pixel) shading operations for fragment or pixel shader programs. The vertex processor  $3005$  performs the vertex processing stage of the 3D graphics pipeline and generates primitives and vertex data. The fragment processor(s) 3015A-3015N use the primitive and vertex data generated by the vertex pro cessor 3005 to produce a framebuffer that is displayed on a display device. In one embodiment, the fragment processor  $\delta$  3015A-3015N are optimized to execute fragment shader programs as provided for in the OpenGL API, which may be used to perform similar operations as a pixel shader program as provided for in the Direct 3D API. 3010 can be a variant of the graphics processor 2910 of FIG.
MMUs within the system, including one or more MMUs<br>associated with the one or more application processor(s)<br>**2905**, image processor **2915**, and/or video processor **2920** of [0329] Graphics processor 3010 additionally includes one or more memory management units (MMUs) 3020A-3020B,  $cache(s)$  3025A-3025B, and circuit interconnect( $s$ ) 3030A-3030B. The one or more MMU( $s$ ) 3020A-3020B provide for virtual to physical address mapping for graphics processor 3010, including for the vertex processor 3005 and/or fragment processor $(s)$  3015A-3015N, which may reference vertex or image/texture data stored in memory, in addition to vertex or image/texture data stored in the one or more  $cache(s) 3025A-3025B$ . In one embodiment, the one or more  $MMU(s)$  3020A-3020B may be synchronized with other MMUs within the system, including one or more MMUs FIG . 29 , such that each processor 2905-2920 can participate in a shared or unified virtual memory system . The one or more circuit interconnect(s)  $3030A - 3030B$  enable graphics processor 3010 to interface with other IP cores within the SoC, either via an internal bus of the SoC or via a direct connection, according to embodiments.

more  $MMU(s)$  3020A-3020B, cache(s) 3025A-3025B, and [ $0330$ ] FIG. 31 is a block diagram illustrating an additional exemplary graphics processor  $3110$  of a system on a chip integrated circuit that may be fabricated using one or more IP cores, according to an embodiment. Graphics processor 3110 can be a variant of the graphics processor 2910 of FIG. 29. Graphics processor 3110 includes the one or circuit interconnect(s)  $3030A - 3030B$  of the integrated circuit 3000 of FIG. 30.

[ $0331$ ] Graphics processor  $3110$  includes one or more shader core(s)  $3115A.3115N$  (e.g.,  $3115A.3115B.3115C$ , 3115D, 3115E, 3115F, through 3015N-1, and 3015N), which provides for a unified shader core architecture in which a mable shader code, including shader program code to implement vertex shaders, fragment shaders, and/or compute shaders. The exact number of shader cores present can vary among embodiments and implementations. Additionally, graphics processor 3110 includes an inter-core task manager 3105, which acts as a thread dispatcher to dispatch execution threads to one or more shader core(s) 3115A-3115N. Graphics processor 3110 additionally includes a tiling unit 3118 to accelerate tiling operations for tile-based rendering, in which rendering operations for a scene are subdivided in image space. Tile-based rendering can be used to exploit local spatial coherence within a scene or to optimize use of internal caches .

[0332] References to "one embodiment", "an embodiment", "example embodiment", "various embodiments", etc., indicate that the embodiment(s) so described may include particular features, structures, or characteristics, but not every embodiment necessarily includes the particular features, structures, or characteristics. Further, some embodiments may have some, all, or none of the features described for other embodiments.

[0333] In the foregoing specification, embodiments have been described with reference to specific exemplary embodi-<br>ments thereof. It will, however, be evident that various modifications and changes may be made thereto without departing from the broader spirit and scope of embodiments<br>as set forth in the appended claims. The Specification and drawings are, accordingly, to be regarded in an illustrative rather than a restrictive sense.

[0334] In the following description and claims, the term "coupled" along with its derivatives, may be used. " Coupled" is used to indicate that two or more elements co-operate or interact with each other, but they may or may not have intervening physical or electrical components between them.

[0335] As used in the claims, unless otherwise specified the use of the ordinal adjectives "first", "second", "third", etc., to describe a common element, merely indicate that different instances of like elements are being referred to, and<br>are not intended to imply that the elements so described must be in a given sequence, either temporally, spatially, in ranking, or in any other manner.

[ $0336$ ] The following clauses and/or examples pertain to further embodiments or examples. Specifics in the examples may be used anywhere in one or more embodiments . The various features of the different embodiments or examples may be variously combined with some features included and others excluded to suit a variety of different applications. Examples may include subject matter such as a method, means for performing acts of the method, at least one machine-readable medium including instructions that, when performed by a machine cause the machine to perform acts hybrid communication according to embodiments and examples described herein.<br>[ 0337] Some embodiments pertain to Example 1 that

includes an apparatus to facilitate sharing of data and compression and expansion of models, the apparatus comprising: detection/observation logic to detect a first processor processing information relating to a neural network at the apparatus, wherein the first processor comprises a first graphics processor and the apparatus comprises a first autonomous machine; and data sharing/retrieval logic to facilitate the first processor to store one or more portions of the information in a library at a database, wherein the one or more portions are accessible to a second processor of a

computing device.<br> **[0338]** Example 2 includes the subject matter of Example 1, wherein the data sharing/retrieval logic is further to facilitate the second processor to access and retrieve the one or more portions of the information from the library when the second processor performs tasks relating to the neural network, wherein the neural network includes a convolu-

tional neural network (CNN), wherein second processor<br>comprises a second graphics processor and the computing<br>device comprises a second autonomous machine.<br>[0339] Example 3 includes the subject matter of Examples<br>1-2, furt generation/mapping logic is further to map the one or more portions of the information to first processor.

[0340] Example 4 includes the subject matter of Examples 1-3, wherein the first and second autonomous machines include autonomous vehicles in communication over one or more communication mediums including a cloud network, wherein the database includes a cloud database.

[ 0341 ] Example 5 includes the subject matter of Examples 1-4 , further comprising compression / expansion logic to compress a model with an item, wherein the compressed model along with the item is communicated over to the second autonomous machine over the one or more commu nication mediums .

[ 0342 ] Example 6 includes the subject matter of Examples 1-5 , wherein the compression / expansion logic is further to facilitate reception of the compressed model and the item at the second autonomous machine, wherein the compressed model is uncompressed using the item, wherein the item incudes one or more of an artefact, a light, and a hint.

[0343] Example 7 includes the subject matter of Examples 1-6, wherein the first graphics processor is co-located with an application processor on a common semiconductor pack age .

includes a method for facilitating sharing of data and com-[0344] Some embodiments pertain to Example 8 that includes a method for facilitating sharing of data and compression and expansion of models, the method comprising: detecting a first processor processing information relati first processor comprises a first graphics processor and the first computing device comprises a first autonomous machine; and facilitating the first processor to store one or more portions of the information in a library at a database, wherein the one or more portions are accessible to a second processor of a computing device.

[ $0345$ ] Example 9 includes the subject matter of Example 8, further comprising facilitating the second processor to access and retrieve the one or more portions of the information from the library when the second processor performs tasks relating to the neural network, wherein the neural network includes a convolutional neural network (CNN). wherein second processor comprises a second graphics processor and the computing device comprises a second

[0346] Example 10 includes the subject matter of Examples 8-9 , further comprising generating the library at the database; and mapping the one or more portions of the information to first processor.

[0347] Example 11 includes the subject matter of Examples 8-10, wherein the first and second autonomous machines include autonomous vehicles in communication over one or more communication mediums including a cloud network, wherein the database includes a cloud database .

[0348] Example 12 includes the subject matter of Examples 8-11, further comprising compressing a model with an item, wherein the compressed model along with the item is communicated over to the second autonomous machine over the one or more communication mediums performance counters that are regarded as faulty or outside

[0349] Example 13 includes the subject matter of Examples 8-12 , further comprising facilitating reception of the compressed model and the item at the second autono mous machine, wherein the compressed model is uncompressed using the item, wherein the item incudes one or more of an artefact, a light, and a hint.

[0350] Example 14 includes the subject matter of Examples 8-13 , wherein the first graphics processor is co-located with an application processor on a common semiconductor package.

[0351] Some embodiments pertain to Example 15 that includes a graphics processing system comprising a computing device having memory coupled to a processor, the processor to: detect a first processor processing information relating to a neural network at the first computing device, wherein the first processor comprises a first graphics processor and the first computing device comprises a first autonomous machine; and facilitating the first processor to store one or more portions of the information in a library at a database, wherein the one or more portions are accessible to a second processor of a computing device.

[0352] Example 16 includes the subject matter of Example 15, wherein the processor is further to facilitate the second processor to access and retrieve the one or more portions of the information from the library when the second processor performs tasks relating to the neural network, wherein the neural network includes a convolutional neural network graphics processor and the computing device comprises a

second autonomous machine.<br>[0353] Example 17 includes the subject matter of Example 15-16, wherein the processor is further to facilitate generating the library at the database; and mapping the one or more portions of the information to first processor.

more communication mediums including a cloud network, wherein the database includes a cloud database. [0354] Example 18 includes the subject matter of Example 15-17 , wherein the first and second autonomous machines include autonomous vehicles in communication over one or more communication mediums including a cloud network,

[ 0355 ] Example 19 includes the subject matter of Examples 15-18 , wherein the operations further comprise compressing a model with an item , wherein the compressed model along with the item is communicated over to the second autonomous machine over the one or more commu nication mediums .

[0356] Example 20 includes the subject matter of Examples 15-19, wherein the operations further comprise facilitating reception of the compressed model and the item<br>at the second autonomous machine, wherein the compressed model is uncompressed using the item, wherein the item incudes one or more of an artefact, a light, and a hint.

[0357] Example 21 includes the subject matter of Examples 15-20, wherein the first graphics processor is co-located with an application processor on a common

semiconductor package.<br>
[0358] Example 22 includes at least one non-transitory or tangible machine-readable medium comprising a plurality of instructions, when executed on a computing device, to implement or perform a method as claimed in any of claims or examples  $8-14$ .

[0359] Example 23 includes at least one machine-readable medium comprising a plurality of instructions, when executed on a computing device, to implement or perform a method as claimed in any of claims or examples 8-14.

[0360] Example 24 includes a system comprising a mechanism to implement or perform a method as claimed in any of claims or examples  $8-14$ .

[ $0361$ ] Example 25 includes an apparatus comprising means for performing a method as claimed in any of claims or examples  $8-14$ .

 $[0362]$  Example 26 includes a computing device arranged to implement or perform a method as claimed in any of

[0363] Example 27 includes a communications device arranged to implement or perform a method as claimed in

[0364] Example 28 includes at least one machine-readable medium comprising a plurality of instructions, when executed on a computing device , to implement or perform a method or realize an apparatus as claimed in any preceding claims .

[0366] Example 30 includes a system comprising a mechanism to implement or perform a method or realize an apparatus as claimed in any preceding claims.

[0367] Example 31 includes an apparatus comprising means to perform a method as claimed in any preceding claims .

[0368] Example 32 includes a computing device arranged to implement or perform a method or realize an apparatus as claimed in any preceding claims.

[0369] Example 33 includes a communications device arranged to implement or perform a method or realize an

[0370] The drawings and the forgoing description give examples of embodiments. Those skilled in the art will appreciate that one or more of the described elements may well be combined into a single functional element. Alternatively, certain elements may be split into multiple functional elements. Elements from one embodiment may be added to another embodiment. For example, orders of processes described herein may be changed and are not limited to the manner described herein. Moreover, the actions of any flow diagram need not be implemented in the order shown; nor do all of the acts necessarily need to be performed. Also, those acts that are not dependent on other acts may be performed in parallel with the other acts. The scope of embodiments is by no means limited by these specific examples. Numerous variations, whether explicitly given in the specification or not, such as differences in structure, dimension, and use of material, are possible. The scope of embodiments is at least as broad as given by the following claims.

1-20. (canceled)  $21$ . An apparatus comprising:

- a transmitter and a receiver for a first autonomous vehicle to transmit and receive data over one or more communication mediums with one or more other autonomous vehicles, the one or more other autonomous vehicles including a second autonomous vehicle; and<br>one or more processors including a graphics processor,
- wherein the one or more processors are to:<br>establish communications with the second autonomous
	- establish communications with the second autonomous vehicle and determine to share data with the second autonomous vehicle , process a neural network model for transmission to the
	- second autonomous vehicle, the processing of the neural network model including compressing the neural network model to generate a compressed neural network model, and
	- transmit the compressed neural network model from the first autonomous vehicle to the second autono mous vehicle over the one or more communication mediums .

22. The apparatus of claim 21, wherein processing the neural network model for transmission further includes assigning an artefact to the compressed neural network model, the artefact to identify the compressed neural network model, and wherein the one or more processors are to communicate the artefact and the compressed neural net work model to the second autonomous vehicle.

23. The apparatus of claim 21, wherein the one or more communication mediums include a cloud network.

24. The apparatus of claim  $21$ , wherein the one or more processors are further to detect the second autonomous vehicle.

25. The apparatus of claim 21, wherein determining to share the data with the second autonomous vehicle is based on one or more of :

- location of the first autonomous vehicle and the second autonomous vehicle;
- the first autonomous vehicle and the second autonomous vehicle having a similar view; or
- the first autonomous vehicle and the second autonomous vehicle experiencing same environmental conditions.

26. The apparatus of claim 21, wherein the data to be shared with the second autonomous vehicle includes one or more of traffic data, weather information, and emergency alerts .

27. The apparatus of claim 21, wherein the neural network model includes a convolutional neural network (CNN) model.

28. The apparatus of claim 21, wherein the compression of the neural network model includes reducing a number of layers from the neural network model to generate the compressed neural network model.<br>29. The apparatus of claim 21, wherein the graphics

processor is co-located with an application processor on a common semiconductor package.

30. At least one non-transitory machine-readable medium comprising instructions that when executed by one or more processors , cause the one or more processors to perform operations comprising :

- establishing communications between a first autonomous vehicle and a second autonomous vehicle;<br>determining to share data between the first autonomous
- vehicle and the second autonomous vehicle;
- processing a neural network model at the first autonomous vehicle for transmission to the second autonomous<br>vehicle, the processing of the neural network model including compressing the neural network model to generate a compressed neural network model , the com ing a number of layers from the neural network model to generate the compressed model; and
- transmitting the compressed neural network model from the first autonomous vehicle to the second autonomous<br>vehicle over one or more communication mediums.

31. The at least one non-transitory machine-readable medium of claim 30, wherein processing the neural network model for transmission further includes assigning an artefact<br>to the compressed neural network model, the artefact to identify the compressed neural network model, and wherein the one or more processors are to communicate the artefact and the compressed neural network model to the second

32. The at least one non-transitory machine-readable medium of claim 30, wherein the one or more communication mediums include a cloud network.

33. The at least one non-transitory machine-readable medium of claim 30, wherein the instructions further include instructions that when executed by the one or more proces sors, cause the one or more processors to perform operations comprising:

detecting the second autonomous vehicle.

34. The at least one non-transitory machine-readable medium of claim 30, wherein determining to share the data with the second autonomous vehicle is based on one or more  $\alpha$ f

- location of the first autonomous vehicle and the second autonomous vehicle:
- the first autonomous vehicle and the second autonomous vehicle having a similar view; or
- the first autonomous vehicle and the second autonomous<br>vehicle experiencing same environmental conditions.

35. The at least one non-transitory machine-readable medium of claim 30, wherein the data to be shared between the first autonomous vehicle and the second autonomous vehicle includes one or more of traffic data, weather information, and emergency alerts.

36. A method comprising:

- establishing communications between a first autonomous vehicle and a second autonomous vehicle;
- receiving a compressed neural network model at the first autonomous vehicle from the second autonomous vehicle over one or more communication mediums , the

compressed neural network model including data to be shared between the first autonomous vehicle and the second autonomous vehicle; and

- expanding the compressed neural network model to generate an original neural network model.
- 37. The method of claim 36, wherein the one or more communication mediums include a cloud network.

38. The method of claim 36, further comprising:

- receiving, at the first autonomous vehicle, an artefact from the second autonomous vehicle;
- wherein expanding the received compressed neural network model including applying the artefact in expanincluding applying applying applying the compressed neural network model.<br> **39**. The method of claim 38, wherein the artefact and the

compressed model are received separately by the first

40. The method of claim 38, wherein the artefact serves as both an extension to the compressed model and a form of a identification for the compressed neural network model in communication .

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