

US011839086B2

(12) United States Patent Petti et al.

(54) 3-DIMENSIONAL MEMORY STRING ARRAY OF THIN-FILM FERROELECTRIC TRANSISTORS

(71) Applicant: SUNRISE MEMORY

CORPORATION, San Jose, CA (US)

(72) Inventors: Christopher J. Petti, Mountain View,

CA (US); Vinod Purayath, Sedona, AZ (US); George Samachisa, Atherton, CA (US); Wu-Yi Henry Chien, San Jose, CA (US); Eli Harari, Saratoga, CA

(US)

(73) Assignee: SUNRISE MEMORY
CORPORATION, San Jose, CA (US)

(*) Notice: Subject to any disclaimer, the term of this

patent is extended or adjusted under 35 U.S.C. 154(b) by 0 days.

(21) Appl. No.: 17/812,375

(22) Filed: Jul. 13, 2022

(65) **Prior Publication Data**

US 2023/0027837 A1 Jan. 26, 2023

Related U.S. Application Data

- (60) Provisional application No. 63/222,926, filed on Jul. 16, 2021.
- (51) Int. Cl. H10B 51/30 (2023.01) G11C 11/22 (2006.01) H10B 51/20 (2023.01)
- (52) U.S. Cl.

(10) Patent No.: US 11,839,086 B2

(45) **Date of Patent:**

Dec. 5, 2023

(58) Field of Classification Search

CPC H10B 51/30; H10B 51/20; H10B 51/50; H10B 51/10; G11C 11/223;

(Continued)

(56) References Cited

U.S. PATENT DOCUMENTS

4,213,139 A 7/1980 Rao 4,984,153 A 1/1991 Kregness et al. (Continued)

FOREIGN PATENT DOCUMENTS

CN 111799263 A * 10/2020 JP 1998-269789 A2 10/1998 (Continued)

OTHER PUBLICATIONS

"PCT Search Report and Written Opinion, PCT/US2022/039473", dated Dec. 6, 2022, 14 pages.

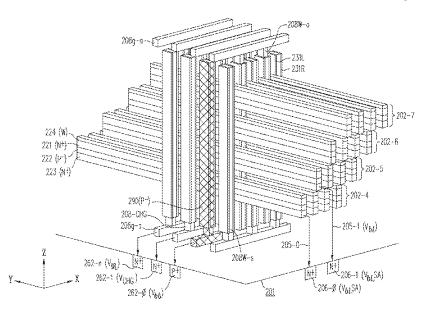
(Continued)

Primary Examiner — Fernando Hidalgo

(57) ABSTRACT

Thin-film Ferroelectric field-effect transistor (FeFET) may be organized as 3-dimensional NOR memory string arrays. Each 3-dimensional NOR memory string array includes a row of active stack each including a predetermined number of active strips each provided one on top of another and each being spaced apart from another by an isolation layer. Each active strip may include a shared source layer and a shared drain layer shared by the FeFETs provided along the active strip. Data storage in the active strip is provided by ferroelectric elements that can individually electrically set into one of two polarization states. FeFETs on separate active strips may be configured for read, programming or erase operations in parallel.

66 Claims, 49 Drawing Sheets



US 11,839,086 B2 Page 2

(50) E		1e 41	6 1	0.210.225	D2	12/2015	Wanda akal
· /		ssification		9,219,225 9,230,985			Karda et al. Wu et al.
CPG	~ ·······		1/2273; G11C 11/2275; G11C /2259; H01L 29/78391; H01L	9,231,206			Tao et al.
		11		9,263,577			Ramaswamy et al.
TIO	20		29/40111; H01L 29/516	9,281,044	B2 *		Ramaswamy H01L 29/78391
				9,299,580			Kong et al.
See	applicati	on file to	r complete search history.	9,337,210 9,362,487			Karda et al. Inumiya et al.
(50)		D. C	C' I	9,391,084		7/2016	
(56)		Referen	ces Cited	9,412,752			Yeh et al.
	II S	DATENT	DOCUMENTS	9,455,268	B2		Oh et al.
	0.5.	17111111	DOCOMENTS	9,472,560			Ramaswamy et al.
5,388	,246 A	2/1995	Kasai	9,530,794 9,558,804			Ramaswamy et al.
	,808 A		Brahmbhatt	9,538,804		1/2017 4/2017	Liang et al.
5,646	,886 A		Brahmbhatt	9,633,944		4/2017	
	,842 A		Iwamatsu	9,698,152	B2	7/2017	Peri
	,192 A ,776 A	6/1998 8/1998	Ancaster et al.	9,711,529			Hu et al.
	,993 A		Kramer et al.	9,748,172	B2	8/2017	Takaki Ramaswamy et al.
	,167 A	6/1999	Leedy	9,786,684 9,799,761			Or-Bach et al.
	,605 A		Sano et al.	9,818,468		11/2017	
	,497 A	4/2000		9,818,848	B2		Sun et al.
	,862 A ,133 A		Margulis Furukawa et al.	9,830,969			Slesazeck et al.
	,171 A		Davies et al.	9,842,651		12/2017	Harari
	,838 A		Kim et al.	9,865,680 9,875,784			Okumura et al. Li et al.
6,313	,518 B1		Ahn et al.	9,876,018			Chavan et al.
	,046 B1		Kamiya et al.	9,892,800		2/2018	
	,508 B1	3/2002 5/2002	Rasovaky et al.	9,911,497		3/2018	
	,744 B1 ,053 B1		Fujiwara	9,941,299			Chen et al.
	,124 B1		Cleeves et al.	10,014,317 10,038,092		7/2018	Chen et al.
	,365 B1		Salling	10,033,557			Slesazeck et al.
	,503 B2	9/2003		10,056,393			Schroeder et al.
	,094 B2 ,105 B1	6/2004 6/2004	Chang et al.	10,074,667			Higashi
	,458 B2		Fricke et al.	10,090,036 10,096,364		10/2018	Van Houdt
6,873	,004 B1		Han et al.	10,030,364			Harari et al.
	,994 B2		Lee et al.	10,157,780			Wu et al.
	,703 B2 ,350 B2		Ryu et al. Walker et al.	10,211,223			Van Houdt et al.
	,977 B2		Chen et al.	10,211,312			Van Houdt et al.
	,653 B2		Cheng et al.	10,217,667 10,249,370		4/2019	Or-Bach et al.
	,308 B2	12/2007		10,254,968			Gazit et al.
	,141 B2		Takeuchi	10,283,493		5/2019	Nishida
	,980 B2 ,002 B2		Arimoto et al. Forbes et al.	10,319,696			Nakano
	,963 B2		Edahiro et al.	10,355,121		2/2019	Or-Bach et al. Gupta et al.
	,012 B2	3/2009	Kuo	10,373,956 10,381,370			Shin et al.
	,725 B2	4/2009		10,381,378		8/2019	
	,348 B1	6/2009 11/2009		10,395,737	B2	8/2019	
	,411 B2 ,359 B2		Boes et al.	10,403,627			Van Houdt et al.
	,145 B2		Shimizu et al.	10,418,377 10,424,379			Van Houdt et al. Slesazeck et al.
	,295 B2		Park et al.	10,424,579			Herner et al.
	,614 B2		Kang et al.	10,438,645			Muell et al.
	,009 B2 ,521 B1		Wilson et al.	10,460,788		10/2019	
	,321 B1 ,418 B2		Or-Bach et al. Carman	10,475,812		11/2019	
	,396 B2		Sinha et al.	10,510,773 10,600,808			Ramaswamy et al. Schröder et al.
8,237	,213 B2	8/2012	Liu	10,608,008			Harari et al.
	,504 B2	8/2012		10,608,011			Harari et al.
	,183 B2 ,823 B2	10/2012	Lerner Boescke	10,622,051			Mueller et al.
	,823 B2 ,482 B2		Kim et al.	10,622,377 10,636,471			Harari et al.
8,395	,942 B2		Samachisa et al.	10,650,892		5/2020	Ramaswamy et al.
	,618 B2		Cooney et al.	10,651,153			Fastow et al.
	,114 B2 ,612 B2	1/2014	Lue Choi et al.	10,651,182	B2	5/2020	Morris et al.
	,612 B2 ,436 B2		Scalia et al.	10,651,196			Sharangpani et al.
	,473 B2		Shim et al.	10,692,837			Nguyen et al.
8,848	,425 B2	9/2014	Schloss	10,692,874 10,700,093			Harari et al. Kalitsov et al.
	,278 B2		Alsmeier et al.	10,700,093		7/2020	
	,801 B2 ,802 B2		Sandhu et al. Mueller	10,720,437		11/2020	
	,802 B2 ,310 B2		Li et al.	10,872,905	B2	12/2020	Muel
9,158	,622 B2	10/2015	Lee et al.	10,879,269			Zhang et al.
	,293 B2		Wang et al.	10,896,711			Lee et al.
9,202	,694 B2	12/2015	Konevecki et al.	10,937,482	B2	5/2021	Sharma et al.

US 11,839,086 B2 Page 3

(56)	Referen	ces Cited	2013/0031325			Nakamoto et al.	
	TTO 1	DATENT	DOCUMENTS	2013/0126957 2013/0256780			Higashitani et al.
	U.S.	PATENT	DOCUMENTS	2014/0015036			Kai et al. Fursin et al.
10,950,616	B2	3/2021	Harari et al.	2014/0040698			Loh et al.
11,043,280			Prakash et al.	2014/0070289	A1	3/2014	Tanaka et al.
11,171,157			Lai et al.	2014/0070290			Inumiya et al.
11,411,025			Lai et al.	2014/0075135			Choi et al. Dunga et al.
2001/0030340			Fujiwara	2014/0112075 2014/0117366		5/2014	
2001/0053092 2002/0012271		1/2001	Kosaka et al.	2014/0151774		6/2014	
2002/0012271			Lee et al.	2014/0173017			Takagi et al.
2002/0051378			Ohsawa	2014/0213032			Kai et al.
2002/0193484		12/2002		2014/0229131			Cohen et al. Karda et al.
2003/0038318		2/2003		2014/0247674 2014/0328128			Louie et al.
2004/0214387 2004/0246807		10/2004	Madurawe et al.	2014/0340952			Ramaswamy et al.
2004/0262681			Masuoka et al.	2014/0355328			Muller et al.
2004/0262772	A1	12/2004	Ramanathan et al.	2015/0054507			Gulaka et al.
2004/0264247		12/2004		2015/0098272 2015/0113214			Kasolra et al. Sutardja
2005/0128815			Ishikawa et al.	2015/0155876			Jayasena et al.
2005/0236625 2005/0280061		10/2005	Schuele et al.	2015/0194440			Noh et al.
2006/0001083			Bhattacharyya	2015/0220463			Fluman et al.
2006/0080457			Hiramatsu et al.	2015/0249143		9/2015	
2006/0155921			Gorobets et al.	2015/0263005 2015/0340371		9/2015	Zhao et al.
2006/0212651		9/2006 11/2006	Ashmore	2015/0340371			Chen et al.
2006/0261404 2007/0012987			McTeer	2016/0013156			Zhai et al.
2007/0023817		2/2007		2016/0019951	Al	1/2016	Park et al.
2007/0045711			Bhattacharyya	2016/0035711		2/2016	
2007/0134876			Lai et al.	2016/0086970		3/2016	
2007/0192518			Rupanagunta et al.	2016/0118404 2016/0141294		4/2016 5/2016	Peri et al.
2008/0022026 2008/0054346			Yang et al. Saitoh et al.	2016/0225860			Karda et al.
2008/0034340			Lee et al.	2016/0276360			Doda et al.
2008/0173930			Watanabe	2016/0300724			Levy et al.
2008/0178794			Cho et al.	2016/0314042		10/2016	
2008/0212358	A1*	9/2008	Mitsui H01L 21/76849	2016/0321002 2016/0358934			Jung et al. Lin et al.
2009/0220912	A 1	10/2009	Maafumi et al	2017/0053906			Or-Bach et al.
2008/0239812 2008/0266960		10/2008	Naofumi et al.	2017/0062456			Sugino et al.
2008/0291723			Wang et al.	2017/0092370		3/2017	
2008/0301359		12/2008		2017/0092371		3/2017	
2009/0057722			Masuoka et al.	2017/0140807 2017/0148517		5/2017	Sun et al.
2009/0140318		6/2009	e e e e e e e e e e e e e e e e e e e	2017/0148810			Kai et al.
2009/0157946 2009/0237996		6/2009 9/2009	Kirsch et al.	2017/0194341	A1		Yamada
2009/0268519		10/2009		2017/0213731			Yoon et al.
2009/0279360	A1	11/2009	Peter et al.	2017/0213821 2017/0358594			Or-Bach et al. Lu et al.
2009/0290442		11/2009		2018/0006044			Chavan et al.
2009/0316487			Lee et al.	2018/0095127			Pappu et al.
2010/0013001 2010/0121994			Cho et al. Kim et al.	2018/0108416		4/2018	Harari
2010/0124116			Takashi et al.	2018/0151419	A1		Wu et al.
2010/0128509	A1		Kim et al.	2018/0261613			Ariyoshi et al.
2010/0148215			Schulze et al.	2018/0269229 2018/0286918			Or-Bach et al. Bandyopadhyay et al.
2010/0207185			Lee et al. Awaya et al.	2018/0330791			Li et al.
2010/0219392 2010/0254191			Son et al.	2018/0331042			Manusharow et al.
2010/0327413			Lee et al.	2018/0342455			Nosho et al.
2011/0003418	A1		Sakata et al.	2018/0342544 2018/0366471			Lee et al. Harari et al.
2011/0044113		2/2011		2018/0366485		12/2018	
2011/0047325 2011/0115011			Mishima Masuoka et al.	2018/0366489			Harari et al.
2011/0113011			Jones et al.	2018/0374929	A1	12/2018	Yoo
2011/0143519		6/2011		2019/0006009		1/2019	
2011/0170266			Haensh et al.	2019/0006015 2019/0019564			Norman et al. Li et al.
2011/0208905			Shaeffer et al.	2019/0019304			Herner et al.
2011/0298013 2011/0310683			Hwang et al. Gorobets	2019/0007327			Or-Bach et al.
2011/0310083		3/2012		2019/0157296			Harari et al.
2012/0146126			Lai et al.	2019/0180821	A1	6/2019	Harari
2012/0182801	A1	7/2012	Lue	2019/0206890			Harari et al.
2012/0208347			Hiwang et al.	2019/0214077			Oh et al.
2012/0223380			Lee et al.	2019/0237470			Mine et al.
2012/0243314 2012/0307568			Takashi Banna et al.	2019/0238134 2019/0244971		8/2019 8/2019	Lee et al.
2012/0307308		12/2012		2019/0259769			Karda et al.
2012/052//17		12,2012		2017, 0207 107		5, 2017	

(56)References Cited U.S. PATENT DOCUMENTS 2019/0303042 A1 10/2019 Kim et al 2019/0304988 A1 10/2019 Dong et al. pp. 235-248. 2019/0325945 A1 10/2019 Linus 2019/0325964 A1 10/2019 Harari 2019/0319044 A1 11/2019 Harari 2019/0348424 A1 11/2019 Karda et al. 2019/0355747 A1 11/2019 Herner et al. 2019/0370117 A1 12/2019 Fruchtman et al. 2020/0020718 A1 1/2020 Harari et al. 2020/0051990 A1 2/2020 Harari et al. 2020/0063263 A1 2/2020 Technology, Jun. 2021, 2 p. Yang et al. 2020/0065647 A1 2/2020 Mulaosmanovic et al. 2020/0075631 A1 3/2020 Dong 2020/0098738 A1 3/2020 Herner et al. 2020/0098779 A1 3/2020 Cernea et al. 3060589., Apr. 2021, pp. 617-620. 2020/0098881 A1* 3/2020 Vellianitis H01L 29/66628 2020/0176468 A1 6/2020 Herner et al. 2020/0201718 A1 6/2020 Richter et al. 2020/0203378 A1 6/2020 Harari et al. 2020/0219572 A1 7/2020 Harari Papers, 2010, pp. 131-132. 2020/0243486 A1 7/2020 Ouader et al. 2020/0258897 A1 8/2020 Yan et al. Hoffmann H10B 53/30 2020/0350324 A1* 11/2020 2020/0357453 A1 11/2020 Slesazeck et al. 2020/0357455 A1 11/2020 Noack et al. 2020/0357470 A1 11/2020 Noack 2020/0357822 A1* 11/2020 Chen G11C 11/223 2020/0365609 A1 11/2020 Harari 2020/0388313 A1 12/2020 Cho et al. 2020/0388711 A1 12/2020 Doyle et al. 2020/0403002 A1 12/2020 Harari 2020/0411533 A1 12/2020 Alsmeier et al. 2021/0005238 A1 Mueller 1/2021 pages. 2021/0013224 A1 1/2021 Purayath et al. 2021/0020659 A1 1/2021 Chen 2021/0066502 A1 3/2021 Karda et al. 2021/0074725 A1 3/2021 Lue 2021/0074726 A1 3/2021 Lue 2021/0111179 A1 4/2021 Shivaraman et al. 2021/0175251 A1 6/2021 Zhang et al. pages 2021/0247910 A1 8/2021 Kim et al. 2021/0248094 A1 8/2021 Norman et al. 2021/0265308 A1 8/2021 Norman et al. 2021/0375933 A1* 12/2021 Lu H01L 29/40111 (2015), Feb. 21, 2015, pp. N29-N35. 2022/0028876 A1 1/2022 Purayath et al. 2022/0028886 A1 1/2022 Pur et al. 2022/0139933 A1* 5/2022 Noack G11C 11/221 2022/0231049 A1* $7/2022 \quad Lin \quad \qquad H10B \; 51/20$ 10.8.1-10.8.4. 2022/0246777 A1* 8/2022 Kim B32B 17/10 2022/0254390 A1* 8/2022 Gans G11C 5/04 2022/0351776 A1* 11/2022 Nam G11C 11/223 2022/0384459 A1 12/2022 Lu et al. 2022/0393031 A1* Ando H01L 29/66484 12/2022 4199-4205. 2023/0052477 A1* 2/2023 Ha H01L 29/0847 FOREIGN PATENT DOCUMENTS 2000-243972 A2 9/2000 JP JP 2000339978 A 12/2000 2004079606 A 3/2004 JP JP 2006099827 A1 4/2006 2009206451 A1 9/2009 JР 2010108522 A1 5/2010 JΡ 2010251572 A 11/2010 JР 2011028540 A1 2/2011

JΡ

KR

KR

KR

WO

WO

WO

2012-150876 A2

20120085591 A

20120085591 A1

2015025357 A1

2018236937 A1

2019066948 A1

20120085603 A

8/2012

8/2012

8/2012

8/2012

2/2015

12/2018

4/2019

OTHER PUBLICATIONS

Lee, Sang Yeol, "Comprehensive Review on Amorphous Oxide Semiconductor Thin Film Transistor", Trans. Electr. Electron. Mater. 21, https://doi.org/10.1007/s42341-020-00197-w, Mar. 28, 2020,

Lee, Chea-Young, et al., "Etching Characteristics and Changes in Surface Properties of IGZO Thin Films by O2 Addition in CF4/Ar Plasma", Coatings. 2021, 11(8):906. https://doi.org/10.3390/ coatings11080906, Jul. 29, 2021.

Liao, P.J., et al., "Characterization of Fatigue and Its Recovery Behavior in Ferroelectric HfZrO", 2021 Symposium on VLSI

Liao, C.Y., et al., "Multibit Ferroelectric FET Based on Nonidentical Double HfZrO2 for High-Density Nonvolatile Memory", IEEE Electron Device Letters, vol. 42, No. 4, doi: 10.1109/LED.2021.

Lue, H.T., et al., "A Highly Scalable 8-Layer 3D Vertical-gate $\{VG)$ TFT NANO Flash Using Junction-Free Buried Channel BE-SONOS Device", Symposium on VLSI: Tech. Dig. Of Technical

Materano, Monica, et al., "Interplay between oxygen defects and dopants: effect on structure and performance of HfO2-based ferroelectrics", Inorg. Chem. Front. (8), Apr. 9, 2021, pp. 2650-2672. Mo, Fei, et al., "Experimental Demonstration of Ferroelectric HfO2 FET with Ultrathin-body IGZO for High-Density and Low-Power Memory Application", 2019 Symposium on VLSI Technology, 3, doi: 10.23919/VLSIT.2019.8776553., 2019, pp. T42-T4.

Mo, Fei, et al., "Experimental Demonstration of Ferroelectric HfO2 FET with Ultrathin-body IGZO for High-Density and Low-Power Memory Application", Symposium on VLSI Technology, 2019, pp. T42-T43, doi: 10.23919/VLSIT.2019.8776553., Jul. 13, 2020, 7

Mo, Fei, et al., "Low-Voltage Operating Ferroelectric FET with Ultrathin IGZO Channel for High-Density Memory Application", IEEE Journal of the Electron Devices Society, vol. 8, doi: 10.1109/ JEDS.2020.3008789., Jul. 13, 2020, pp. 717-723.

Mueller, S., et al., "Development Status of Gate-First FeFET Technology", 2021 Symposium on VLSI Technology, Jun. 2021, 2

Mueller, J., et al., "Ferroelectric Hafnium Oxide Based Materials and Devices: Assessment of Current Status and Future Prospects", ECS Journal of Solid State Science and Technology, 4 (5) N30-N35

Mueller, J., et al., "Ferroelectric hafnium oxide: A CMOScompatible and highly scalable approach to future ferroelectric memories", IEEE International Electron Devices Meeting, Washington, DC, USA, doi: 10.1109/IEDM.2013.6724605., 2013, pp.

Mueller, Stephen, et al., "From MFM Capacitors Toward Ferroelectric Transistors: Endurance and Disturb Characteristics of HfO2-Based FeFET Devices", IEEE Transactions on Electron Devices, vol. 60, No. 12, doi: 10.1109/TED.2013.2283465., Dec. 2013, pp.

Müller, J., et al., "Ferroelectric hafnium oxide: A CMOS-compatible and highly scalable approach to future ferroelectric memories", 2013 IEEE International Electron Devices Meeting, Washington, DC, USA, doi: 10.1109/IEDM.2013.6724605., 2013, pp. 10.8.1-10.

Nguyen, Manh-Cuong, et al., "Wakeup-Free and Endurance-Robust Ferroelectric Field-Effect Transistor Memory Using High Pressure Annealing", IEEE Electron Device Letters, vol. 42, No. 9, doi: 10.1109/LED.2021.3096248., Sep. 2021, pp. 1295-1298.

On, Nuri, et al., "Boosting carrier mobility and stability in indiumzinc-tin oxide thin-film transistors through controlled crystallization", Sci Rep 10, 18868, https://doi.org/10.1038/s41598-020-76046w, 2020, 16 pages.

Onuki, Tatsuya, et al., "Embedded memory and ARM Cortex-MO core using 60-nm C-axis aligned crystalline indium-gallium-zinc oxide FET integrated with 65-nm Si CMOS", IEEE Journal of Solid-State Circuits, vol. 52, No. 4,, Apr. 2017, pp. 925-932.

(56) References Cited

OTHER PUBLICATIONS

Park, Goon-Ho, et al., "Electrical Characteristics of SiO2/High-k Dielectric Stacked Tunnel Barriers for Nonvolatile Memory Applications", Journal of the Korean Physical Society, vol. 55, No. 1, Jul. 2009, pp. 116-119.

Park, Min Hyuk, et al., "Review of defect chemistry in fluorite-structure ferroelectrics for future electronic devices", J. Mater. Chem. C, vol. 8, No. 31 10526-10550, Jun. 9, 2020., Aug. 21, 2020, pp. 10526-10550.

Rios, Rafael, et al., "A Physically Based Compact Model for IGZO Transistors", IEEE Transactions on Electron Devices, vol. 68, No. 4, doi: 10.1109/TED.2021.3059387., Apr. 2021, pp. 1664-1669. Rzehak, Volker, "Low-Power FRAM Microcontrollers and Their Applications", Texas Instruments White Paper, SLAA502, Jul.

2019, 7 pages.

Saitoh, Masumi, et al., "HfO2-based FeFET and FTJ for Ferroelectric-Memory Centric 3D LSI towards Low-Power and High-Density Storage and AI Applications", 2020 IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553.2020. 9372106., 2020, pp. 18.1.1-18.1.4.

Saitoh, Masumi, et al., "HfO2-based FeFET and FTJ for Ferroelectric-Memory Centric 3D LSI towards Low-Power and High-Density Storage and AI Applications", IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553.2020.9372106., 2020, pp. 18.1.1-18.1.4.

Salahuddin, Sayeef, et al., "FeFETs for Near-Memory and In-Memory Compute", 2021 IEEE International Electron Devices Meeting (IEDM), Department of Electrical Engineering and Computer Sciences, Dec. 2021, 4 pages.

Sharma, Abhishek A., et al., "High Speed Memory Operation in Channel-Last, Back-gated Ferroelectric Transistors", IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553. 2020.9371940., 2020, pp. 18.5.1-18.5.4.

Sheng, Jiazhen, et al., "Review Article: Atomic layer deposition for oxide semiconductor thin film transistors: Advances in research and development", J. Vac. Sci. Technol. A 36, 060801, https://doi.org/10.1116/1.5047237., Nov. 2, 2018, 14 pages.

Si, Mengwei, et al., "Why In2O3 Can Make 0.7 nm Atomic Layer Thin Transistors?", https://arxiv.org/ftp/arxiv/papers/2012/2012. 12433.pdf., 2012, 26 pages.

Si, Mengwei, et al., "Why In2O3 Can Make 0.7 nm Atomic Layer Thin Transistors?", available at https://arxiv.org/ftp/arxiv/papers/2012/2012.12433.pdf, School of Electrical and Computer Engineering and Birck Nanotechnology Center, Purdue University, West Lafayette, IN 47907, United States, 2012, pp. 1-26.

Su, Nai-Chao, et al., "A Nonvolatile InGaZnO Charge-Trapping-Engineered Flash Memory With Good Retention Characteristics", IEEE Electron Device Letters, vol. 31, No. 3, Mar. 2010, pp. 201-203.

Sun, Chen, "First Demonstration of BEOL-Compatible Ferroelectric TCAM Featuring a-IGZO Fe-TFTs with Large Memory Window of 2.9 V, Scaled Channel Length of 40 nm, and High Endurance of 108 Cycles", 2021 Symposium on VLSI Technology Digest of Technical Papers, Kyoto, doi: 978-4-86348-779-6., Jun. 2021, 2 pages.

Tan, Ava J., et al., "A Nitrided Interfacial Oxide for Interface State Improvement in Hafnium Zirconium Oxide-Based Ferroelectric Transistor Technology", IEEE Electron Device Letters, vol. 39, No. 1, doi: 10.1109/LED.2017.2772791., Jan. 2018, pp. 95-98.

Tan, Ava J., et al., "Experimental Demonstration of a Ferroelectric HfO2-Based Content Addressable Memory Cell", IEEE Electron Device Letters, vol. 41, No. 2, doi: 10.1109/LED.2019.2963300., Feb. 2020, pp. 240-243.

Tan, Ava Jiang, et al., "Ferroelectric HfO2 Memory Transistors with High-κ Interfacial Layer and Write Endurance Exceeding 1010 Cycles", arXiv:2103.08806 [physics.app-ph], available at https://arxiv.org/abs/2103.08806., Mar. 16, 2021.

Tan, Ava J., et al., "Hot Electrons as the Dominant Source of Degradation for Sub-5nm HZO FeFETs", IEEE Symposium on

VLSI Technology, Honolulu, HI, USA, doi: 10.1109/VLSITechnology18217.2020.9265067., 2020, pp. 1-2.

Tan, Yan-Ny, et al., "Over-Erase Phenomenon in SONOS-Type Flash Memory and its Minimization Using a Hafnium Oxide Charge Storage Layer", IEEE Transactions on Electron Devices, vol. 51, No. 7, Jul. 2004, pp. 1143-1147.

Tanaka, T., et al., "A 768 GB 3b/cell 3D-Floaling-Gate NANO

Tanaka, T., et al., "A 768 GB 3b/cell 3D-Floaling-Gate NANO Flash Memory", Digest of Technical Papers, the 2016 EEE International Solid-Slate Circuits Conference, 2016, pp. 142-144.

Van Houdt, Jan, "The 3D FeFET: contender for 3D-NAND Flash memory and machine learning", available at https://www.imec-int.com/en/imec-magazine/imec-magazine-october-2017/the-vertical-ferroelectric-fet-a-new-contender-for-3d-nand-flash-memory-and-machine-learning, Sep. 30, 2019, 8 pages.

Wann, H.C., et al., "High-Endurance Ultra-Thin Tunnel Oxide in Monos Device Structure for Dynamic Memory Application", IEEE Electron Device letters, vol. 16, No. 11, Nov. 1995, pp. 491-493. Yang, Jin, et al., "Highly Optimized Complementary Inverters Based on p-SnO and n-InGaZnO With High Uniformity", IEEE Electron Device Letters, vol. 39, No. 4, doi: 10.1109/LED.2018. 2809796., Apr. 2018, pp. 516-519.

Wu, Jixuan, et al., "A Monolithic 3D Integration of RRAM Array with Oxide Semiconductor FET for In-Memory Computing in Quantized Neural Network AI Applications", 2020 IEEE Symposium on VLSI Technology Digest of Technical Papers, Honolulu, HI, USA, Jun. 2020, 4 pages.

"EP Extended Search Report EP168690149.3", dated Oct. 18, 2019. "European Search Report, EP 16852238.1", dated Mar. 28, 2019. "European Search Report, EP17844550.8", dated Aug. 12, 2020, 11 pages.

"Imec Demonstrates Capacitor-less IGZO-Based DRAM Cell With >400s Retention Time", IMEC, Press release, available at https://www.imec-int.com/en/press/imec-demonstrates-capacitor-less-igzo-based-dram-cell-400s-retention-time, Dec. 15, 2020, 15 pages.

"Invitation to Pay Additional Fees (PCT/ISA/206), PCT/US2020/015710", dated Mar. 20, 2020, 2 pages.

"Invitation to Pay Additional Fees, PCT/US2019/065256", dated Feb. 13, 2020, 2 pages.

"Notification of Reasons for Refusal, Japanese Patent Application 2018-527740", (English translation), dated Nov. 4, 2020, 8 pages. "Partial European Search Report EP 16869049.3", dated Jul. 1, 2019, pp. 1-12.

"PCT Search Report and Written Opinion, PCT/US2018/038373", dated Sep. 10, 2018.

"PCT Search Report and Written Opinion, PCT/US2018/067338", dated May 8, 2019.

"PCT Search Report and Written Opinion, PCT/US2019/014319", dated Apr. 15, 2019.

"PCT Search Report and Written Opinion, PCT/US2019/041678", dated Oct. 9, 2019.

"PCT Search Report and Written Opinion, PCT/US2019/052164", dated Feb. 27, 2020.

"PCT Search Report and Written Opinion, PCT/US2019/052446", dated Dec. 11, 2019.

"PCT Search Report and Written Opinion, PCT/US2019/065256", dated Apr. 14, 2020.

"PCT Search Report and Written Opinion, PCT/US2020/015710", dated Jun. 9, 2020.

"PCT Search Report and Written Opinion, PCT/US2020/017494", dated Jul. 20, 2020, 13 pages.

"PCT Search Report and Written Opinion, PCT/US2020/065374", dated Mar. 15, 2021, 17 pages.

"PCT Search Report and Written Opinion, PCT/US2020/065670", dated Apr. 5, 2021, 12 pages.

"PCT Search Report and Written Opinion, PCT/US2021/016964", dated Jun. 15, 2021, 19 pages.

"PCT Search Report and Written Opinion, PCT/US2021/025722", dated Jun. 15, 2021, 10 pages.

"PCT Search Report and Written Opinion, PCT/US2021/042607", dated Nov. 4, 2021, 17 pages.

"PCT Search Report and Written Opinion, PCT/US2021/064844", dated Mar. 8, 2022, 15 paged.

(56) References Cited

OTHER PUBLICATIONS

"PCT Search Report and Written Opinion, PCT/US2021/42620", dated Oct. 28, 2021, 18 pages.

"PCT Search Report and Written Opinion, PCT/US2022/016729", Applicant: SunRise Memory Corporation, dated May 17, 2022, 20 pages.

Ahn, Min-Ju, et al., "Transparent multi-level-cell nonvolatile memory with dual-gate amorphous indiumgallium-zinc oxide thin-film transistors", Appl. Phys. Lett. 109, 252106; doi: 10.1063/1. 4972961., 2016, 6 pages.

Alessandri, Cristobal, et al., "Monte Carlo Simulation of Switching Dynamics in Polycrystalline Ferroelectric Capacitors", IEEE Transactions on Electron Devices, vol. 66, No. 8, doi: 10.1109/TED. 2019.2922268., Aug. 2019, pp. 3527-3534.

Ali, T., et al., "A Multilevel FeFET Memory Device based on Laminated HSO and HZO Ferroelectric Layers for High-Density Storage", IEEE International Electron Devices Meeting (IEDM), 2019, doi: 10.1109/ EDM19573.2019.8993642., Dec. 2019, pp. 28.7.1-28.7.4.

Ali, T., et al., "High Endurance Ferroelectric Hafnium Oxide-Based FeFET Memory Without Retention Penalty,", IEEE Transactions on Electron Devices, vol. 65, No. 9, doi: 10.1109/TED.2018.2856818.—paper cited in Bae (Berkeley) paper, Sep. 2018, pp. 3769-3774. Bae, Jong-Ho, et al., "Highly Scaled, High Endurance, Q-Gate,

Nanowire Ferroelectric FET Memory Transistors", IEEE Electron Device Letters, vol. 41, No. 11, doi: 10.1109/LED.2020.3028339.— Sayeef-Berkeley paper on FeFET memory, Nov. 2020, pp. 1637-1640.

Beyer, Sven, et al., "FeFET: A versatile CMOS compatible device with game-changing potential", IEEE International Memory Workshop (IMW), doi: 10.1109/IMW48823.2020.9108150., 2020, pp. 1-4.

Böscke, T.S., et al., "Ferroelectricity in hafnium oxide: CMOS compatible ferroelectric field effect transistors", 2011 International Electron Devices Meeting, Washington, DC, USA, doi: 10.1109/IEDM.2011.6131606., 2011, pp. 24.5.1-24.5.4.

Chang, Sou-Chi, et al., "Anti-ferroelectric HfxZr1-xO2 Capacitors for High-density 3-D Embedded-DRAM", IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553.2020. 9372011., 2020, pp. 28.1.1-28.1.4.

Cho, Min Hoe, et al., "Achieving a Low-Voltage, High-Mobility IGZO Transistor through an ALD-Derived Bilayer Channel and a Hafnia-Based Gate Dielectric Stack", ACS Applied Materials & Interfaces, Apr. 1, 2021, 13 (14), DOI: 10.1021/acsami.0c22677, 2021, pp. 16628-16640.

Cho, Min Hoe, et al., "Comparative Study on Performance of IGZO Transistors With Sputtered and Atomic Layer Deposited Channel Layer", IEEE Transactions on Electron Devices, vol. 66, No. 4, doi: 10.1109/TED.2019.2899586., Apr. 2019, pp. 1783-1788.

Choi, Seonjun, et al., "A novel three-dimensional NAND flash structure for improving the erase performance", IEICE Electronics Express, 2019 vol. 16 Issue 3, 2019, 6 pages.

De, Sourav, et al., "Ultra-Low Power Robust 3bit/cell Hf0.5Zr0. 5O2 Ferroelectric FinFET with High Endurance for Advanced Computing-In-Memory Technology", 2021 Symposium on VLSI Technology, Kyoto, Japan, Jun. 13-19, 2021, Jun. 2021, 2 pages. Dünkel, "A FeFET based super-low-power ultra-fast embedded NVM technology for 22nm FDSOI and beyond", IEEE International Electron Devices Meeting (IEDM), San Francisco, CA, USA, 2017, doi: 10.1109/IEDM.2017.8268425., 2017, pp. 19.7.1-19.7.4. Dutta, Sourav, et al., "Logic Compatible High-Performance Ferroelectric Transistor Memory", available at https://arxiv.org/abs/2105.11078, Department of Electrical Engineering, University of Notre Dame, Notre Dame, IN 46556, USA, May 24, 2021, 28 pages. Florent, K., "First demonstration of vertically stacked ferroelectric AI doped HfO2 devices for NAND applications", Symposium on VLSI Technology, doi: 10.23919/VLSIT.2017.7998162., 2017, pp.

Florent, K., et al., "Vertical Ferroelectric HfO2 FET based on 3-D NAND Architecture: Towards Dense Low-Power Memory", IEEE

International Electron Devices Meeting (IEDM), San Francisco, CA, 2018, doi: 10.1109/IEDM.2018.8614710., 2018, pp. 2.5.1-2.5.

Hendy, Ian , "The Role of Increased Semiconductor Mobility—an IGZO Case Study", Display Daily, available at https://www.displaydaily.com/article/display-daily/the-role-of-increased-semiconductor-mobility-an-igzo-case-study, Apr. 2021, 11 pages. Hou, S. Y., et al., "Wafer-Leval Integration of an Advanced Logic-Memory System Through the Second-Generation CoWoS Technology", IEEE Transactions on Electron Devices, vol. 64, No. 10, Oct. 2017, 4071-4077.

Hsain, Hanan Alexandra, et al., "Many routes to ferroelectric HfO2: A review of current deposition methods", J. Vac. Sci. Technol. A 40, 010803 (2022), published Dec. 1, 2021., Dec. 1, 2021, pp. 010803-1-010803-36.

Jeewandara, Thamarasee , et al., "High-speed 3-D memory with ferroelectric NAND flash memory", available at https://techxplore.com/news/2021-01-high-speed-d-memoryferroelectric-nand.html, Jan. 26, 2021, 5 pages.

Kawai, H., et al., "Design Principle of Channel Material for Oxide-Semiconductor Field-Effect Transistor with High Thermal Stability and High On-current by Fluorine Doping", IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553. 2020.9372121., 2020, pp. 22.2.1-22.2.4.

Kim, Min-Kyu, et al., "CMOS-compatible ferroelectric NAND flash memory for high-density, low-power, and high-speed three-dimensional memory", Science Advances, vol. 7, No. 3, eabe 1341, DOI: 10.1126/sciadv.abe1341. (Replacement of polysilicon channel material—use IZO for channel), Jan. 13, 2021, 10.

Kim, Taeho, et al., "Effects of high pressure nitrogen annealing on ferroelectric Hf0.5Zr0.5O2 films", Appl. Phys. Lett. 112, 092906 https://doi.org/10.1063/1.5003369, Mar. 2, 2018, 5 pages.

Kim, N., et al., "Multi-layered Vertical gate NANO Flash Overcoming Stacking Limit for Terabit Density Storage", Symposium on VLSI Tech. Dig. of Technical Papers, 2009, pp. 188-189.

Kunitake, Hitoshi, et al., "High-temperature Electrical Characteristics of 60nm CAAC-IGZO FET: Comparison with Si FET", International Conference on Solid State Devices and Materials, https://doi.org/10.7567/SSDM.2018.N-3-04, Sep. 2018, pp. 787-788

Chen, Chi-Yu, et al., "FeFET Memory Featuring Large Memory Window and Robust Endurance of Long-Pulse Cycling by Interface Engineering using High-k AION", 2020 IEEE Symposium on VLSI Technology, doi: 10.1109/VLSITechnology18217.2020.9265103., 2020, pp. 1-2.

Kim, Hyungwoo, et al., "Two-step deposition of TiN capping electrodes to prevent degradation of ferroelectric properties in an in-situ crystallized TiN/Hf0.5Zr0.5O2/TIN device", Nano Express 3 015004., 2022, 12 pages.

McBriarty, Martin E., et al., "Crystal Phase Distribution and Ferroelectricity in Ultrathin HfO2—ZrO2 Bilayers", Phys. Status Solidi B, 257: 1900285. https://doi.org/10.1002/pssb.201900285, Aug. 21, 2019, pp. 1-25.

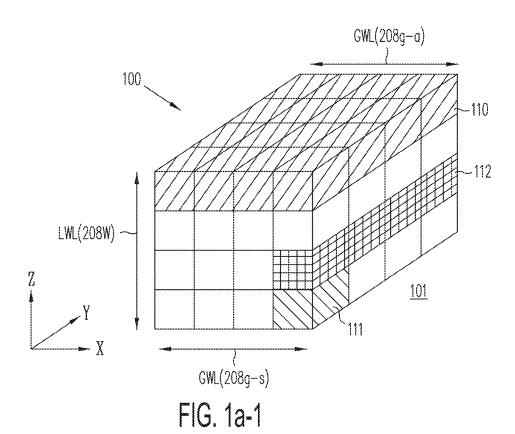
Mittmann, T., et al., "Impact of Oxygen Vacancy Content in Ferroelectric HZO films on the Device Performance", 2020 IEEE International Electron Devices Meeting (IEDM), doi: 10.1109/IEDM13553.2020.9372097., 2020, pp. 18.4.1-18.4.4.

Onaya, Takashi, et al., "Improvement in ferroelectricity of HfxZr1—xO2 thin films using top- and bottom-ZrO2 nucleation layers", APL Materials 7, 061107; https://doi.org/10.1063/1.5096626, 2019, pp. 1.8

Sato, Yuta, et al., "Source/Drain Contact Engineering of InGaZnO Channel BEOL Transistor for Low Contact Resistance and Suppressing Channel Shortening Effect,", 20th International Workshop on Junction Technology (IWJT), doi: 10.23919/IWJT52818.2021. 9609366., 2021, 3 pages.

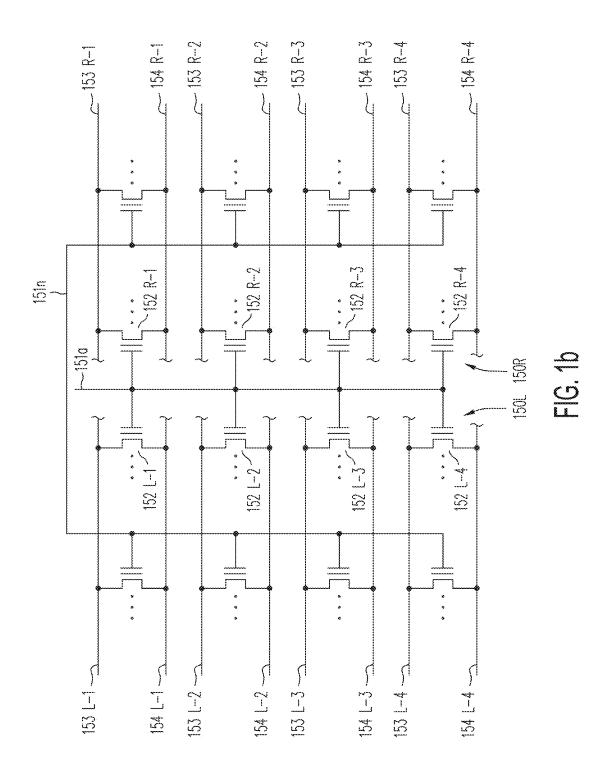
Wang, Chin-I, et al., "Atomic layer deposited TIN capping layer for sub-10 nm ferroelectric Hf0.5Zr0.5O2 with large remnant polarization and low thermal budget", Applied Surface Science, vol. 570, 2021, 151152, ISSN 0169-4332, https://doi.org/10.1016/j.apsusc. 2021.151152., Aug. 9, 2021, pp. 1-8.

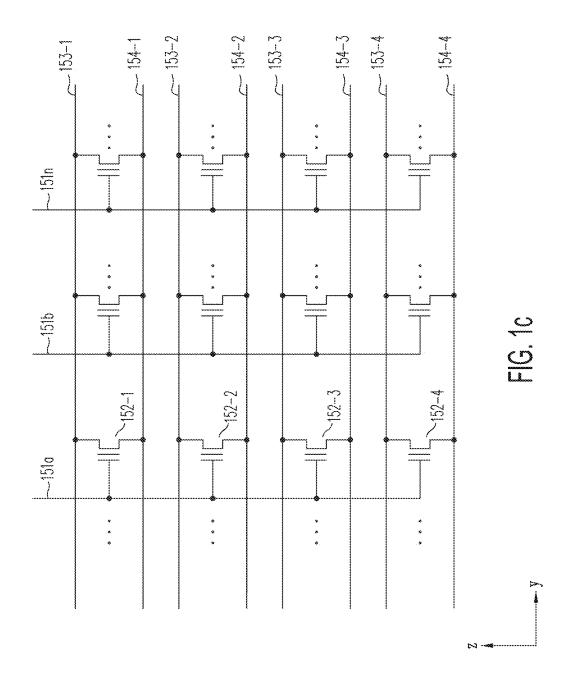
* cited by examiner

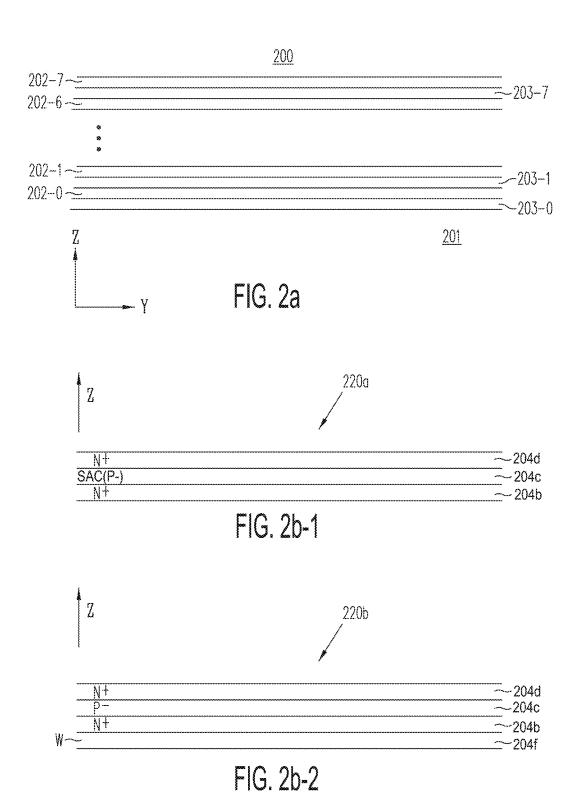


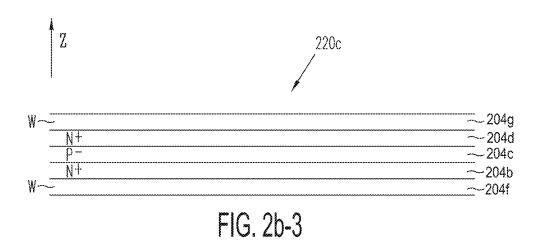
114 Y 100 Y 101

FIG. 1a-2









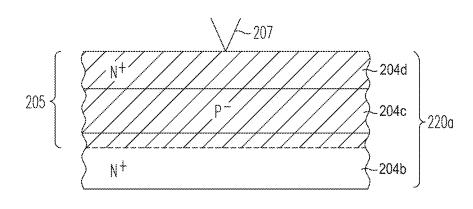


FIG. 2b-4

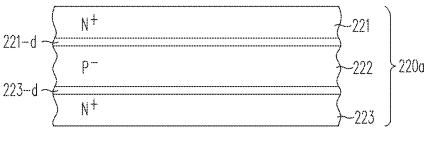
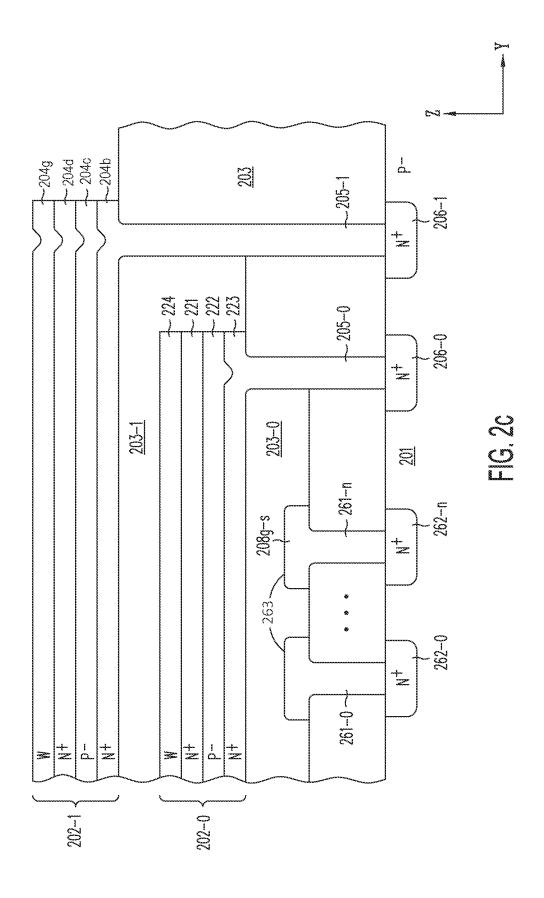
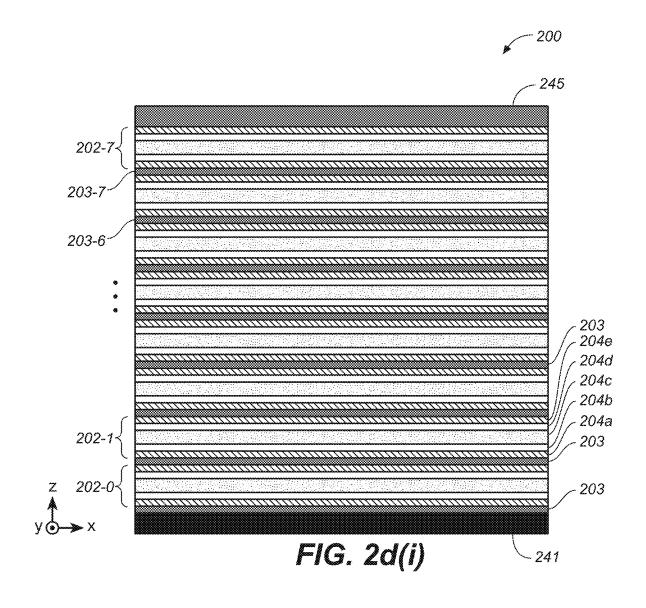
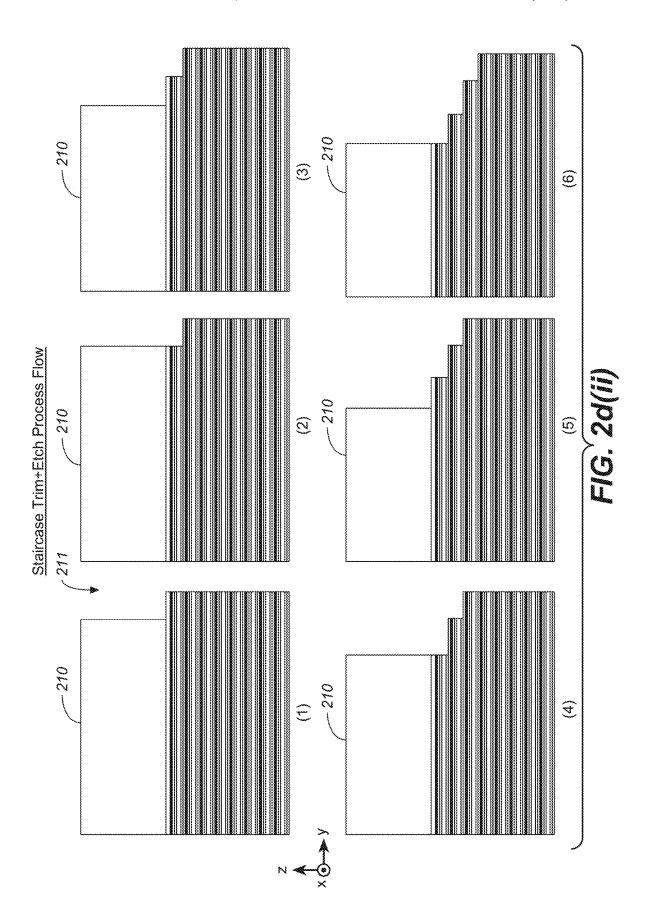
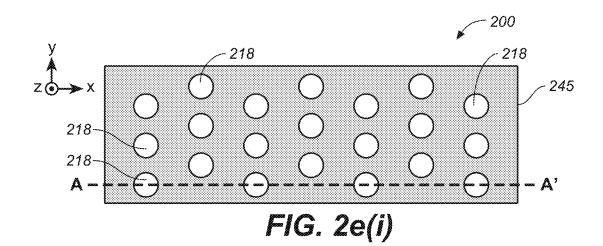


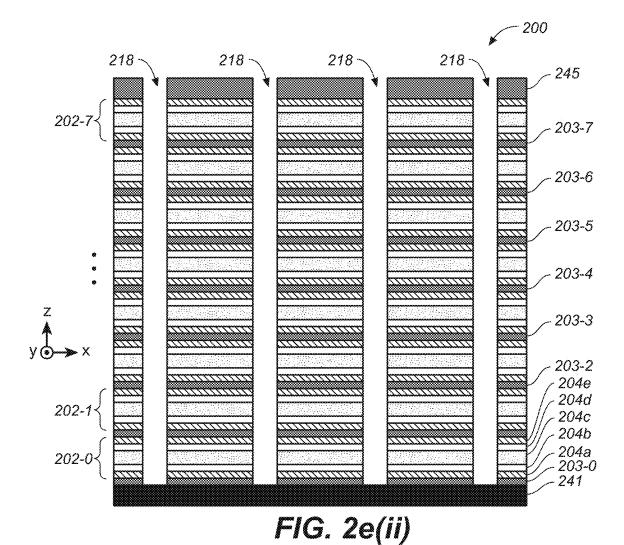
FIG. 2b-5

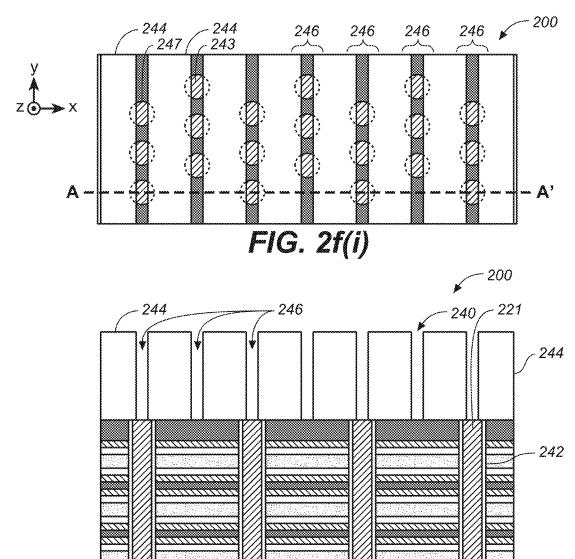


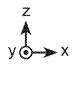












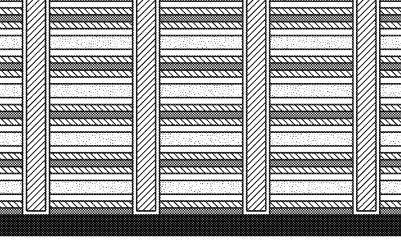
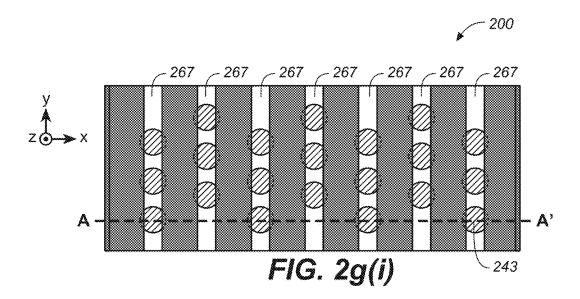


FIG. 2f(ii)



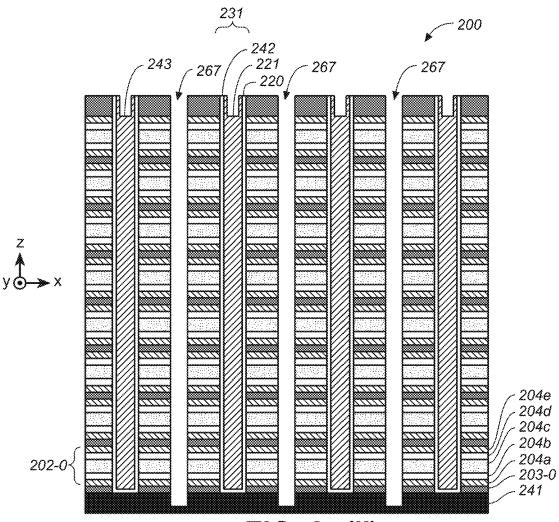
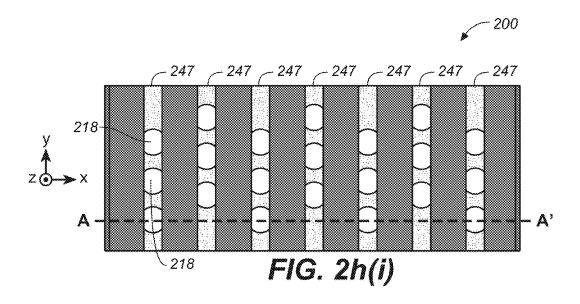


FIG. 2g(ii)



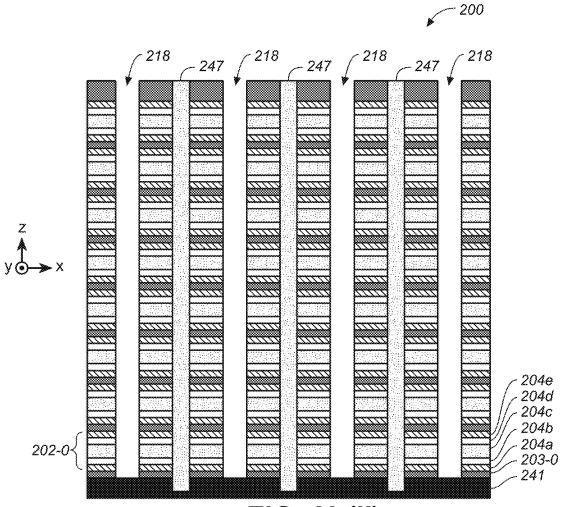
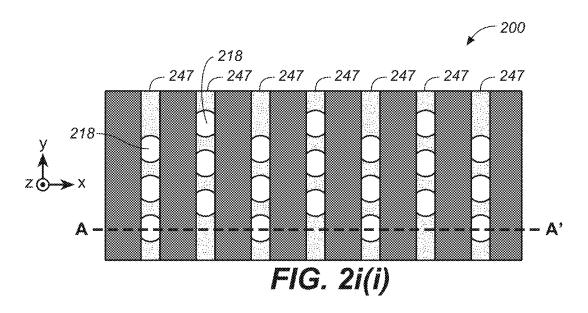


FIG. 2h(ii)



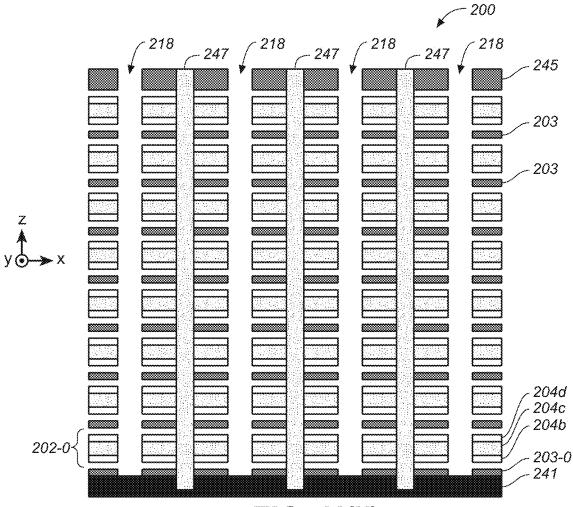
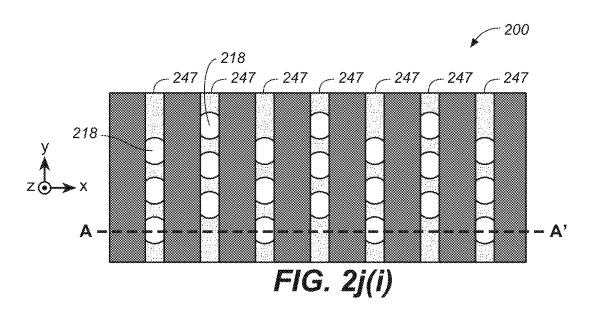


FIG. 2i(ii)



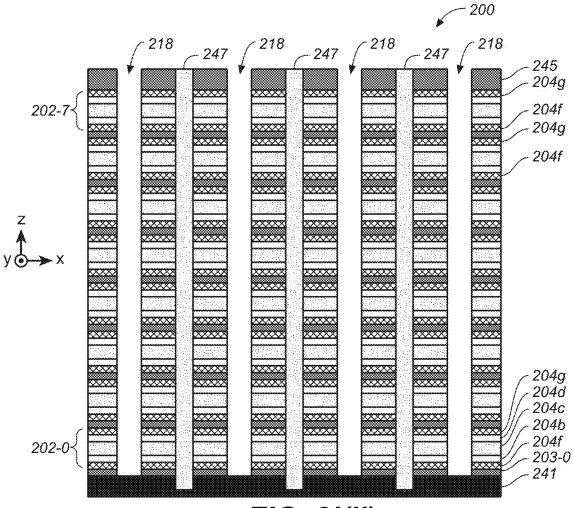
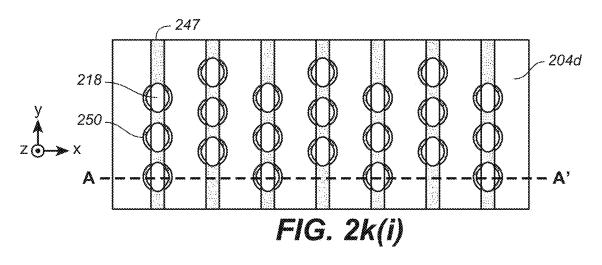


FIG. 2j(ii)



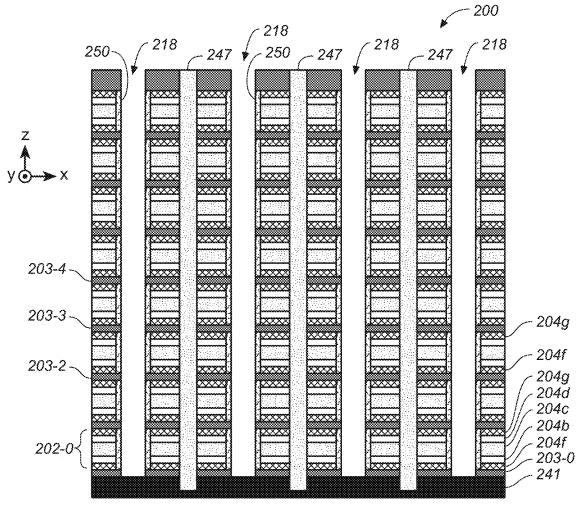


FIG. 2k(ii)

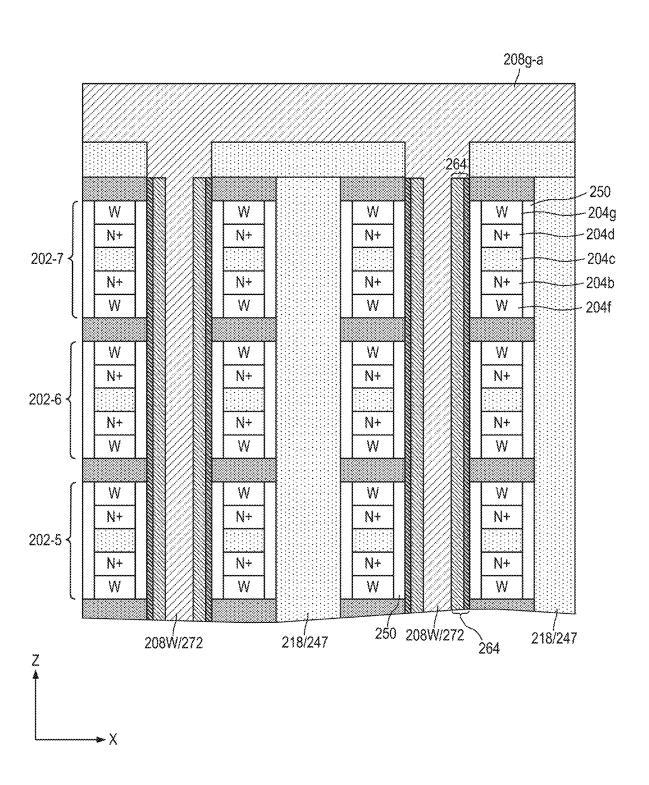
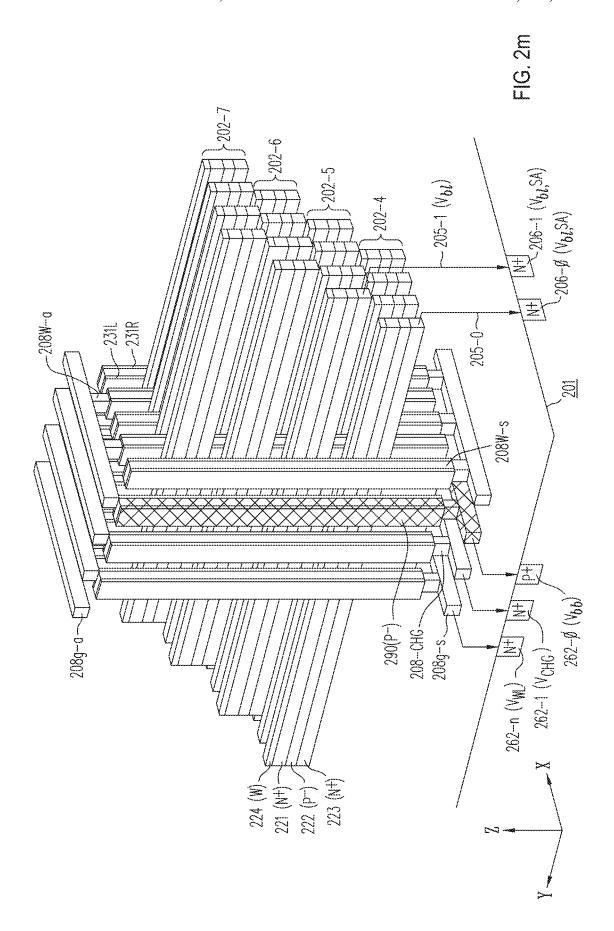


FIG. 21



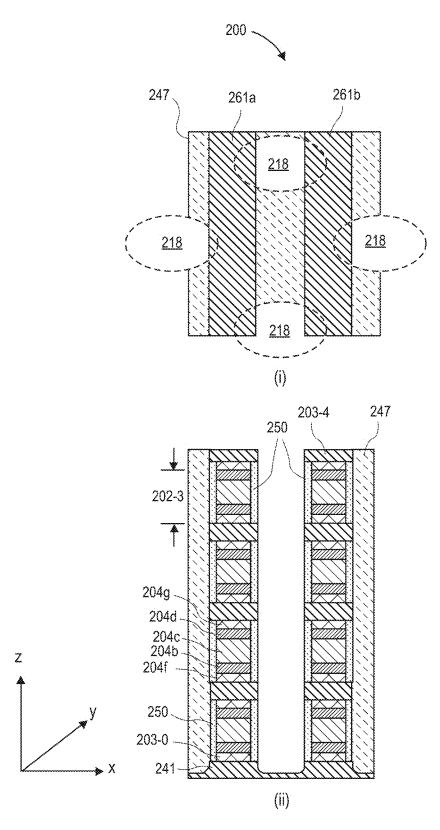
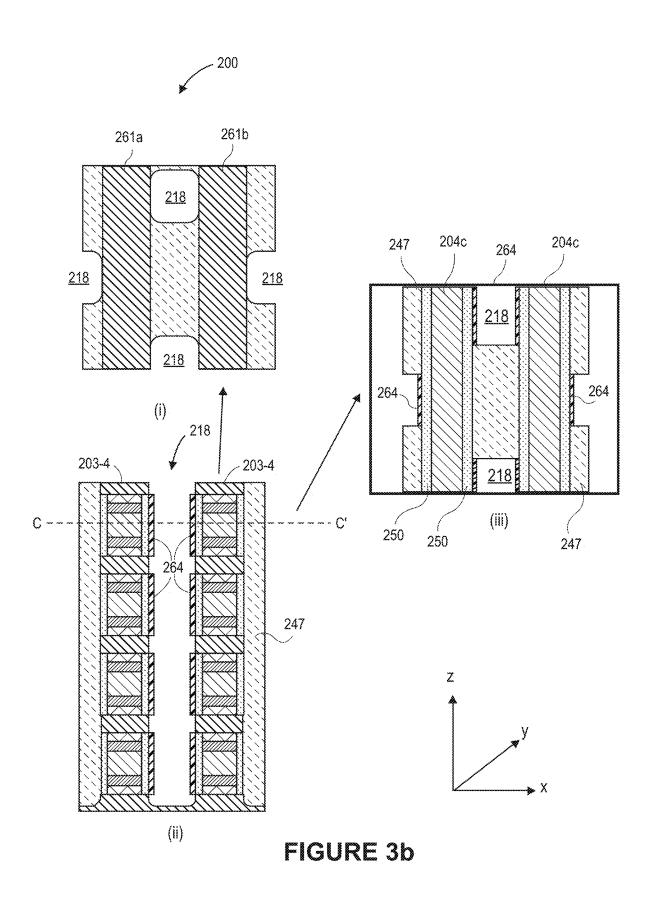
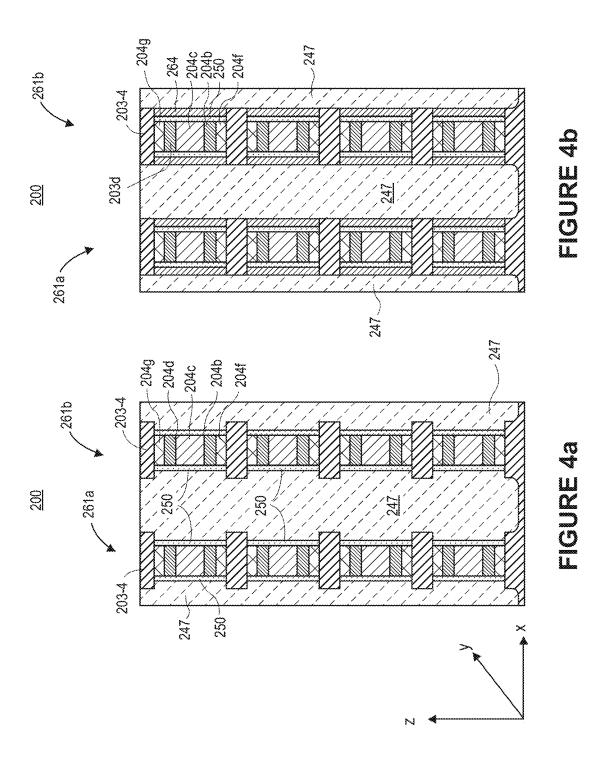
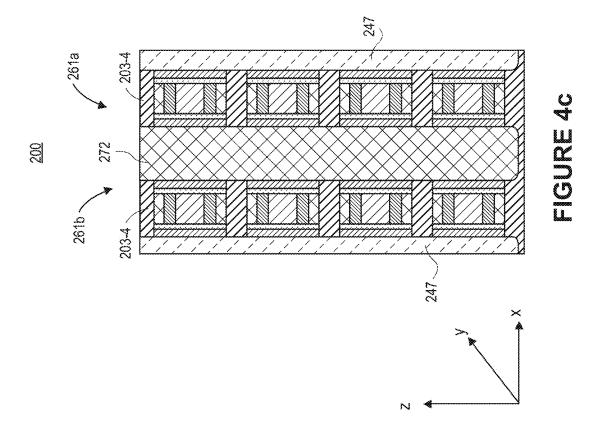
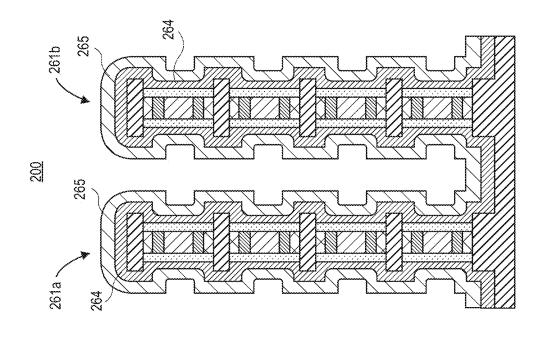


FIGURE 3a

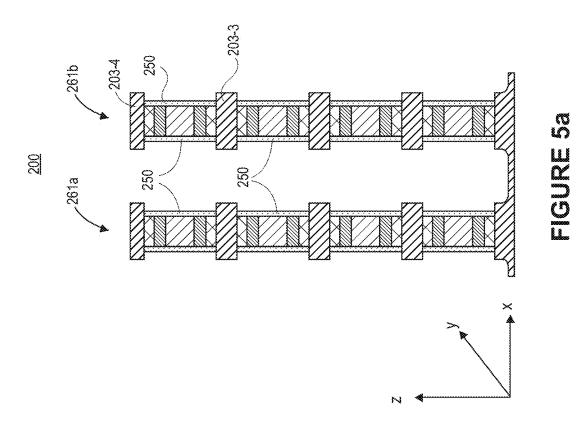


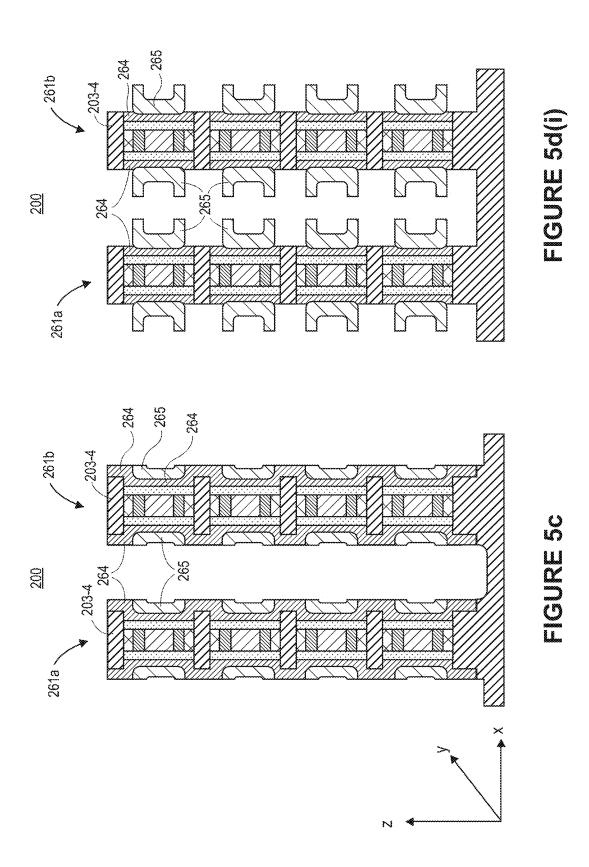


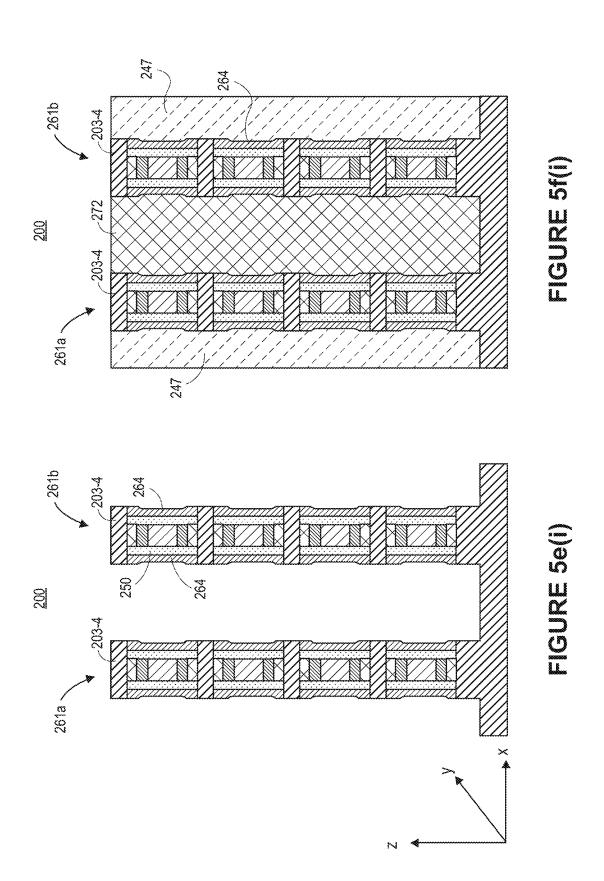


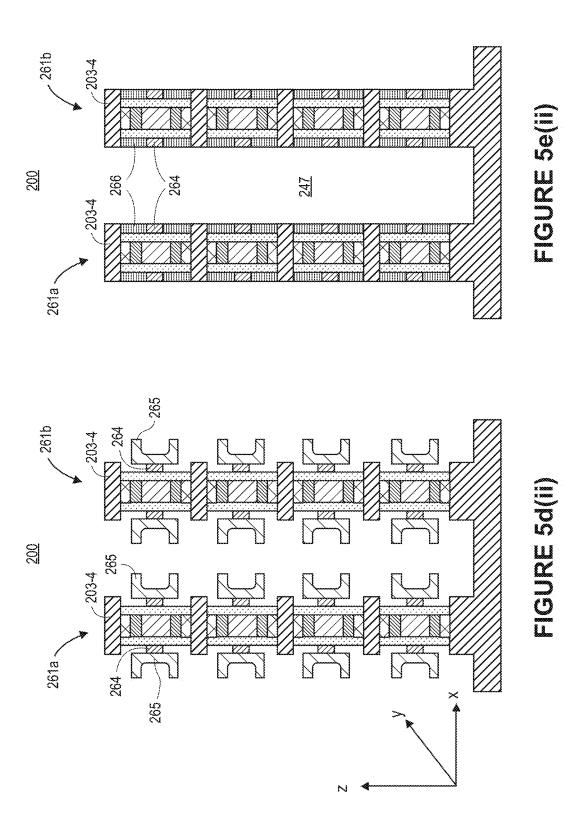


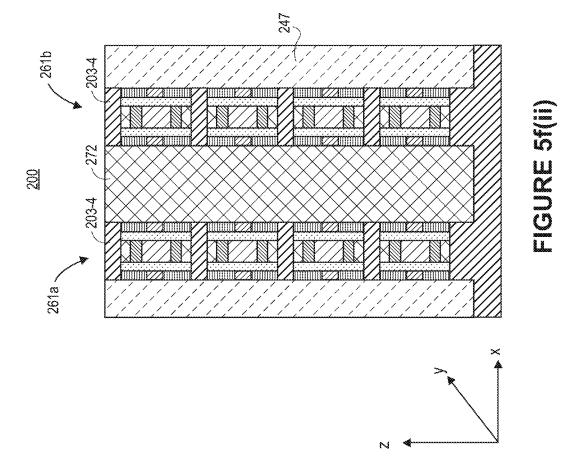
S S C L

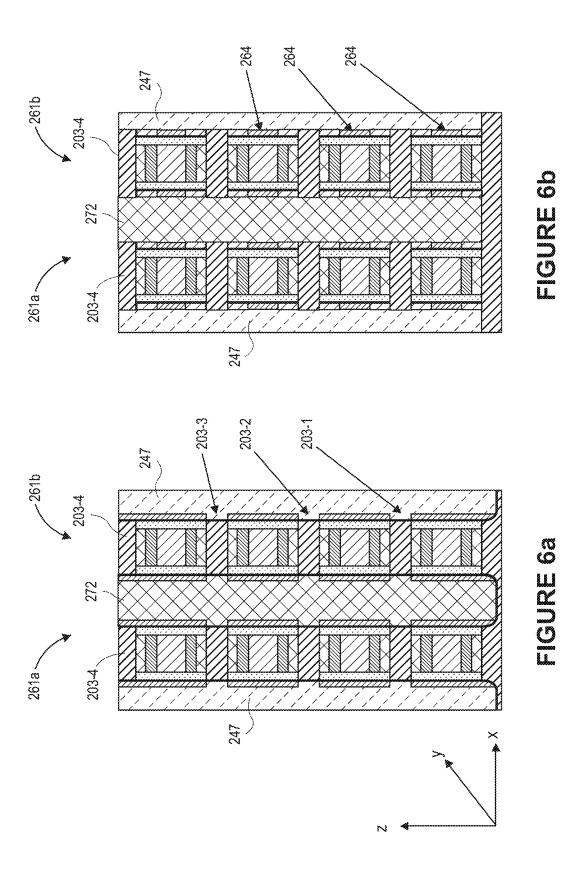












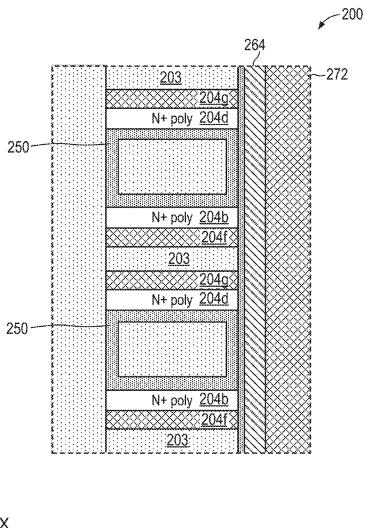


FIG. 6c

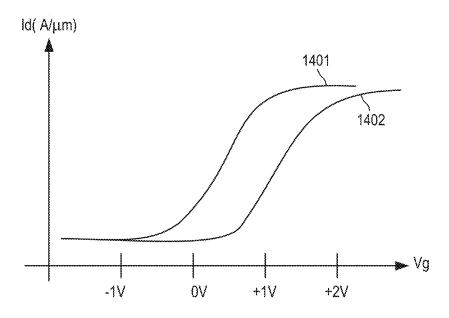


FIG. 7a

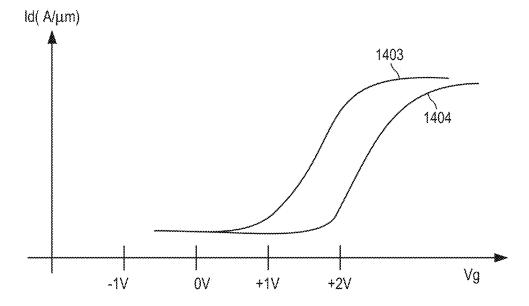


FIG. 7b

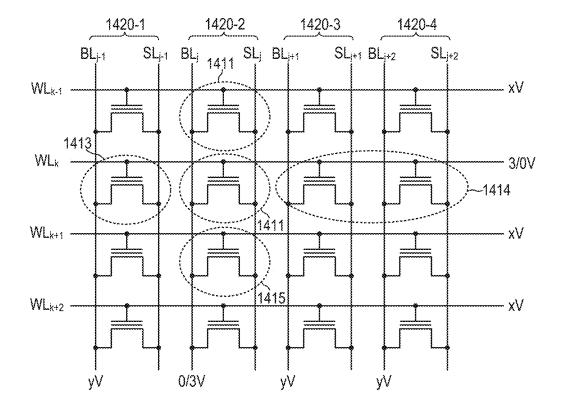
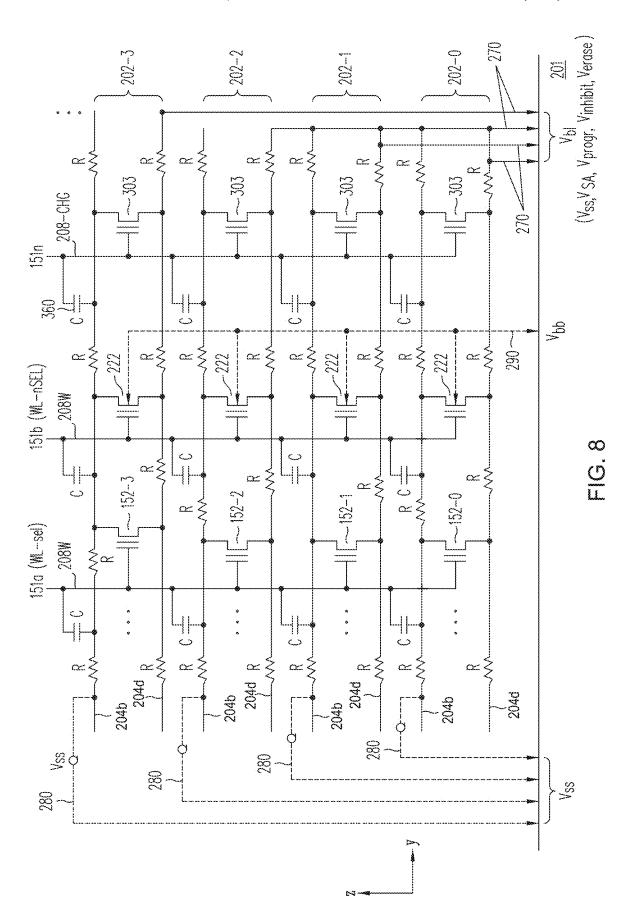


FIG. 7c



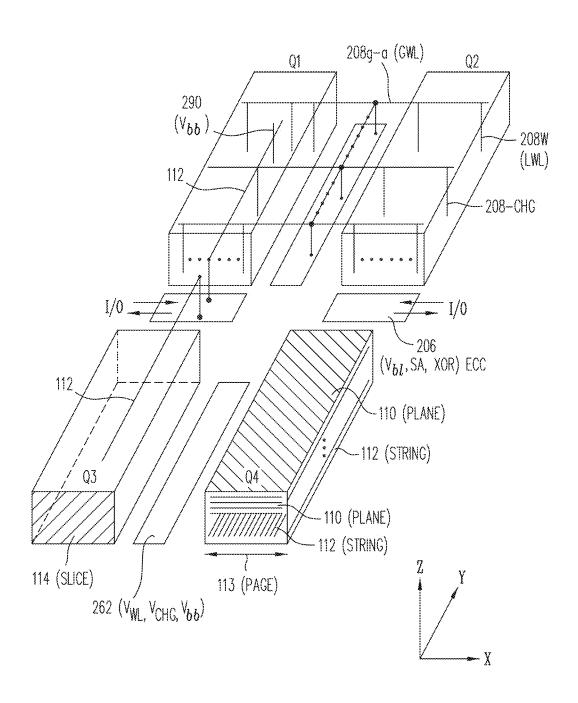


FIG. 9A

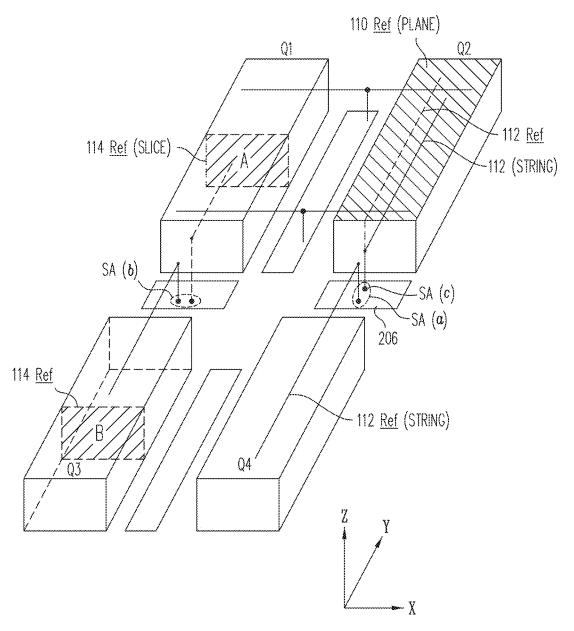


FIG. 9b

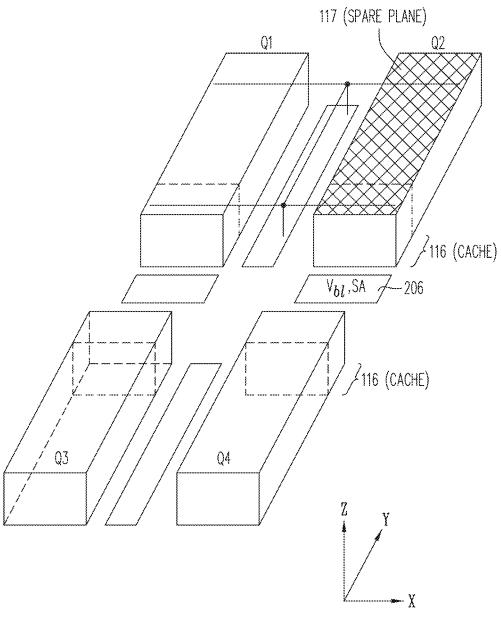
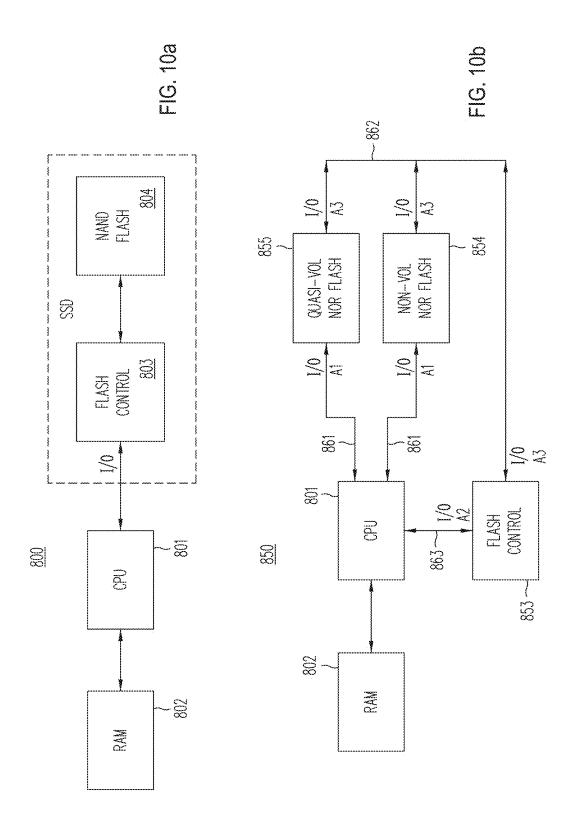
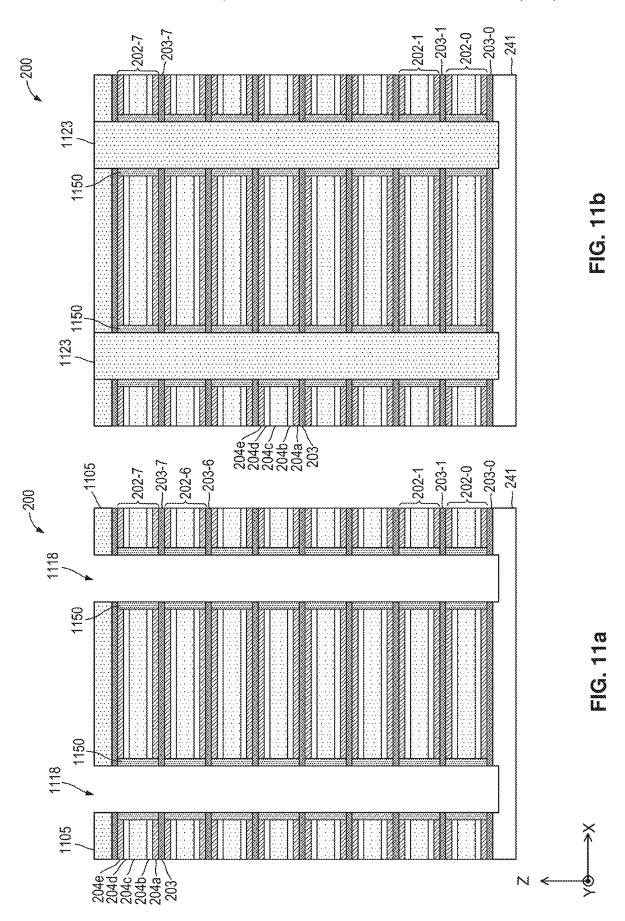
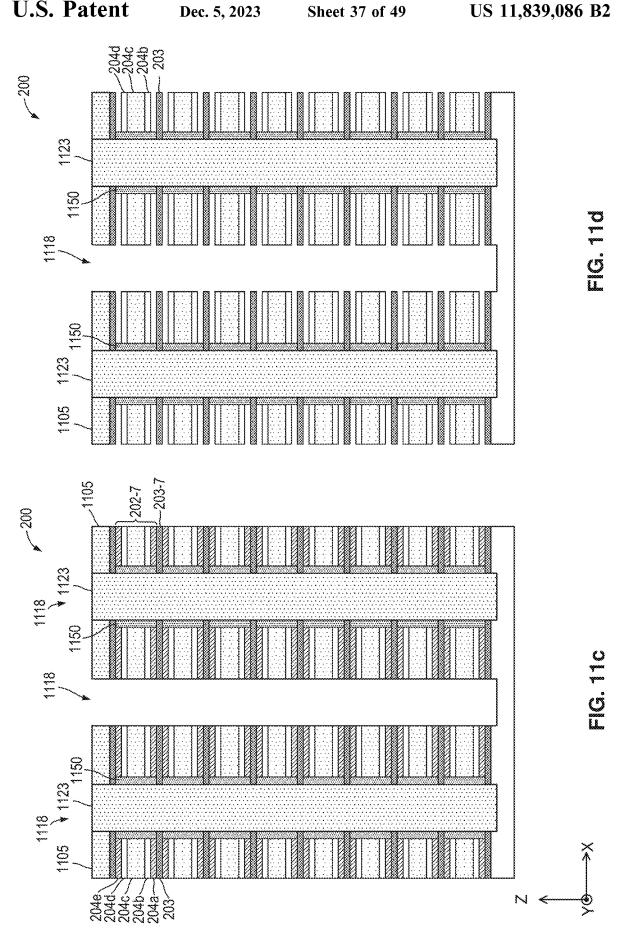
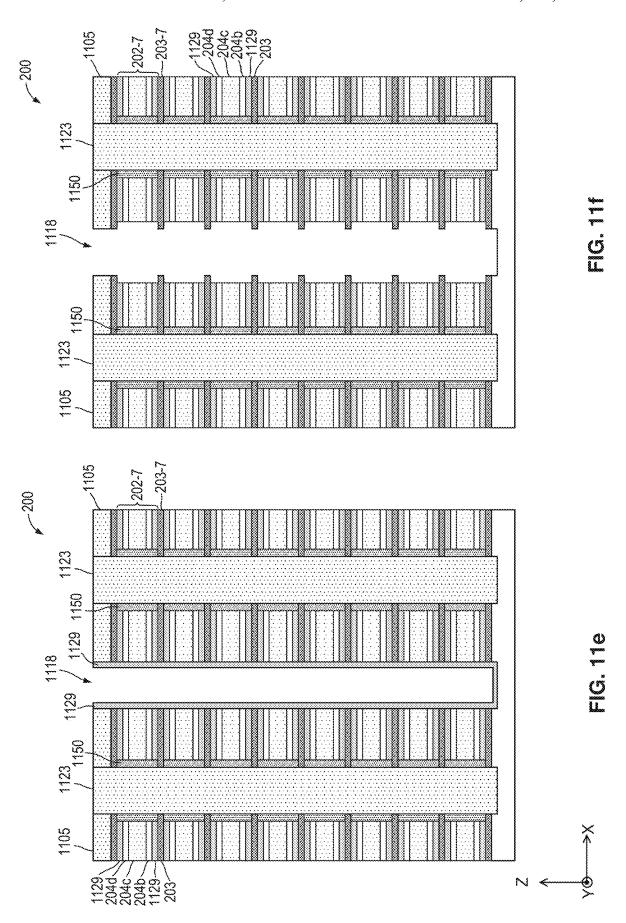


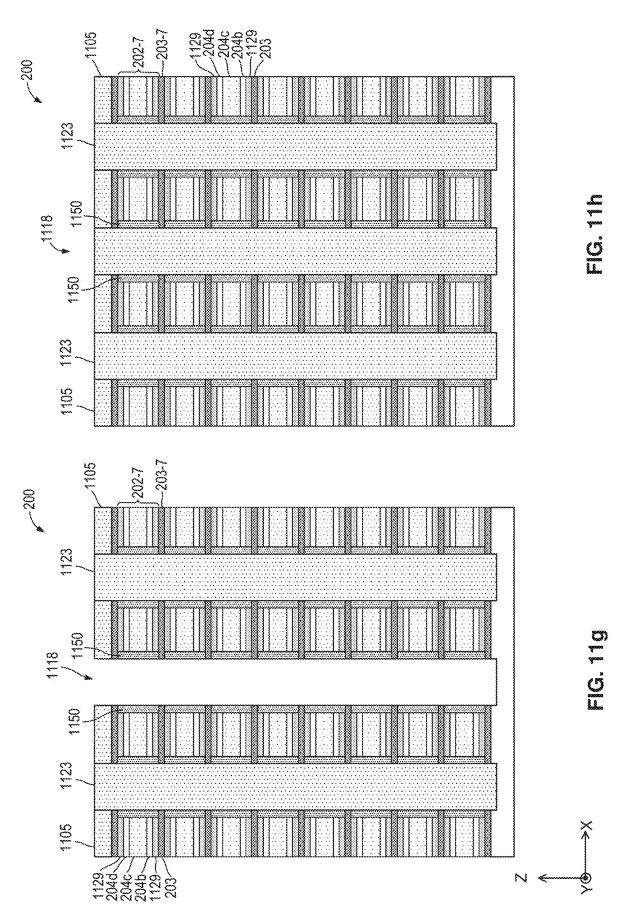
FIG. 9c

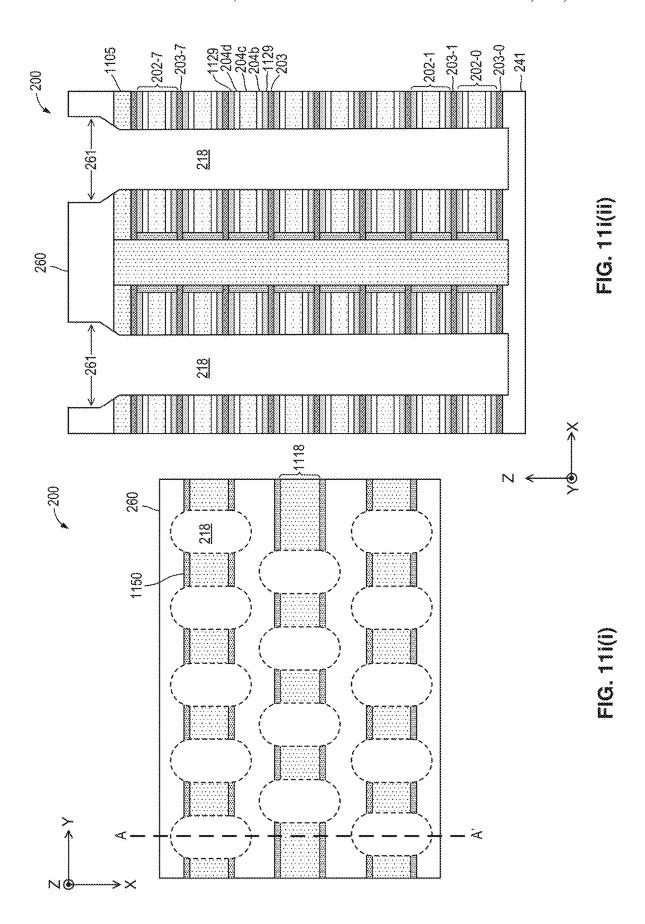


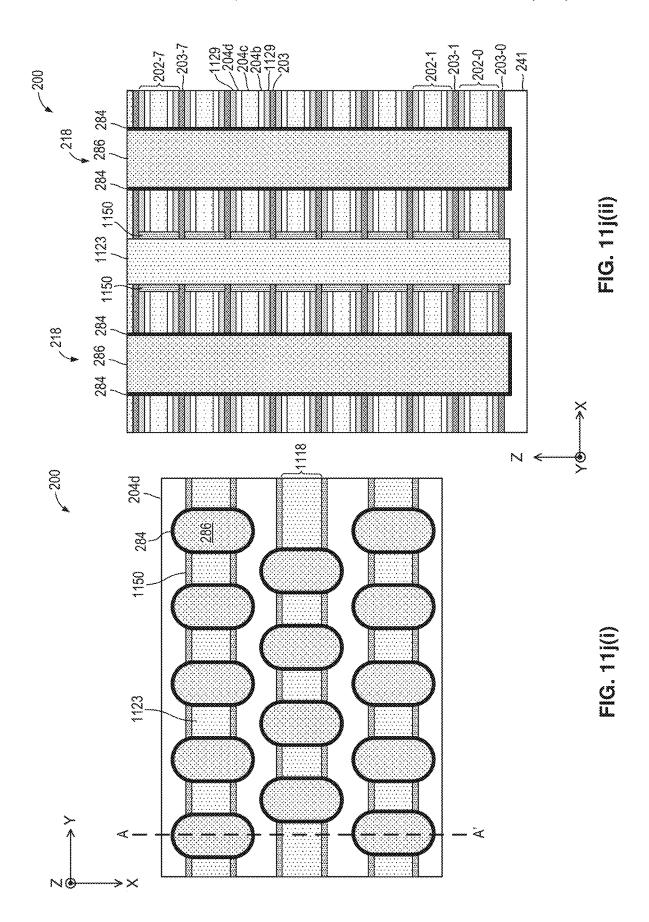


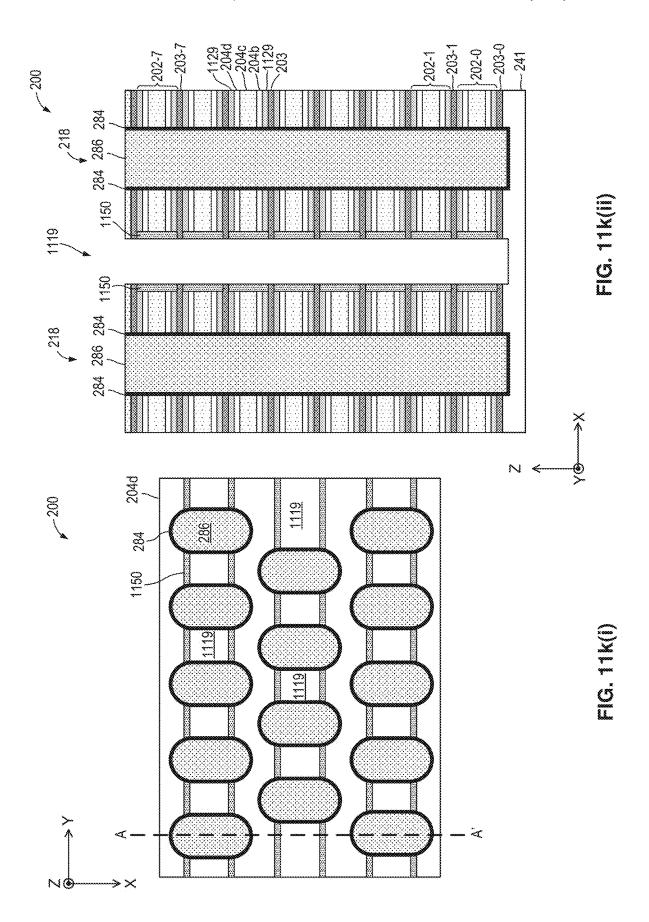


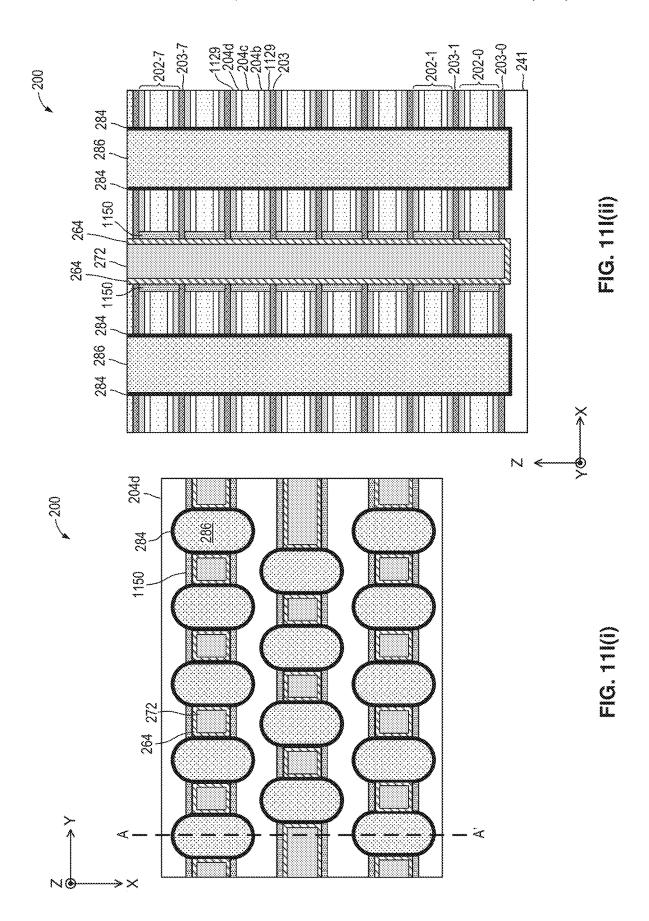


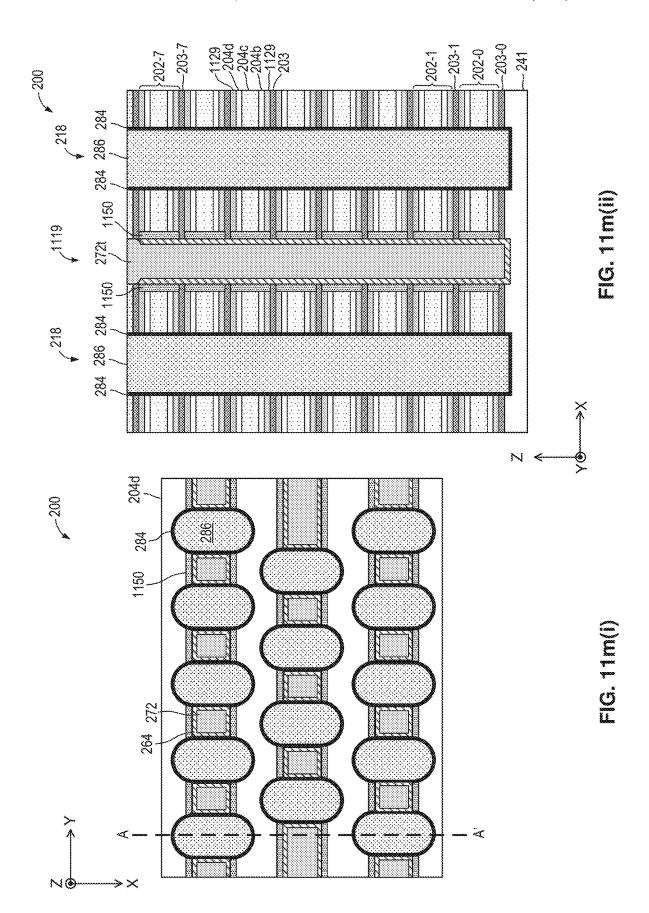


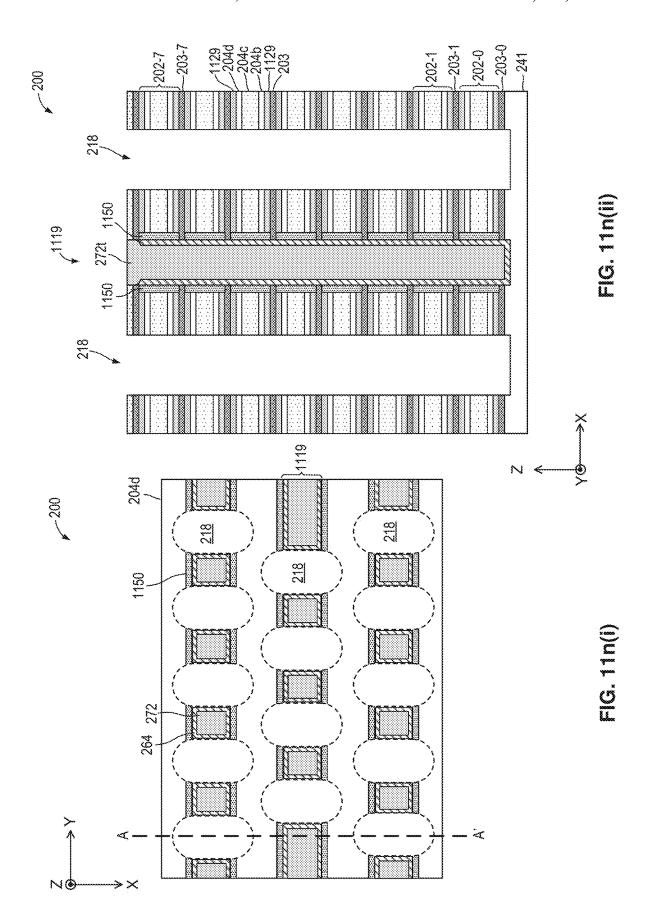


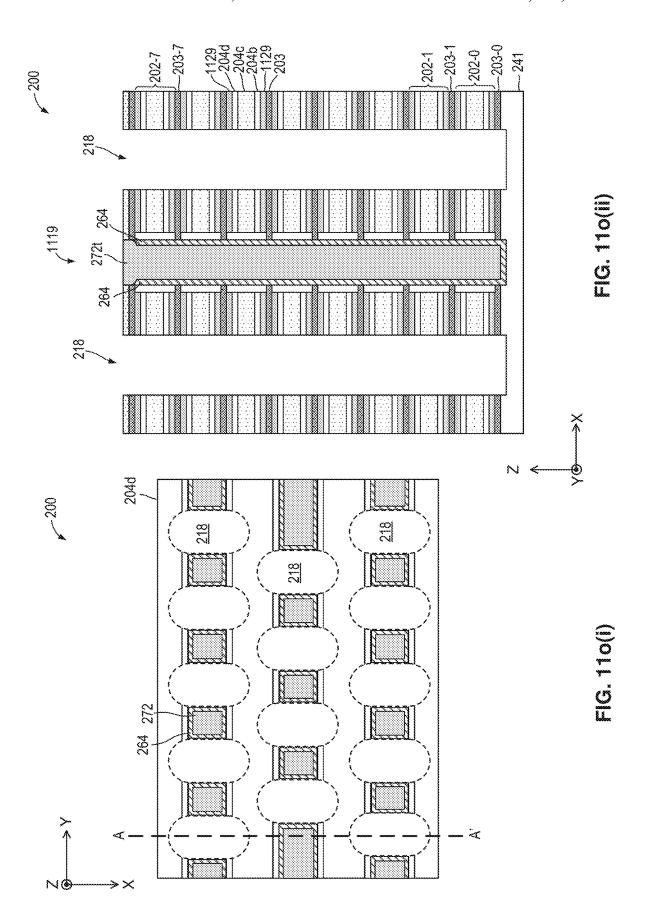


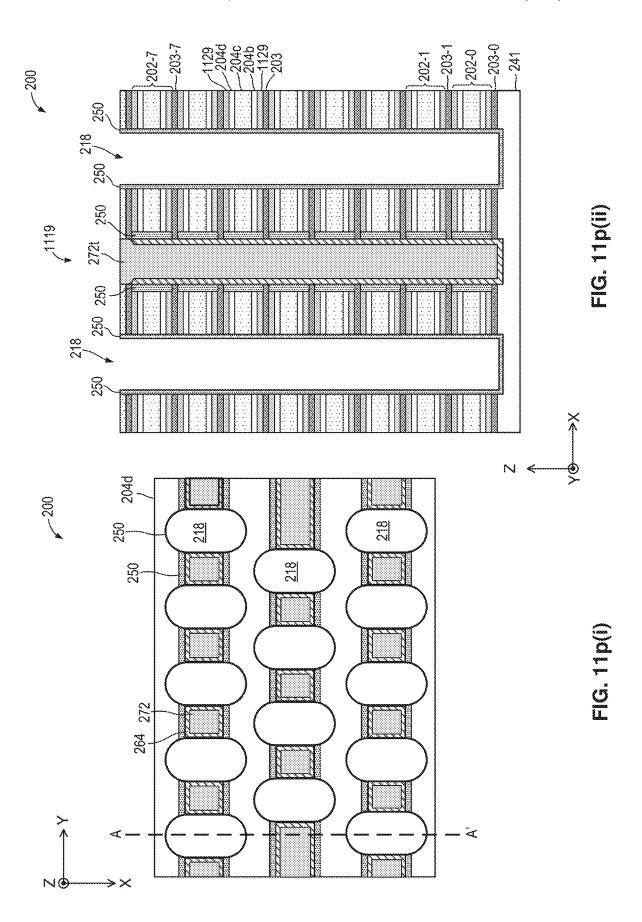


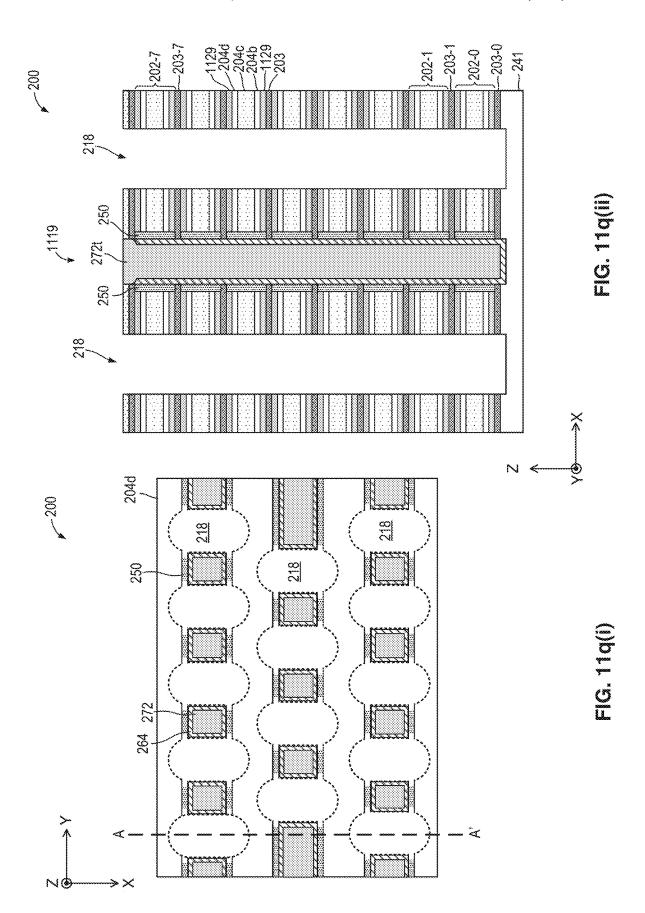


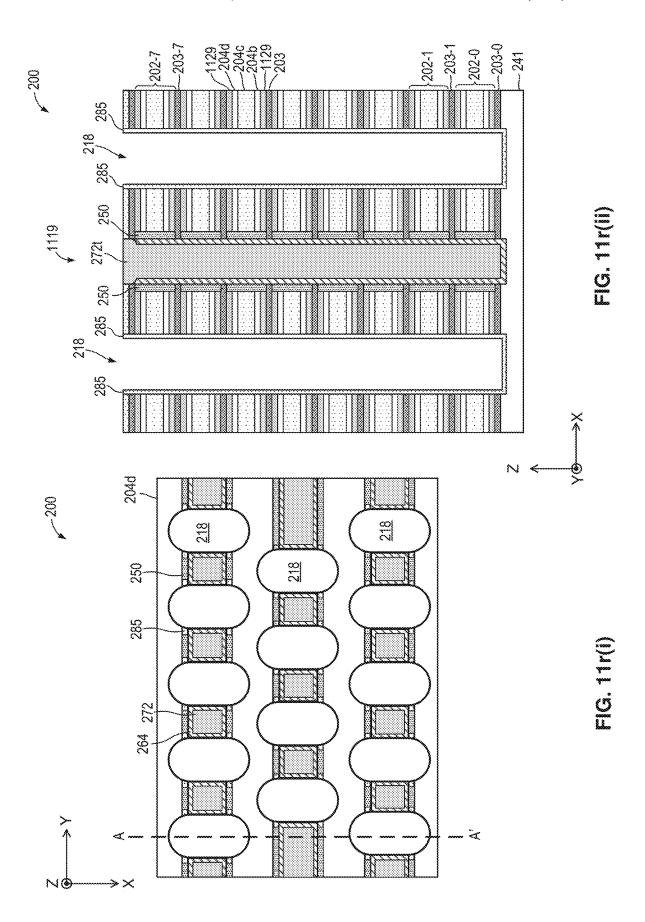












3-DIMENSIONAL MEMORY STRING ARRAY OF THIN-FILM FERROELECTRIC TRANSISTORS

CROSS REFERENCE TO RELATED APPLICATIONS

The present application claims priority to U.S. provisional application ("Parent Application"), Ser. No. 63/222,926, entitled "3-Dimensional Memory String Array of Thin-Film 10 Ferroelectric Transistors," filed on Jul. 16, 2021.

The present application is related to U.S. patent application ("Non-Provisional Application I"), Ser. No. 17/674,137, entitled "Thin-film Storage Transistor With Ferroelectric Storage Layer," filed on Feb. 17, 2022, which claims priority 15 to (i) U.S. provisional application ("Provisional Application I"), Ser. No. 63/152,266, entitled "Thin-Film Storage Transistor With Ferroelectric Storage Layer," filed on Feb. 22, 2021, and (ii) U.S. provisional application ("Provisional Application II"), Ser. No. 63/153,273, entitled "Thin-Film 20 strings"). Storage Transistor With Ferroelectric Storage Layer," filed on Feb. 24, 2021.

The present application is related to U.S. patent application ("Non-Provisional Application II"), Ser. No. 17/382, Memory Structure Of NOR Memory Strings," filed on Jul. 221, 2021, which claims priority to U.S. provisional application ("Provisional Application III"), Ser. No. 63/054,743, entitled "Methods For Fabricating A 3-Dimensional Memory Structure Of NOR Memory Strings," filed on Jul. 30 21, 2020. The present application is also related to U.S. provisional application ("Provisional Application IV"), Ser. No. 63/054,750, entitled "Methods for Fabricating A 3-Dimensional Memory Structure of NOR Memory Strings, filed on Jul. 21, 2020, now U.S. patent application Ser. No. 35 17/382,126, filed Jul. 21, 2021; U.S. provisional application ("Provisional Application V"), Ser. No. 63/139,435, entitled "Vertical NOR Thin-film Transistor Strings and Fabrication Thereof," filed on Jan. 20, 2021, now U.S. patent application sional application ("Provisional Application VI"), Ser. No. 63/114,958, entitled "Methods for Reducing Disturb Errors by Refreshing Data Alongside Programming or Erase Operations," filed on Nov. 17, 2020, now U.S. patent application Ser. No. 17/525,712, filed Nov. 12, 2021.

The present application is also related to U.S. nonprovisional patent application ("Related Application"), Ser. No. 16/894,596, entitled "Capacitive-Coupled Non-Volatile Thin-Film Transistor Strings in Three Dimensional Arrays," filed on Jun. 5, 2020, which is a continuation of U.S. patent 50 application Ser. No. 16/107,118, entitled "Capacitive-Coupled Non-Volatile Thin-Film Transistor Strings in Three Dimensional Arrays," filed on Aug. 21, 2018, which is a divisional application of U.S. non-provisional patent application Ser. No. 15/248,420, entitled "Capacitive-Coupled 55 Non-Volatile Thin-Film Transistor Strings in Three Dimensional Arrays," filed on Aug. 26, 2016, which is related to and claims priority of (i) U.S. provisional application Ser. No. 62/235,322, entitled "Multi-gate NOR Flash Thin-film Transistor Strings Arranged in Stacked Horizontal Active 60 Strips With Vertical Control Gates," filed on Sep. 30, 2015; (ii) U.S. provisional patent application Ser. No. 62/260,137, entitled "Three-dimensional Vertical NOR Flash Thin-film Transistor Strings," filed on Nov. 25, 2015; (iii) U.S. nonprovisional patent application Ser. No. 15/220,375, "Multi- 65 Gate NOR Flash Thin-film Transistor Strings Arranged in Stacked Horizontal Active Strips With Vertical Control

2

Gates," filed on Jul. 26, 2016; and (vi) U.S. provisional patent application Ser. No. 62/363,189, entitled "Capacitive Coupled Non-Volatile Thin-film Transistor Strings," filed Jul. 15, 2016.

The disclosures of Parent Application, Related Application, Non-Provisional Applications I and II, and Provisional Applications I, II, III, IV, V, and VI are hereby incorporated by reference in their entireties.

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates to high-density memory structures. In particular, the present invention relates to high-density, low read-latency memory structures formed by interconnected thin-film storage elements (e.g., 3-dimensional array of thin-film storage transistors), including those organized as NOR-type memory strings ("NOR memory

2. Discussion of the Related Art

The memory circuit structures of this detailed description 064, entitled "Methods For Fabricating A 3-Dimensional 25 may be fabricated above the planar surfaces of one or more semiconductor substrates (e.g., silicon wafers) using conventional fabrication processes. In this description, "vertical" refers to a direction substantially normal to a planar surface of a semiconductor substrate, and "horizontal" refers to any direction that is orthogonal to "vertical." A Cartesian coordinate system may therefore be adopted, in which "horizontal" refers to any direction substantially parallel to the planar surface of semiconductor substrate, which is a surface spanning any two non-colinear horizontal directions X and Y (also referred to as the "X-Y plane"). Similarly, with respect to that Cartesian coordinate system, "vertical" refers to a normal direction (also referred to as the "Z-direction") to the planar surface.

A NOR-type memory string includes storage transistors Ser. No. 17/559,101, filed Dec. 22, 2021; and U.S. provi- 40 that share a common source region and a common drain region, while allowing each storage transistor to be individually addressed and accessed. The Related Application discloses 3-dimensional arrays of NOR memory strings ("NOR memory strings") of charge storage transistors. The Related Application discloses, for example, a NOR memory string that includes (i) a common source region and a common drain region both running lengthwise along a horizontal direction and (i) the gate electrodes for the charge storage transistors each running along a vertical direction.

Advances in electrically polarizable materials ("ferroelectric materials"), especially those used in electronic circuits, suggest new potential applications in ferroelectric memory circuits. For example, the article "Ferroelectricity in Hafnium Oxide: CMOS compatible Ferroelectric Field Effect Transistors," by T. S. Böscke et al., published in 2011 International Electron Devices Meeting (IEDM), pp. 24.5.1-24.5.4 (doi: 10.1109/IEDM.2011.6131606), discloses a ferroelectric field effect transistor ("FeFET") that uses hafnium oxide as a gate dielectric material. By controlling polarization in a ferroelectric gate dielectric layer, the FeFET may be programmed to have one of two or more selectable threshold voltages. Each threshold voltage of the FeFET constitutes a state (e.g., a "programmed" state, an "erased" state, or any other state that represents a designated binary value). Such an FeFET has application in high-density memory circuits. For example, U.S. patent application Ser. No. 13/897,037, entitled "Apparatuses having a ferroelectric field-effect tran-

sistor memory array and related method," by D. V. Nirmal Ramaswamy et al., filed on May 17, 2013, now U.S. Pat. No. 9,281,044, discloses a 3-dimensional array of FeFETs.

The FeFETs of the prior art, however, suffer from low endurance. For example, the article "Vertical Ferroelectric 5 HfO₂ FET based on 3-D NAND Architecture: Towards Dense Low-Power Memory," by K. Florent et. al., published in 2018 IEEE International Electron Devices Meeting pp. 2.5.1-2.5.4 2018, (doi: IEDM.2018.8614710), discloses an endurance of merely 10^4 10 cycles. Such low endurance renders the memory circuits practically unsuitable for many memory applications. It is highly desirable to have a memory circuit that can serve as a building a block of a main memory system. Current building blocks of main memory systems (e.g., dynamic 15 random-access memory (DRAM)) have low retention time and low memory density.

SUMMARY

According to one embodiment of the present invention, an array of memory cells includes FeFETs formed in stacks of horizontal active strips, with vertical control gate electrodes (also referred to as "local word lines") provided along one or both sidewalls of the active strips, the control gate 25 electrodes being separated from their associated active strips by one or more polarizable or ferroelectric elements (also referred to as "ferroelectric or polarization layers"). Polarization of the polarizable element in an FeFET changes the threshold voltage the FeFET. Each active strip may include 30 a channel layer formed adjacent both a common source region and a common drain region. In one embodiment, in a given stack of active strips (also referred to as an "active stack"), the polarizable or ferroelectric elements associated with the FeFETs of each active strip are separated from the 35 polarizable or ferroelectric elements associated with the FeFETs of other active strips. The FeFETs of each active strip may be organized as one or more NOR memory strings. Further, independently addressable FeFETs may be formed along opposite sides of an active strip to achieve twice the 40 memory density than if the FeFETs are formed as a single file along the length of the common bit line. Furthermore, by forming channel regions of the FeFETs along the vertical sidewalls of each active strip (i.e., between the common drain region and the common source region), a memory 45 density that is substantially higher than is achieved in the prior art is possible because the volume of the channel region, even with a large extent along the Z-direction, has an insignificant footprint in the X-Y plane.

In one embodiment, only the common drain region (also 50 referred to as a "common bit line") of the FeFETs in a NOR memory string is electrically and selectively connected by one or more conductors to one of several voltage sources. In one embodiment, the common source region may be precharged to a predetermined voltage prior to a read, program 55 (sometimes also referred to as "write"), or erase operation, and is held to that predetermined voltage by the parasitic capacitance associated with the common source region during the read, program (sometimes also referred to as "write"), or erase operation. In one embodiment, one or 60 more arrays of NOR memory strings may be organized and formed as an independently addressable modular memory circuit (also referred herein as a "tile"). A group of tiles may be formed above the same substrate but are configured to each be independently controlled and operated, thereby 65 allowing massively parallel read, program or erase operations to be carried out concurrently on the tiles. The NOR

4

memory strings within each tile may also be organized in groups that are configured to each be independently controlled and operated.

In one embodiment, in a 3-dimensional array of NOR memory strings, FeFETs are formed on both side edges along the lengths of each active strip, with vertical local word lines being provided along both side edges. High density is achieved by sharing the local word lines between their adjacent active stacks, each local word line serving as a common gate electrode to vertically aligned FeFETs in the active strips of both active stacks. The vertical local word lines may be contacted by interconnection conductors (also referred to as "global word lines") provided above or below the 3-dimensional array of NOR memory strings. The global word lines may run along a direction transverse to the lengths of the contacted active strips.

Organizing the FeFETs as 3-dimensional arrays of NOR memory strings—rather than as NAND memory strings, for example—results in (i) reduced read-latencies approaching those of dynamic random access memory (DRAM) arrays, (ii) reduced sensitivities to read-disturb and program-disturb conditions, (iii) reduced power dissipation and a lower cost-per-bit metric relative to planar NAND or DRAM memory arrays, and (iv) the ability to read, write or erase FeFETs on multiple active strips concurrently, thereby significantly increasing data throughput.

According to one embodiment of the present invention, undesirable effects on a memory operation due to variations among FeFETs within an array of NOR memory strings may be overcome by processing the memory operation in conjunction with FeFETs in a designated reference NOR memory string within the same array. For example, effects on a read operation of an FeFET of a NOR memory string due to a background leakage current present in NOR memory strings can be substantially mitigated by comparing the sensed result of a concurrently read FeFET within a reference NOR memory string. In some embodiments, the ferroelectric or polarizable elements may have a data retention time that requires refreshing to meet data stability requirements in some applications. Relative to storage transistors in conventional non-volatile memory circuits (e.g., non-volatile NAND memory strings), which do not require refreshing in those applications, the FeFETs may be considered "quasi-volatile" (QV) storage transistors. The FeFETs are considered quasi-volatile because, as compared to conventional volatile memory circuits (e.g., DRAM circuits), the FeFETs require refreshing significantly less frequently. For example, FeFETs in high density arrays need only be refreshed, restored, or replaced, once every few minutes, hours or every few days, and may be accomplished by accessing the FeFETs individually or accessed concurrently as one or more groups. Such refresh operations ensure long term stable and reliable operations. The refresh operations may also be performed in the background, without interfering regular user read, erase or write operations.

The present invention is better understood upon consideration of the detailed description below, in conjunction with the accompanying drawings.

BRIEF DESCRIPTION OF THE DRAWINGS

Various embodiments of the invention are disclosed in the following detailed description and the accompanying drawings. Although the drawings depict various examples of the invention, the invention is not limited by the depicted examples. It is to be understood that, in the drawings, like

reference numerals designate like structural elements. Also, it is understood that the depictions in the FIGS. are not necessarily to scale.

FIG. 1, which includes FIGS. 1a-1 and 1a-2, show conceptualized memory structure 100, which illustrates an 5 organization of ferroelectric field-effect transistors (FeFETs) in an array of NOR memory strings, according to one embodiment of the present invention.

FIG. 1b shows a basic circuit representation of four NOR memory string-pairs, each NOR memory string-pair including two NOR memory strings provided on the same plane (i.e., one of planes 159-1, 159-2, 159-3 and 159-4), according to one embodiment of the present invention.

FIG. 1c shows a basic circuit representation of four NOR memory strings, each NOR memory string being provided 15 on a respective one of planes 159-1, 159-2, 159-3 and 159-4, according to one embodiment of the present invention.

FIG. 2a shows a cross section in a Y-Z plane of semiconductor structure 200, after active layers 202-0 to 202-7 (each separated from the next active layer by one of isolation 20 layers 203-0 to 203-7) have been formed on semiconductor substrate 201, but prior to formation of individual active strips, in accordance with one embodiment of the present invention.

FIGS. 2b-1, 2b-2, 2b-3, 2b-4 and 2b-5 show various 25 examples of active layer 220a, according to one embodiment of the present invention.

FIG. 2c shows cross section in a Y-Z plane of structure 200 of FIG. 2a through buried contacts 205-0 and 205-1, which connect n^+ silicon layers 204b of active layers 202-0 30 and 202-1 to circuitry 206-0 and 206-1 in semiconductor substrate 201.

FIGS. 2d(i), 2d(ii), 2e(i), 2e(ii), 2f(i), 2f(ii), 2g(i), 2g(ii), 2h(ii), 2h(ii), 2i(i), 2i(ii), 2j(i), 2j(ii), 2h(i) and 2h(ii) illustrate steps of an exemplary process for forming a modular circuit 35 or tile that includes an array of NOR memory strings, including an array portion (i.e., memory structure 200), according to one embodiment of the present invention; in particular, FIG. 2h illustrates, preceding steps that form the gate stacks of the FeFETs, a step for providing a channel 40 region for each FeFET; the channel region may be provided by any one group of the groups of process steps described in conjunction with FIGS. 3a to 3b, FIGS. 4a to 4c and FIGS. 5a to 5f, respectively.

FIG. 2*l* shows an X-Z plane cross section through a row 45 of local word lines 208*g-a*, showing active strips formed out of active layers 202-7 and 202-6, according to one embodiment of the present invention.

FIG. 2m shows a 3-dimensional view of horizontal active layers 202-4 to 202-7 in the embodiment of FIG. 2k, with 50 local word lines or local pre-charge word lines connected to global word lines 208g-s and global word lines 208g-a, and showing each active strip as having its N⁺ layer 204d (acting as a drain region) connected through select circuits to any of voltage supplies (e.g., V_{ss} , V_{bl} , V_{pgm} , $V_{inhibit}$ and V_{erase}), 55 decoding, sensing and other circuits arranged either adjacent or directly underneath the memory arrays.

FIGS. 3a and 3b illustrate a first group of process steps for forming the gate stacks in the FeFETs of an array of NOR memory strings, according to one embodiment of the present 60 invention.

FIGS. 4a, 4b and 4c illustrate a second group of process steps for forming the gate stacks in the FeFETs in an array of NOR memory strings, according to one embodiment of the present invention.

FIGS. 5a, 5b, 5c, 5d(i), 5d(ii), 5e(i), 5e(i), 5f(i) and 5f(ii) illustrate a third group of process steps for forming the gate

6

stacks in the FeFETs of an array of NOR memory strings, according to one embodiment of the present invention.

FIG. 6a shows an X-Z cross sectional view of memory structure 200, which represents active stacks 261a and 261b fabricated under any group of process steps illustrated by FIG. 3a-3b, FIGS. 4a to 4c and FIGS. 5a to 5f(i), in accordance with one embodiment of the present invention.

FIG. 6b shows an X-Z cross sectional view of memory structure 200, which represents active stacks 261a and 261b fabricated under the group of process steps illustrated by FIGS. 5a-5f(ii), in accordance with one embodiment of the present invention.

FIG. 6c shows an X-Z cross sectional view of memory structure 200, fabricated under a channel-last fabrication process, in accordance with one embodiment of the present invention.

FIG. 7a shows hysteresis in the drain current (I_d) in response to an applied gate voltage (V_g) in a conventional FeFET.

FIG. 7b shows a desirable hysteresis in the drain current (I_d) response to an applied gate voltage (V_g) in a thin-film FeFET in a NOR memory array, according to one embodiment of the present invention.

FIG. 7c shows a circuit schematic diagram of FeFETs in NOR memory strings **1420-1** to **1420-4**, according to one embodiment of the present invention.

FIG. 8 illustrates the methods and circuit elements used for setting source voltage V_{ss} in n^+ silicon layers **204***d*; specifically, source voltage V_{ss} may be set through hard-wire decoded source line connections **280** (dashed line) or alternatively, by activating pre-charge FeFETs **303** and decoded bit line connections **270**, to any one of voltage sources for bit line voltages V_{ss} , V_{bl} , V_{pgm} , $V_{inhibit}$ and V_{erase} .

FIG. 9a shows semiconductor structure 600, which is a three-dimensional representation of an array of NOR memory strings organized into quadrants Q1-Q4; in each quadrant, (i) numerous NOR memory strings are each formed in an active strip extended along the Y-direction (e.g., NOR memory string 112), (ii) pages extending along the X-direction (e.g., page 113), each bit in the page corresponds to an FeFET from each NOR memory string at a corresponding Y-position, the NOR memory strings in the page being of the same corresponding Z-position (i.e., of the same active layer); (iii) slices extending in both the X- and Z-directions (e.g., slice 114), with the bits in each slice provided by the pages of the same corresponding Y-position, one page from each of the planes, and (iv) planes extending along both the X- and Y-directions (e.g., plane 110), the bits in each plane being provided by all pages at a given Z-position (i.e., formed out of the same active layer).

FIG. 9b shows structure 600 of FIG. 9a, showing FeFETs in programmable reference string 112-Ref in quadrant Q4 and FeFETs in NOR memory string 112 in quadrant Q2 coupled to sense amplifiers SA(a), Q2 and Q4 being "mirror image quadrants"; FIG. 9b also shows (i) programmable reference slice 114-Ref (indicated by area A) in quadrant Q3 similarly providing corresponding reference FeFETs for slice 114 in mirror image quadrant Q1, sharing sense amplifiers SA(b), and (ii) programmable reference plane 110-Ref in quadrant Q2 providing corresponding reference FeFETs to plane 110 in mirror image quadrant Q1, sharing sense amplifiers SA(c), and also providing corresponding reference FeFETs for NOR memory strings in the same quadrant (e.g., NOR memory string 112).

FIG. 9c shows structure 600 of FIG. 9a, showing slices 116 being used as a high speed cache because of their close proximity to their sense amplifiers and voltage sources 206;

FIG. 9c also show spare planes 117, which may be used to provide replacement or substitution NOR memory strings or pages in quadrant O2.

FIG. 10a shows in simplified form prior art storage system 800 in which microprocessor (CPU) 801 communicates with system controller 803 in a flash solid state drive (SSD) that employs NAND flash chips 804; the SSD emulates a hard disk drive and NAND flash chips 804 do not communicate directly with CPU 801 and have relatively long read latency.

FIG. 10b shows in simplified form system architecture 850 using the memory devices of the present invention, in which FeFETs of NOR memory string arrays 854 and 855 communicate directly with CPU 801 through one or more input and output (I/O) ports 861, and indirectly through 15 controller 863.

FIGS. 11a, 11b, 11c, 11d, 11e, 11f, 11g, 11h, 11i(i), 11i(ii), 11j(i), 11j(ii), 11k(i), 11k(ii), 11l(i), 11l(ii), 11m(i), 11m(ii), 11n(ii), 11o(i), 11o(ii), 11p(i), 11p(ii), 11q(i), 11q(ii), 11r(i), and 11r(ii) illustrate an additional process in which, 20 subsequent to forming ferroelectric layers, the channel regions for the FeFETs of an array of NOR memory strings are formed by replacing an earlier-formed sacrificial layer, according to one embodiment of the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

In this detailed description, process steps described for one embodiment may be used in a different embodiment, 30 even if the process steps are not expressly described in the different embodiment. When reference is made herein to a method including two or more defined steps, the defined steps can be carried out in any order or simultaneously, except where the context dictates or specific instruction 35 otherwise are provided herein. Further, unless the context dictates or express instructions otherwise are provided, the method can also include one or more other steps carried out before any of the defined steps, between two of the defined steps, or after all the defined steps

The present inventors realize that certain polarizable materials (also referred to as "ferroelectric materials") may be used to form storage elements in memory applications. Such ferroelectric materials include, for example, hafnium oxide, which have been used in forming charge-trapping 45 layers in non-volatile or quasi-volatile memory circuits. According to one embodiment of the present invention, high-density data storage may be achieved using ferroelectric field-effect transistors ("FeFETs") as the storage transistors of a 3-dimensional memory array (e.g., an array of 50 NOR memory strings). The present invention, as illustrated herein, provides high endurance, long data retention, and relatively low voltage operations for both erase (under 7.0 volts) and programming (e.g., under -7.0 volts) operations. By combining their ferroelectric or polarization character- 55 istics with the 3-dimensional organization (e.g., as thin-film NOR memory strings described herein), memory arrays formed out of the FeFETs of the present invention achieve the additional benefits of high-density, low-cost, high-speed randomly accessed memory circuits (e.g., low read latency). 60

FIGS. 1a-1 and 1a-2 show conceptualized memory structure 100, which illustrates an organization of ferroelectric field-effect transistors (FeFETs) in an array of NOR memory strings, according to one embodiment of the present invention. This detailed description is facilitated by adopting as 65 reference a Cartesian coordinate system, as indicated in FIGS. 1a-1 and 1a-2. Conceptualized memory structure 100,

8

provided herein as an abstraction of certain salient characteristics of a memory structure of the present invention, represents an array of NOR memory strings fabricated over a surface of substrate 101. Conceptualized memory structure 100, being an abstraction, is not drawn to scale along any of the X-, Y-, and Z-directions. Substrate 101 may be, for example, a conventional silicon wafer having a top layer with a planar surface for fabricating integrated circuits, as is familiar to those of ordinary skill in the art. Thus, under this adopted Cartesian coordinate system, the term "horizontal" refers to any direction parallel to the planar surface of substrate 101, which is a surface of an X-Y plane. Similarly, the term "vertical" refers to the Z-direction or a normal to the planar surface.

As shown in both FIGS. 1a-1 and 1a-2, four planes (e.g., plane 110) of conceptualized memory structure 100 are isolated from each other and stacked one on top of another in the vertical direction. FIGS. 1a-1 and 1a-2 each show four horizontal active strips of NOR memory strings in each plane (e.g., active strip 112). As discussed below in further detail in this detailed description, each active strip—which includes several material layers extending lengthwise along the Y-direction—may have formed therein one or more NOR memory strings. Each NOR memory string includes numerous FeFETs (e.g., FeFET 111) sharing a common drain region and a common source region, the FeFETs being formed side-by-side along one or both sides of the active strip, with the thin-film FeFET current flowing through a channel region along the vertical direction. In an FeFET, the polarization in the ferroelectric element affects its threshold voltage.

In data storage or memory applications, two different threshold voltage values are selected to represent the values of a binary data bit. In such applications, the polarization states of the FeFET may be referred to as being in the "erased state" or the "programmed" state. In a NOR memory string, programming (also referred to as "writing"), reading or erasing any one of the FeFETs does not require activating any other FeFETs in the same NOR memory string. Accord-40 ingly, each FeFET of a NOR memory string is randomly addressable and accessible. FeFETs in close proximity that are not intended for the read, program or erase operation, however, are protected from unintended programming or erasing by having their transistor terminals suitably biased, often referred to as "program-inhibit" or "erase-inhibit" operations. Furthermore, an FeFET that has been accessed repeatedly may experience a gradual shift in its threshold voltage due to a phenomenon commonly referred to as "read-disturb." Indeed, even programmed or erased FeFETs need to be protected from shifts in their threshold voltages ("disturbed") during read, write or erase operations on nearby FeFETs. One effective way to mitigate these "disturb" conditions is a "refresh" operation on the FeFET. The refresh operation is designed to read and restore the programmed or erased state of the FeFET before its threshold voltage shifted too much. In some embodiments, an error detection and correction (ECC) operation may accompany the refresh operation, thereby allowing the correct data to be restored, even if the disturb condition has shifted the FeFET to an incorrect polarization state.

FIG. 1a-2 introduces additional randomly addressable units of memory cells that may be organized in an array of NOR memory strings: "columns," "pages" and "slices". For example, each column (e.g., column 115) represents the vertically-aligned FeFETs of multiple NOR memory strings in the active strips of an active stack; these vertically-aligned FeFETs share a common control gate or local word line.

Each "page" refers to a fixed number FeFETs (e.g., a fixed number of FeFETs arranged along the X-direction, indicated in FIG. 1a-2 as page 113), and each "slice" refers to an organization of contiguous FeFETs that extend along both the X- and Z-directions and one FeFET deep along the 5 Y-direction, indicated in FIG. 1a-2 by reference numeral 114

Although illustrated by conceptualized memory structure 100, which is shown in the figures as an array having a 4 planes by 4 active strips organization, a memory structure of 10 the present invention may have any number of FeFETs along any of the X-, Y- and Z-directions. For example, there may be 1, 2, 4, 8, 16, 32, 64 . . . planes in the Z direction, 2, 4, 8, 16, 32, 64, ... active strips along the X-direction, and each NOR memory string in each active strip may have 2, 4, 8, 15 16, . . . 8192 or more FeFETs formed along the Y-direction. The use of numbers that are integer powers of 2 (i.e., 2^n , where n is an integer) herein follows a customary practice in conventional memory design. It is customary to access each addressable unit of memory by decoding a binary address. 20 Thus, for example, a memory structure of the present invention may have M NOR memory strings along each of the X and Z directions, with M being a number that is not necessarily 2ⁿ, for any integer n. FeFETs of the present invention may be organized such that reading, programming 25 or erasing may be performed simultaneously on an individual page or individual slice basis. In some embodiments, an erase operation can also be performed in simultaneously for FeFETs in an even larger addressable unit of the array of NOR memory strings.

FIG. 1b shows a basic circuit representation of four NOR memory string-pairs, with each NOR memory string-pair including two NOR memory strings located on the same plane (i.e., one of planes 159-1, 159-2, 159-3 and 159-4), according to one embodiment of the present invention. 35 Within each NOR memory string-pair, corresponding FeFETs of the NOR memory strings share a common local word line (e.g., local word line 151n). The detailed structure of this configuration is discussed and illustrated below. FIG. 1b shows a configuration which includes four NOR memory string-pairs on four separate planes (e.g., NOR memory strings 150L and 150R on plane 159-4) that are provided in adjacent columns of conceptualized memory structure 100 and share common local word line 151a.

As shown in FIG. 1b, NOR memory strings 150L and 45 150R may be NOR memory strings formed along two active strips located on opposite sides of shared local word line **151***a*. FeFETs **152**R-**1** to **152**R-**4** and **152**L-**1** to **152**L-**4** may be FeFETs located in the four pairs of active strips on opposite sides of local word line 151a. In this embodiment, 50 local word line 151a controls FeFETs 152R-1, 152R-2, 152R-3 and 152R-4 from four NOR memory strings located on planes 159-1, 159-2, 159-3 and 159-4, respectively. Local word line 151a also controls FeFETs 152L-1, 152L-2, 152L-3 and 152L-4 from four NOR memory strings on 55 corresponding planes 159-1, 159-2, 159-3 and 159-4, respectively. As discussed in greater detail below, in some embodiments, the parasitic capacitance C intrinsic to each NOR memory string (e.g., the distributed capacitance between the common n⁺ source region of a NOR memory 60 string and its multiple associated local word lines) may be used as a virtual voltage source, under some operating conditions, to provide common source voltage V_{ss} . For that matter, the common source voltage can be set through a pre-charge operation that provides it the voltage of the 65 common n+ drain, for example. This flexibility in readily setting the source voltage and the drain voltage to a voltage

10

that is more positive than the voltage on the local word line (e.g., during an erase operation), or to a voltage that is more negative than the voltage on the local word line (e.g., during a programming operation) is particularly advantageous because the polarization states (e.g., the "erased" and "programmed" states) of the FeFET are primarily symmetrical. That is, a first polarization state of the FeFET can be set by imposing a first voltage on both the source and the drain regions—as a first electrode—and a second voltage on the local word line, serving as a second electrode. This first polarization state can be readily reversed (i.e., set to the other polarization state by exchanging the voltages on the first and the second electrodes (i.e., imposing (i) the second voltage on both the source and the drain regions, and (ii) the first voltage on the local word line). According to one embodiment of the present invention, the NOR memory string does not require a hard-wired connection to a ground potential (Vss) for its common source region, common drain region, or any of its channel regions. Consequently, both polarization states can be achieved on any of the FeFETs of the NOR memory string using only positive voltages (or only negative voltages) applied to the corresponding local word line and the common source and the common drain regions. This single voltage polarity capability is advantageous because the typical FeFET has polarization states characterized by a voltage hysteresis loop that between a negative voltage in the "programmed" state to a positive voltage in the "erased" state.

FIG. 1c shows a basic circuit representation of four NOR memory strings, with each NOR memory string being located on a respective one of planes 159-1, 159-2, 159-3 and 159-4, according to one embodiment of the present invention. In FIG. 1c, corresponding FeFETs (e.g., FeFETs 152-1, 152-2, 152-3 and 152-4) from the NOR memory strings share a common local word line (e.g., 151a). Each NOR memory string may run horizontally along the Y-direction, with storage elements (i.e., FeFETs) of the NOR memory string connected between common source line 153-m and common bit line 154-m, where in takes a value between 1 and 4. Corresponding FeFETs in the four NOR memory strings share corresponding one of local word lines **151**-*n*, where n is the index of a local word line. The FeFETs in the NOR memory strings of the present invention are variable threshold voltage thin-film ferroelectric or polarization transistors wherein current to voltage hysteresis is created when the FeFET is electrically programmed, erased, or read, using programming, inhibition, erasure and read voltages that may be readily generated in the "circuitry under array" ("CUA") formed in the semiconductor substrate.

Fabrication Processes

FIG. 2a shows a cross section in a Y-Z plane of semiconductor structure 200, after active layers 202-0 to 202-7 (each separated from the next active layer respectively by isolation layers 203-0 to 203-7) have been formed on semiconductor substrate 201, but prior to forming individual active strips, in accordance with one embodiment of the present invention. Semiconductor substrate 201 represents, for example, a p⁻ doped bulk silicon wafer on which support circuits for memory structure 200 may be formed as CUA prior to forming the active layers. Such support circuits, which may be formed alongside contacts 206-0 and 206-1 in FIG. 2c, may include both analog and digital circuits. Some examples of such support circuits include shift registers, latches, sense amplifiers, reference cells, power supply lines, bias and reference voltage generators, inverters, NAND, NOR, Exclusive-Or and other logic gates, input/output driv-

ers, address decoders (e.g., bit line and word line decoders), other memory elements, sequencers and state machines. These support circuits may be formed out of the building blocks for conventional devices (e.g., n-wells, p-wells, triple wells, n^+ , p^+ diffusions, isolation regions, low and high 5 voltage transistors, capacitors, resistors, vias, interconnects and conductors), as is known to those of ordinary skill in the

After the support circuits have been formed in and on semiconductor substrate **201**, an isolation layer is provided, 10 which may be a deposited or grown thick silicon oxide, for example.

Next, in some embodiments, one or more layers of interconnect conductors (e.g., metallic interconnect lines) may be formed, including "global word lines" for intercon- 15 necting the support circuits to the "local word lines" within the array of NOR memory strings. The layers of interconnect conductors are further discussed below. In other embodiments, global word lines are not provided under the memory array. Such metallic interconnect lines (e.g., global word line 20 landing pads 263 of FIG. 2c, discussed below) may be provided as horizontal long narrow conductive strips running along a predetermined direction that may be perpendicular to the NOR memory strings to be formed at a later step. To facilitate discussion in this detailed description, the 25 global word lines are presumed to run along the X-direction. The metallic interconnect lines may be formed by applying photo-lithographical patterning and etching steps on one or more deposited metal layers. (Alternatively, these metallic interconnect lines can be formed using a conventional 30 damascene process, such as a copper or tungsten damascene process). A thick oxide is deposited to form isolation layer, followed by a planarization step using conventional chemical mechanical polishing (CMP) techniques.

Over isolation layer 203-0, active layers 202-0 to 202-7 35 are then successively formed, with each active layer being electrically insulated from the previous active layer underneath by a corresponding one of isolation layers 203-1 to 203-7. In this detailed description, the term "active layer" refers to material layers that form circuit elements (e.g., 40 thin-film ferroelectric or polarization transistors) in the array of NOR memory strings. In FIG. 2a, although eight active layers are shown, any number of active layers may be provided. In practice, the number of active layers to be formed may depend on the process technology, such as the 45 availability of a well-controlled anisotropic etching process that allows cutting through a tall stack of active layers to reach semiconductor substrate 201.

FIGS. 2b-1 to 2b-5 show examples of active layer 220a, according to one embodiment of the present invention. In the 50 example of FIG. 2b-1, active layer 220a includes the material layers: (i) n⁺ polysilicon or silicon layers 204d and 204b and (ii) sacrificial dielectric layer 204c (which is to be subsequently replaced at least in part by a p-doped channel polysilicon layer). In one implementation, n⁺ silicon layers 55 204d-204b may be deposited successively in the same process chamber without removal in between. N⁺ silicon layer 204b may be formed by depositing 10-100 nm of in-situ doped n⁺ polysilicon. Dielectric layer 204c may be formed by depositing a sacrificial material (e.g., silicon 60 oxide or SiO₂) in the thickness range of 10-100 nm. N⁺ silicon layer 204d (i.e., the top portion of the deposited polysilicon) may be formed by depositing 10-100 nm of in-situ doped n⁺ polysilicon. Both silicon layers 204b and 204d may be in situ n⁺-doped. N⁺ dopant concentrations in 65 silicon layers 204d and 204b may be, for example between 1×10^{20} /cm³ and 1×10^{21} /cm³. One consideration for the

dopant concentration is the sheet resistivity in n+ silicon layers 204d and 204b. The n⁺ doping may be achieved by either (i) a low-energy shallow high-dose ion implantation of phosphorus, arsenic or antimony, or (ii) in-situ phosphorus or arsenic doping of the deposited polysilicon, forming a 10-100 nm thick n^+ silicon layer **204***d* on top. The boron or p-dopant concentration of the channel polysilicon layer—replacing at least in part sacrificial dielectric layer **204**c—can be in the range of 1×10^{16} /cm³ to 1×10^{18} /cm³. The actual boron concentration in layer 204c determines the native transistor threshold voltage, channel mobility, n⁺p⁻n⁺ punch-through voltage, n+p- junction leakage and reverse diode conduction characteristics, and channel depletion depth under the various operating conditions for the n+p-n+ FeFETs formed along active strips formed out of active layers 202-0 to 202-7.

12

Thermal activation of the n⁺ and p⁻ implanted species and recrystallization of silicon layers 204d, 204c and 204b should preferably take place all at once after all active layers 202-0 to 202-7 have been formed, using a conventional rapid thermal annealing technique (e.g., at 700° C. or higher) or a conventional rapid laser annealing technique, thereby ensuring that all active layers experience elevated temperature processing in roughly the same amount.

Alternatively, n^+ and p^- dopants of each of active layers 202-0 to 202-7 can be activated individually by shallow rapid thermal annealing using, for example, an excimer laser annealing (ELA) technique at an ultraviolet wavelength (e.g., 308 nanometer). The annealing energy which is absorbed by the polysilicon or amorphous silicon to partially melt layer 204d and part or all of layer 204c, optionally penetrating layer 204b to affect volume 205 (see FIG. 2b-4) without unduly heating other active layers lying below portion 223 of n^+ silicon layer 204b of the annealed active layer 220a.

Although the use of successive layer-by-layer excimer laser shallow rapid thermal anneal is more costly than a single deep rapid thermal anneal step, ELA has the advantage that the localized partial melting of polysilicon (or amorphous silicon) can result in recrystallization of annealed volume 205 to form larger silicon polycrystalline grains having substantially improved mobility and uniformity, and reduced FeFET leakage due to reduced segregation of n+ dopants at the grain boundaries of the affected volume. Such shallow excimer laser low-temperature anneal technique is well-known to those of ordinary skill in the art. For example, such technique is used to form polysilicon or amorphous silicon films in solar cell and flat panel display applications. See, for example, the article "Comprehensive Study of Lateral Grain Growth in Poly-Si Films by Excimer Laser Annealing (ELA) and its applications to Thin Film Transistors" by H. Kuriyama et al., published in Japanese Journal of Applied Physics, Vol. 33, Part 1, Number 10, 20 Aug. 1994, or the article "Annealing of Silicon Backplanes with 540W Excimer Lasers", published in technical publication by Coherent Inc. on their website.

The thickness of layer 204c roughly corresponds to the channel length of the FeFETs to be formed, which may be 10 nm thick or less over long active strips. In one embodiment, for example, diffusion in active layer 220a of FIG. 2b-5 may be controlled so that the channel length of the FeFETs remains near the target of 20.0-100.0 nm, even after several thermal process cycles, by depositing an ultra-thin (from one or a few atomic layers thick) film of silicon nitride (e.g., SiN or Si_3N_4), or another suitable diffusion-blocking film following the formation of n^+ silicon layer 204b (e.g., layer 204b-d in FIG. 2b-5). A second ultra-thin film of

silicon nitride, or another suitable diffusion-blocking film (e.g., layer 204d-d in FIG. 2b-5), may optionally be deposited following deposition of p doped layer 204c, before depositing n⁺ silicon layer 204d. The ultra-thin dopant diffusion-blocking layers 204d-d and 204b-d can be depos- 5 ited by chemical vapor deposition, atomic layer deposition or any other suitable means (e.g., high pressure nitridation at low temperature). Each ultra-thin dopant diffusion-blocking layer acts as a barrier that prevents the n+ dopants in n+ silicon layers 204d and 204b from diffusing into p⁻ layer 10 204c, yet are sufficiently thin to only marginally impede the MOS transistor action in the channel region between n+ silicon layer 204d (acting as a source) and n⁺ silicon layer **204***b* (acting as a drain). These additional ultra-thin dopant diffusion-blocking layers increase the manufacturing cost, 15 but such ultra-thin dopant diffusion-blocking layers may serve to significantly reduce the cumulative leakage current from the multiple FeFETs along the active strips that are in the "off" state. However, if that leakage current is tolerable then these ultra-thin layers can be omitted. Indeed, relative 20 to conventional non-ferroelectric field-effect transistors, the "off" state leakage currents in FeFETs are significantly smaller because of the large bandgap in the ferroelectric storage material (e.g., HfO₂).

Sheet resistance in n⁺ silicon layers 204b and n⁺ silicon 25 layers 204d may be reduced by adding optional conductive or metallic layers 204g or 204f to each active layer, adjacent to n^+ silicon layers 204d or 204b (e.g., metal layer 204a, labeled as W (for tungsten) in FIGS. 2b-2 and 2b-3). Metal layer 204g or 204f may be provided by one or more 30 deposited metal layers. For example, metal layer 204f or 204g may be provided by depositing 1-2 nm thick liner layer of TiN followed by depositing a 1-40 nm thick layer of tungsten, a similar refractory metal, or a polycide or silicide (e.g., nickel silicide). Metal layer 204f or 204g is more 35 preferably in the 10.0-40.0 nm thickness range. Preferably, a sacrificial material (e.g., silicon nitride or SiN) is deposited to stand-in for metal layers 204f and 204g, which is subsequently replaced by a conductive material to form lowresistivity layers 204f and 204g. Even very thin layer 204f 40 and 204g (e.g., 20.0-50.0 nm) can significantly reduce the line resistance of a long active strip, while allowing the use of less heavily doped n⁺ silicon layers 204d and 204b.

As shown in FIG. 2c, the conductor inside contact opening 205-1 can become quite long for a tall stack, which may 45 adversely increase the line resistance. In that case, metal layer 204a (e.g., a tungsten layer) may preferably be included below n⁺ silicon layer 204b, so as to substantially fill contact opening 205-1, rather than placing it above n⁺ silicon layer 204d, as is shown in FIG. 2c.

In the embodiments where metal layers 204f and 204g are not present, there are several tradeoffs that can be made: for example, longer active strips are possible if the resulting increased read latency is acceptable. In general, the shorter the active strip, the lower the line resistance and therefore 55 the shorter the latency. (The trade-off is in array efficiency). In the absence of metallic layers 204f and 204g, the thickness of n⁺ silicon layers 204d and 204b can be increased (for example, to 100.0 nanometers) to reduce the intrinsic line resistance, at the expense of a taller stack to etch through. 60 The line resistance can be further reduced by increasing the n⁺ doping concentration in n⁺ silicon layers 204d and 204b and by applying higher anneal temperatures in excess of 1,000° C. (e.g., by rapid thermal anneal, deep laser annealing or shallow excimer laser annealing techniques) to 65 enhance recrystallization and dopant activation and to reduce dopant segregation at the grain-boundaries.

14

Shorter active strips also have superior immunity to leakage between n⁺ silicon layer 204b and n⁺ silicon layer 204d. A thicker n⁺ layer provides reduced strip line resistance and increased strip parasitic capacitance, which is not desirable for dynamic sensing. The integrated circuit designer may opt for a shorter active strip (with or without metal layer 224) when low read latency is most valued. Alternatively, the strip line resistance may be reduced by contacting both ends of each active strip, rather than just at one end

Block-formation patterning and etching steps define separate blocks in each of the active layers formed. Each block occupies an area in which a large number (e.g., thousands) of active strips running in parallel may be formed, as discussed below, with each active strip running along the Y-direction, eventually forming one or more NOR memory strings that provide a large number (e.g., thousands) of individually randomly accessible FeFETs.

Each of active layers 202-0 to 202-7 may be successively formed by repeating the steps described above. In addition, in the block-formation patterning and etching steps discussed above, each next higher active layer may be formed with an extension slightly beyond the previous active layer (see, e.g., as illustrated in FIG. 2c, discussed below, layer 202-1 extends beyond layer 202-0) to allow the upper active layer to access its specific decoders and other circuitry in semiconductor substrate 201 through designated buried contacts.

As shown in FIG. 2c, buried contacts 205-0 and 205-1 connect contacts 206-0 and 206-1 in semiconductor substrate 201, for example, to the local bit lines or source lines formed out of n⁺ silicon layer 204b in each of active layers 202-0 and 202-1. Buried contacts for active layers 202-2 to 202-7 (not shown) may be similarly provided to connect active layers 202-2 to 202-7 to contacts 206-2 to 206-7 in semiconductor substrate 201 in an inverted staircase-like structure in which the active layer closest to the substrate has the shortest buried contact, while the active layer furthest from the substrate has the longest buried contact. Alternatively, in lieu of buried contacts, conductor-filled vias extending from the top of the active layers may be etched through isolation layers 203-0 and 203-1. These vias establish electrical contact from substrate circuitry 206-0, for example, to top n⁺ silicon layers 204d-0 (or metal layer 224, if provided). The vias may be laid out in a "staircase" pattern with the active layer closest the substrate connected by the longest via, and the active layer closest to the top connected by the shortest via. The vias (not shown) have the advantage that more than one plane can be contacted in one maskingand-etch step, as is well-known to a person of ordinary skill in the art.

In an integrated circuit incorporating multiple arrays of NOR memory strings, the arrays of NOR memory strings may each be provided as a modular circuit or tile, as discussed above. Each array of NOR memory strings may include an "array" portion, in which the FeFETs of the NOR memory strings are provided, and one or more "staircase" portions for routing signals into or out of the NOR memory strings, using either the buried contact inverted staircase-like structure, or by conductor-filled vias in a staircase structure, as known to those of ordinary skill in the art. Each tile may have, for example, a regular-shape footprint (e.g., a rectangular footprint) on the planar surface of the semiconductor substrate, thus allowing many tiles to be formed and arranged on a substrate in a regular manner (e.g., a m×n matrix, where m and n are integers.)

Through a switch circuit, each of contacts 206-0 to 206-7 may apply a pre-charge voltage V_{bl} to the respective bit line or source line of the corresponding NOR memory strings. During a read operation, each selected bit line may be connected by the switch circuit to an input terminal of a 5 sense amplifier or a latch. The switch circuit may selectively connect each of contacts 206-0 to 206-7 to any of a number of specific voltage sources, such as a programming voltage (V_{pgm}) , inhibit voltage $(V_{inhibit})$, erase voltage (V_{erase}) , or any other suitable predetermined or pre-charge reference voltage V_{bl} or V_{ss} . In some embodiments, discussed below, taking advantage of the relatively large parasitic distributed capacitance along a bit line or source line in an active strip, a virtual voltage reference (e.g., a virtual ground, providing ground voltage V_{ss}) may be created in the source line (i.e., 15 n⁺ silicon layer **204***d*, or a metal line) of each active strip by pre-charging the source line, as discussed below. The virtual ground approach eliminates the need for hard-wiring n+ silicon layer 204d to a voltage source in the substrate, making it possible to more efficiently use the staircase via 20 structure to connect each active strip from the top to the substrate.

FIG. 2c also shows buried contacts 261-0 to 261-n for connecting global word lines 208g-s—which are to be formed running along the X-direction—to contacts 262-0 to 25 262-n in semiconductor substrate 201. Global word lines 208g-s are provided to connect corresponding local word lines 208W-s yet to be formed to circuits 262-n in substrate 201. Landing pads 263 are provided on the global word lines to allow connection to local word lines 208W-s, which are 30 yet to be formed vertically on top of horizontally running global word lines 208g-s. Through a switch circuit and a global word line decoder, each of global word line contacts 262-0 to 262-n may be selectively connected, either individually, or shared among several global word lines, to any 35 one of a number of reference voltage sources, such as stepped programming voltages (V_{program}), program-inhibit voltage (V_{inhibit}), read voltages (V_{read}) and erasure voltages (V

(V_{erase}).

The buried contacts, the global word lines and the landing 40 pads may be formed using conventional photo-lithographical patterning and etching steps, followed by deposition of one or more suitable conductors or by alloying (e.g., tungsten metal, alloy or tungsten silicide).

FIGS. 2d to 2k illustrates a portion of an exemplary 45 process for forming a modular circuit or tile that includes an array of NOR memory strings, including the array portion (i.e., memory structure 200), according to one embodiment of the present invention; FIG. 2k illustrates a step the exemplary process that provides a channel layer for the 50 FeFETs. The immediate steps following the step providing the channel layer in the exemplary process are illustrated by three groups of process steps, each group representing a different approach to "gate stack" formation. The groups of process steps are illustrated, respectively, in FIGS. 3a to 3b, 55 4a to 4c and 5a to 5f.

FIG. 2d(i) shows a cross section in the Z-X plane of memory structure 200 after depositions of active layers 202-0 to 202-7, in accordance with one embodiment of the present invention. Initially, in this embodiment, pad oxide 60 240 (e.g., silicon oxide) is provided over planar surface 201 of the semiconductor substrate. Etch-stop layer 241 (e.g., tungsten (W), tungsten nitride (WN), aluminum oxide (AlO), aluminum nitride (AlN), or any suitable etch-stop material) is then provided. Silicon oxycarbide (SiOC) layer 65 203-0 is then provided to isolate etch-stop layer 241 from the next active layer (i.e., active layer 202-0) to be deposited.

16

Thereafter, active layers 202-0 to 202-7 (eight in total in this example; generically, each may be referred to as "active layer 202") are successively deposited, each active layer being isolated from the next active layer by a corresponding one of isolation layers 203-1 to 203-7 (active layers 203-0 to 203-7 are also each generically referred to as "isolation layer 203" and, collectively, "isolation layers 203"). Active layers 202 each include, in order of deposition, (i) silicon nitride (SiN) layer 204a, (ii) n*-doped amorphous silicon (or polysilicon) layer 204b, (iii) sacrificial oxide layer 204c, (iv) n*-doped amorphous silicon (or polysilicon) layer 204e. Isolation layer 245 (e.g., an additional SiOC layer) is then deposited on top of active layer 202-7. The resulting structure is shown as memory structure 200 of FIG. 2d(i).

FIG. 2d(ii) illustrates successive recessing and etching steps to create staircase portion 200-sc, in accordance with one embodiment of the present invention. As shown in FIG. 2d(ii), on the top surface of memory structure 200 (i.e., isolation layer 245) is formed mask layer 210, which exposes first portion 211 of memory structure 200, as shown in FIG. 2d(ii)(1). The exposed portion of isolation SiOC layer 245 is then removed to expose a portion of active layer 202-7 underneath. That exposed portion of active layer 202-7 is then removed, exposing a portion of SiOC layer **203-6** underneath. The resulting structure is shown in FIG. 2d(ii)(2). Mask layer 210 is then recessed to expose a further portion of isolation SiOC layer 245. Removal of the exposed SiOC layers (e.g., SiOC layers 245 and 203-7), removal of exposed active layer 202 and recessing mask layer 210 are then repeated 7 more times, thereby forming a staircase structure. Thereafter, an oxide is deposited to fill the portions of active layers 204-1 to 204-7 removed. A chemicalmechanical polishing (CMP) step is carried out to remove mask layer 210 and to planarize the top surface of memory structure 200.

The steps illustrated in FIG. 2d form the staircase structures (e.g., on opposite sides of the array portion) prior to processing the array portion. However, the array portion may also be processed prior to formation of the staircase structures.

The surface of SiOC layer 245 is patterned and etched to create openings 218. FIG. 2e(i) is a top view of memory structure 200, showing shafts 218, after patterning and etching. In FIG. 2e(i), shafts 218 are shown to have a circular X-Y plane cross section. However, any suitable X-Y plane cross section shape (e.g., oval) may also be used. A series of etching steps remove the exposed portions of isolation layers 203-0 to 203-7, and active layers 202-0 to 202-7 until etch-stop layer 241 is reached, thereby extending shafts 218 along the Z-direction. FIG. 2e(ii) shows an X-Z plane cross-section of memory structure 200 along line A-A' in FIG. 2e(i), after the etching steps through active layers 202-0 to 202-7 and isolation layers 203-0 to 203-7, in accordance with one embodiment the present invention. (In this detailed description, a "column" of objects denotes objects aligned along the Y-direction, whereas a "row" of objects denotes objects aligned along the X-direction.) In FIG. 2e(i), shafts of each column is staggered relative to the shafts of each of its immediate neighbor column, such that the closest shafts in adjacent columns have a greater separation between them than if such shafts are aligned along the X-direction. In one embodiment adjacent columns of shafts are 160 nm apart along the X-direction, and each of shafts 218 may have a diameter of 100.0 nm. Shafts 218 may each be up to 2.0 microns deep, so that the etching steps has an aspect ratio that is less than 20.

Shafts 218 may then be lined using an oxide liner 242 and filled using sacrificial material 243 (e.g., carbon or tungsten (W) with an adhesive layer of titanium nitride (TiN)). The surface of memory structure 200 may then be planarized using CMP or, when carbon is used as the sacrificial mate- 5 rial, etched back. After shafts 218 are filled with sacrificial material 243, hard mask 244 (e.g., a carbon hard mask layer) is then deposited and photo-lithographically patterned over memory structure 200. Hard mask 244 may be patterned to provide rectilinear openings 246, each extending lengthwise 10 along the Y-direction, exposing portions of memory structure 200 (e.g., each 60.0 nm wide, in one embodiment), and overlapping but not removing the columns of filled shaft 218. A top view and a cross-section view of resulting memory structure 200, after hard mask 244 is patterned, are 15 shown in FIGS. 2f(i) and 2f(ii), respectively, according to one embodiment of the present invention. The cross-section view of FIG. 2f(ii) cuts through the X-Z plane along line A-A' of FIG. 2f(i).

From underneath rectilinear openings **246**, successive 20 etching steps remove portions of SiOC layer 245, isolation layers 203-0 to 203-7, active layers 202-0 to 202-7, and etch-stop layer 241, thereby dividing memory structure 200 into numerous active stacks separated by trenches 267. The etching steps are designed to be selective to oxide liner 242 25 and sacrificial material 243 in filled shafts 218, so that filled shafts 218 become pillars 248 of sacrificial material 243, surrounded on its cylindrical surface by oxide liner layer 242. An additional wet etch may be applied to remove any stringers of the various material. Hard mask 244 may then 30 be removed (e.g., by CMP or by ashing, as appropriate). A top view and an X-Z plane cross section view along line A-A' of resulting memory structure 200, after hard mask 244 has been removed, are shown in FIGS. 2g(i) and 2g(ii), respectively, according to one embodiment of the present 35

Trenches 267 may then be filled by dielectric material 247 (e.g., silicon oxide) using, for example, atomic layer deposition (ALD) or a spin-on technique. Excess oxide on the top surface of memory structure 200 may then be removed 40 using, for example, CMP. Thereafter, pillars 248 are removed using, for example, a suitable wet etch to restore shafts 218. FIGS. 2h(i) and 2h(ii) show a top view and an X-Z plane cross-section view along line A-A' of resulting memory structure 200, respectively, after removal of pillars 45 248, according to one embodiment of the present invention.

Removal of pillars 248 restores shafts 218 between filled portions of trenches 267 and exposes active layers 202-0 to 202-7 as sidewalls of shafts 218. Through shafts 218, SiN layers 204a and 204e of each of active layers 202-0 to 202-7 50 may be removed using, for example, a wet etch of silicon nitride. FIGS. 2i(i) and 2i(ii) show a top view and an X-Z plane cross-sectional view along line A-A' of resulting memory structure 200, respectively, after removal of SiN to 202-7, according to one embodiment of the present

An ALD step then deposits a conductive material into the cavities resulting from removal of SiN layers 204a and 204e in each of active layers 202-0 to 202-7. The conductive 60 material may include, for example, a metal liner (e.g., titanium (Ti), titanium nitride (TiN), tantalum (Ta) or tantalum nitride (TaN)), followed by a refractory metal (e.g., tungsten (W), tungsten nitride (WN) or molybdenum (Mo)). Thereafter, an anisotropic separation etch removes the con- 65 ductor material from shafts 218. Conductor layers 204f and 204g that respectively replace SiN layers 204a and 204e in

18

each active layer may be recessed, for example, by 10.0 nm, to allow a subsequent filler material (e.g., silicon nitride) to prevent corrosion in the conductive material. FIGS. 2j(i) and 2i(ii) show a top view and an X-Z plane cross-sectional view along line A-A' of resulting memory structure 200, respectively, after replacement of SiN layers 204a and 204e in each active layer 202 by conductor layers 204f and 204g, according to one embodiment of the present invention. Conductor layers 204f and 204g are optional conductor layers that are formed adjacent and in contact with n⁺-doped silicon layers 204b and 204d. As n⁺ silicon layers 204b and 204d of each active layer 202 would become the common bit line and the common source line of a NOR memory string to be formed, conductor layers 204f and 204g desirably reduce the resistivities in the common bit line and the common source line.

Then, sacrificial oxide layer 204c, n⁺ silicon layer 204b, and n⁺ silicon layer 204d of each active layer 202 may be recessed (e.g., by 20.0 nm) to create a cavity into which a channel semiconductor material may be deposited. Polysilicon layer 250, intended for providing the channel regions of the thin-film storage transistors, is then deposited, etched back for separation, and recessed to remove stringers. Channel polysilicon layer 250 may be, for example, p-doped poly silicon (e.g., 7.0-14.0 nm thick, preferably about 10.0 nm thick). In some embodiments, the final thickness is preferably the depth of the cavity created in active layer 204 (e.g., 20.0 nm). FIGS. 2k(i) and 2k(ii) show a top view and an X-Z plane cross-sectional view along line A-A' of resulting memory structure 200, respectively, after channel polysilicon layer 250 is recessed, according to one embodiment of the present invention.

The next steps in the exemplary process form the "gate stacks" for FeFETs in the NOR memory strings. The gate stack for an FeFET includes (i) optionally, an interface layer, (ii) a ferroelectric or polarization layer, and (iii) a gate electrode. Optionally, a thin metal oxide (e.g., 1.0-2.0 nm thick aluminum oxide) may also be provided between the gate electrode and the polarization layer to reduce tunneling of charged species into the polarization layer. The inventors discover that superior performance in 3-dimensional arrays of NOR memory strings may be achieved when the polarization layers in the FeFETs are separated from each other. FIGS. 3a to 3b, 4a-4c and 5a-5f illustrate three exemplary, alternative groups of process steps suitable for fabricating the gate stacks of FeFETs in a NOR memory string, in accordance with some embodiments of the present invention. To provide greater clarity, each of FIGS. 3a to 3b, 4a to 4c and 5a to 5f shows only active stacks 261a and 261band only active strips formed out of active layers 202-0 to 202-3. However, it is understood that the teachings provided therein apply to all active stacks and active strips in any array of NOR memory strings in that embodiment of the present invention.

FIGS. 3a to 3b illustrate a first group of process steps for layers 204a and 204e from each of active multi-layers 202-0 55 forming the gate stacks in FeFETs of an array of NOR memory strings, according to one embodiment of the present invention. FIGS. 3a to 3b each provide (a) an X-Y plane view across isolation layer 203-4 between the active strips that are formed out of active layers 202-3 and 202-4 in FIG. 2k(ii), and (b) an X-Z plane cross-sectional view, each showing active stacks 261a and 261b of memory structure **200** after completing the respective steps in this first group. FIG. 3b also includes a third view (iii), which is an X-Y plane cross-sectional view along line C-C' of FIG. 3b(ii).

> After channel polysilicon layer 250 is conformally deposited, an etch-back step is performed, as discussed above in conjunction with FIGS. 2k(i) and 2k(ii). The etch-back step

also exposes isolation layers 203-0 to 203-7 on the sidewalls of shafts 218, as illustrated on in FIG. 3a(ii). In some embodiments, the etch-back step may also be designed to form a recess in active layers 202 between exposed portions of adjacent isolation layers 203. Thereafter, a thin silicon 5 nitride layer (e.g., 1 nm thick) may be formed on the exposed surfaces of channel polysilicon layer 250 in shafts 218. The silicon nitride layer may be formed using an ammonia nitridation step on a native oxide on the surface of polysilicon layer 250. (Alternatively, a thin oxide layer may be formed on polysilicon layer 250 by an oxidation step at a temperature of 400-650° C., and then converted to silicon nitride by any suitable nitridation process.) This nitride layer is used as an interface dielectric layer between channel polysilicon layer 250 and a polarization or polarization layer 15 to be formed. Other suitable materials for the interface dielectric layer includes zirconium oxide (ZrO₂), silicon oxynitride (SiON), silicon nitride (Si₃N₄), or silicon oxide

The native oxide layer—from which the interface dielec- 20 tric layer is formed—may be formed by a chemical clean of the surface of polysilicon channel material 250. (The silicon nitride in interface dielectric layer may also be any combination of silicon oxynitride (SiON) or silicon nitride (Si_3N_4) , and silicon oxide (SiO_2) , 0.0 to 2.0 nm thick, with 25 a refractive index between 1.5 to 2.0.) A densification treatment may be performed on the silicon nitride layer using, for example, pulsed ozone or by thermal annealing in a hydrogen or deuterium ambient, or any other techniques known to a person of ordinary skill in the art. The densification treatment reduces electronic leakage through the interface dielectric layer, reduces the surface states at the interface between the third semiconductor layer and the polarization layer, or both. The inventors believe that the interface dielectric layer should have a refractive index 35 between 1.4 and 1.8, preferably 1.6, which may be achieved, for example, by annealing at 630° C. in ozone or deuterium

Selective deposition techniques may be used to deposit polarization layer **264** on the exposed surfaces, other than on 40 isolation layers **203-0** to **203-7**. Suitable techniques include "surface poisoning" and self-assembled monolayers ("SAMs") techniques. In the SAM technique, for example, the exposed surfaces are treated with reactive species having active hydroxyl (—OH) bonds. Resulting memory structure 45 **200** is shown in FIGS. 3a(i). In this embodiment, as the interface dielectric layer is deposited in a recess in each active strip between adjacent isolation layers, the interface dielectric layer is not seen in the X-Y plane cross-sectional view of FIG. 3a(ii), which cuts through isolation layer 50 **203-4**.

Polarization layer **264** is then selectively deposited on the surfaces of the silicon nitride layer and channel material 250. (e.g., the treatment by SAMs prevents deposition of polarization layer 250 on the sidewalls of isolation dielectric 55 layers 203.) Resulting memory structure 200 is shown in FIG. 3b(i). In this embodiment, the X-Y plane cross-sectional view of FIG. 3b(iii) shows polarization layer 264 deposited in the recessed portion of active layer 203-3, between adjacent exposed portions of isolation layers 203-3 60 and 203-4. Polarization layer 264 may be formed by the selective deposition using ALD techniques in an ozone ambient at temperature between 200° C. to 330° C. (e.g., around 300° C.), with a post-deposition annealing step at a temperature between 400° C. and 1000° C., based on the 65 desired crystallization phase requirement of the ferroelectric material.

20

Deposition of the ferroelectric material may be achieved, for example, using any suitable deposition technique (e.g., chemical or physical vapor deposition, ALD, or evaporation) repeated cycles of hafnium oxide depositions and zirconium oxide depositions (e.g., in a HfO₂:ZrO₂ ratio of 4:1). Polarization layer 264 may be, for example, 2.0 to 8.0 nm thick, preferably, 2.5-5.0 nm thick). For the thicker polarization layers (e.g., greater than 4.0 nm), additional SAM treatments between deposition cycles may be advisable. Post deposition annealing may be performed, for example, in a temperature of 400-1000° C. to achieve a desired crystallization phase. Other suitable ferroelectric materials include aluminum zirconium-doped Hafnium oxide (HfZrAlO), aluminum-doped hafnium oxide (HfO2:Al), silicon-doped hafnium oxide (HfO₂:Si) or lanthanum-doped hafnium oxide (HfO₂:La). The term HZO may encompass hafnium zirconium oxide (HfZrO), hafnium zirconium oxynitride (HfZrON), hafnium zirconium aluminum oxide (HfZrAlO) or any hafnium oxide that includes zirconium impurities.

The "surface poisoning" technique may be used to provide both the interface dielectric layer and polarization layer 264. In one embodiment, HZO is deposited directly onto polysilicon channel layer 250 after its formation. After a few atomic layers of HZO have been deposited, a layer of silicon oxide is formed as native oxide at the interface between channel polysilicon layer 250 and the atomic layers of HZO. Nitridation of the silicon oxide layer can proceed in the manner discussed above to form the interface dielectric layer, which is silicon nitride in this case. The remainder of polarization layer 264 may then proceed to completion.

The interface dielectric layer significantly reduces tunneling of charged species (i.e., electrons or holes) between channel polysilicon layer 250 and polarization layer 264. The charged species that tunnel into polarization layer 264 may adversely affect its polarization. Over time, the interface dielectric layer may break down as a result of a high number of program/erase cycles, thereby reducing the endurance of the FeFET. The interfacial layer may be formed out of a material with a dielectric constant greater than that of silicon oxide ("high-k" material; preferably, greater than 3.9). For example, aluminum oxide is a high-k material, having a dielectric constant as high as 9.0. A high-k material reduces both the electric field during program or erase operations and tunneling from the channel region. For a 0.0 nm thick interface dielectric layer, the polarization layer may be directly deposited by ALD onto the channel region (e.g., polysilicon). A native oxide of a self-limiting thickness (e.g., 1.0 to 10.0 angstrom) would be inherently formed at the interface between the channel region and the polarization layer. While this native oxide may be of low quality or prone to create a leakage current, it can be substantially improved by a self-limiting nitridation step. This approach is particularly advantageous when the channel region is formed after high temperature steps, such that contamination by dopant diffusion is a lesser concern. In some embodiments, a bandgap-engineered tunneling layer (e.g., a silicon oxide (SiO₂) and zirconium oxide (ZrO₂) multi-layer) may serve as the interface dielectric layer, providing the advantage of reduced tunneling into the polarization layer. In some embodiments, aluminum oxide (e.g., Al₂O₃) or silicon oxide nitride may be suitable. The high-k dielectric properties of the zirconium oxide and other high-k material reduce the electric field in the interface dielectric layer.

In other embodiments, polysilicon channel **250** may be replaced by a junction-less oxide semiconductor (e.g., indium gallium zinc oxide ("IGZO"), and indium aluminum

zinc oxide ("IAIZnO")). With such a channel material, the interfacial layer between it and polarization layer **264** (e.g., HZO) may be subdued or even non-existent. The inventors believe that the oxide semiconductor channel region improves the endurance of the FeFET.

Thereafter, conductive material 272 (i.e., gate electrode) is deposited to fill shafts 218. Conductive material 272 may include tungsten (W), molybdenum (Mo), aluminum (Al), ruthenium (Ru), tantalum (Ta), titanium (Ti), or any combination or alloy of these metals.

FIGS. 4a to 4c illustrate a second process flow for forming the gate stacks of the FeFETs in an array of NOR memory strings, according to one embodiment of the present invention. FIG. 4a shows an X-Z plane cross-sectional view of active stacks 261a and 261b of memory structure 200, 15 after the interface layer (e.g., a 1.0 nm thick silicon nitride layer) has been formed, for example, in substantially the same manner as described above in conjunction with FIG. 3a. Relative to memory structure 200 of FIG. 3a, however, the recesses in the active strips that are formed out of active 20 layers 202-0 to 202-3 in FIG. 4a are deeper. For example, for channel polysilicon layer 250 to have an ultimate thickness of 10.0 nm, the recesses in the active strips that are formed out of active layers 202-0 to 202-3 are made to be 20.0 nm deep. An etch-back step applied to polysilicon channel layer 25 250 (e.g., a wet etch) may reduce the thickness of channel polysilicon layer 250 to 10.0 nm.

Thereafter, polarization layer 264 may be formed by depositing using, for example, ALD techniques. Specifically, polarization layer 264 may be deposited on the interface layer in the recesses of active strips formed out of active layers 202-0 to 202-3. As discussed above, the polarization layer may be formed, for example, using repeated cycles of hafnium oxide depositions and zirconium oxide depositions (e.g., in a HfO₂:ZrO₂ ratio of 4:1). Resulting memory 35 structure 200 is shown in FIG. 4b. Other suitable ferroelectric materials that may be used include aluminum-doped hafnium oxide (HfAlO) and aluminum zirconium-doped hafnium oxide (HfZrAlO). As shown in FIG. 4b, because of the recesses, polarization layer 264 in the active strips are 40 separated from each other by the adjacent isolation dielectric layers 203. Thereafter, conductive material 272 is deposited to fill shafts 218, forming thereby the gate electrodes or local word lines.

FIGS. 5a to 5f illustrate a third group of process steps for 45 forming the gate stacks in FeFETs of an array of NOR memory strings, according to one embodiment of the present invention. FIG. 5a shows an X-Z plane cross-sectional view of active stacks 261a and 261b of memory structure 200, after deposition of channel polysilicon layer 250 into deep 50 recesses of the active strips that are formed out of active layers 202-0 to 202-3 (e.g., 20 nm deep), an etch-back step and the formation of an interface dielectric layer in preparation for the ferroelectric material deposition. The etchback step thins channel polysilicon layer 250 (e.g., to 10.0 55 nm thick) to expose the sidewalls of isolation dielectric layers 203-0 to 203-4. Memory structures 200 of FIG. 5a may be formed using substantially the same process steps as those discussed above that are used to form memory structure 200 of FIG. 4a.

Thereafter, polarization layer **264** is formed on memory structure **200**, using the ALD techniques discussed, for example, in conjunction with FIG. **4***b* above, followed by a conformal deposition of amorphous silicon liner **265**. Resulting memory structure **200** is shown in FIG. **5***b*. Then, 65 an anisotropic dry etch step removes portions of amorphous silicon liner **265**, which exposes portions of polarization

layer 264 on the sidewalls of isolation dielectric layers 203 in shaft 218, while allowing the remainder of amorphous silicon liner 265 protect the portions of polarization layer 264 in the recesses of the active strips. The anisotropic dry etch also sputters away amorphous silicon liner 265 and polarization layer 264 from top of memory structure 200.

Resulting memory structure 200 is shown in FIG. 5c.

22

Thereafter, in a first variation, a wet etch step removes the ferroelectric material (e.g., hafnium zirconium oxide (HZO)) of polarization layer 264 from the sidewalls of isolation dielectric layers 203 in shafts 218. In a second variation, the wet etch step is allowed to remove a portion of the ferroelectric material of polarization layer 264 underneath amorphous liner layer 265. The inventors believe that reducing the extent of polarization layer 264 along the Z-direction allow greater control of the resulting polarization properties. In addition, the interface dielectric layer for the diminished polarization layer 264 results in a reduced electric field between the source region and the drain region of an FeFET. Resulting memory structure 200 for the first variation and resulting memory structure 200 for the second variation are respectively shown in FIGS. 5d(i) and 5d(ii). Thereafter, the remainder of amorphous silicon liner 265 may be removed by a wet etch. In the second variation, oxide 266 (e.g., silicon oxide) is deposited using, for example, ALD techniques to fill the recesses in active layers 202. An anisotropic etch-back step exposes diminished polarization layer 264 on the sidewalls of shafts 218. Resulting memory structure 200 for the first variation and resulting memory structure 200 for the second variation are shown in FIGS. 5e(i) and 5e(ii), respectively.

Conductive material 272 is then deposited into and fills shafts 218 (e.g., a titanium nitride (TiN) adhesion layer, followed by tungsten). Conductive material 272 provides a common gate electrode—i.e., the local word line—for vertically aligned FeFETs in the NOR memory strings of an active stack. Resulting memory structure 200 for the first and the second variations are shown in FIGS. 5f(i) and 5f(ii), respectively. In summary, memory structure 200 of FIG. 6a represents active stacks 261a and 261b fabricated under the first group of process steps, the second group of process steps, or the first variation of the third group of process steps, in accordance with one embodiment of the present invention. Memory structure 200 of FIG. 6b represents active stacks 261a and 261b fabricated under the second variation of the third group of process steps, in accordance with one embodiment of the present invention.

The device structures for FeFETs shown in FIGS. 6a and 6b suggest that the 3-dimensional NOR memory strings disclosed herein have a significant advantage in that they offer a relatively large surface area of polarizable material in polarization layer 264. A large surface area in polarization layer 264 provides a tight distribution of voltages in the erased and programmed states, which are difficult to achieve in planar FeFETs. Moreover, that relatively large surface area of polarizable material is provided entirely on the sidewalls of the active strips, and thus provides a high storage density without the penalty of substrate silicon real estate. Specifically, as the normal to polarization layer 264 is substantially horizontal, the polarization layer incurs little silicon real estate. In one embodiment, for example, the footprint of polarization layer 264 of each FeFET, as projected onto the planar surface of the semiconductor substrate, is merely 10.0 nm by 45.0 nm. (However, the Z-Y plane surface area of ferroelectric is 100.0 nm×45.0 nm, with a 100.0 nm channel length.) In some embodiments, the projected footprint of the polarization layer even lies within

the projected footprint of the common bit line, as the ferroelectric or polarization layer is provided in a recess between the common bit line and the common source line. The advantage is amplified in a 3-dimensional array of NOR memory strings, as the footprint of the polarization material is substantially independent of the number of active strips in an active stack. Furthermore, the substantial advantage over planar storage transistors is achieved even if only one polarization layer is provided along the vertical or Z-direction. In that regard, the FeFETs of the present invention is advantageous even in embedded single-layer memory application.

In the process illustrated above, the channel region and the gate stack of each thin-film FeFETs are formed inside one of shafts 218, which may have a curved contour in an X-Y plane cross section. In that embodiment, each rectilinear portion of trenches 267 between shafts 218 provide an isolation between adjacent FeFETs in a NOR memory string. Alternatively, replacement of SiN layers 204a and 204e may begin from the rectilinear portions of trenches 267 between 20 adjacent shafts 218, rather than through shafts 218. Subsequent to metal replacement, channel polysilicon 250 and the gate stacks may be formed as described above, except that the subsequent process steps take place in the rectilinear portions of trenches 267 between shafts 218, rather than in 25 shafts 218.

After the gate stacks are formed using any group of the process steps illustrated herein by FIGS. 3a-3b, FIGS. 4a-4c, FIGS. 5a-5c, 5d(i), 5e(i) and 5f(i), and FIGS. 5a-5c, 5d(ii), 5e(ii) and 5f(ii), an isolation layer is provided over 30 memory structure 200. Thereafter, global interconnect conductors may be formed using any suitable conventional technique to provide connections to the local word lines in the array portion, and to the common bit lines in the staircase portions. FIG. 2l shows a X-Z plane cross section through a 35 row of local word lines 208g-a, showing active strips formed out of active layers 202-7 and 202-6. Each local word line is formed by conductive material 272 in a rectilinear portion of one of trenches 267, between adjacent shafts 218. As layers 204b and 204d, and low-resistivity metal layer 204f. (In FIG. 2k, low-resistivity metal layer 204g is not shown.) Embedded in dielectric layers 231L and 231R are gate stack elements formed out of channel polysilicon layer 250 and polarization layer 264, each element being located between 45 n⁺ silicon layers **204**b and **204**d (along the Z-direction). The gate stack of each FeFET is formed by one such gate stack element and its adjacent local word line 208g-a (connected to a global word line conductor running above the array portion) or 208g-s (connected to a global word line conduc- 50 tor running between the array portion and the silicon substrate). This configuration represents a high packing density configuration for FeFETs.

N⁺ silicon layer **204***b* (i.e., the common source line) can be pre-charged to a suitable voltage required for an operation 55 of the FeFETs at hand (e.g., program voltage \mathbf{V}_{prog} , inhibition voltage V_{inhibit}, erase voltage V_{erase}, or the read reference voltage V_{bl}). During a read operation, any of FeFETs that are in the "on" state conduct current in the vertical or Z-direction between n⁺ silicon layers **204***d* and **204***b*.

Optional metal layers 204f and 204g reduce the resistance of n⁺ silicon layers 204b and 204d, respectively, thus facilitating fast memory device operations. N⁺ silicon layer 204b in any active strip 202 may be left floating. One or more local word lines (referred to as a "pre-charge word line") may be used to bias n⁺ silicon layer 204b. When a suitable voltage is applied to selected pre-charge word lines (i.e.,

rendering the associated pre-charge FeFET conducting), each pre-charge word line momentarily inverts its channel region (i.e., channel polysilicon layer 250 in its associated gate stack), so that n⁺ silicon layer 204b (i.e., the common source line) may be pre-charged to the pre-charge voltage V_{ss} in the associated n⁺ silicon layer 204d from voltage source V_{bl} in the substrate. When the voltage on the precharge word line is withdrawn, (i.e., when the pre-charge FeFET is returned to its non-conducing state) and all the other word lines on both sides of the active strip are also "turned off", device operation may proceed with n⁺ silicon layer 204d left electrically charged to provide a virtual voltage reference at the pre-charged voltage V_{ss} (typically, ~0V) because the distributed parasitic capacitor formed between n⁺ silicon layer 204b and its multiple local word lines is sufficiently large to hold the virtual voltage reference long enough to support the program, program-inhibit or read operation (see below). Although the FeFETs in a NOR memory string may also serve as pre-charge FeFETs along each NOR memory string, to speed up the pre-charge phase for a read operation (read pre-charge requires a lower word line voltages of typically ~1.0 volts), some of the FeFETs (e.g., one in every 32 or 64 FeFETs along the NOR memory string) may also be activated. It is preferable that FeFETs dedicated entirely to serve as pre-charge FeFETs are provided, as they are more tolerant of program-disturb conditions than the FeFETs used in memory operations.

FIG. 2m shows a three-dimensional view of horizontal active layers 202-4 to 202-7 in the embodiment of FIG. 2k, with local word lines or local pre-charge word lines connected to global word lines 208g-s and global word lines 208g-a, and showing each active strip as having its N⁺ layer 204d (acting as a drain region) connected through select circuits to any of voltage supplies (e.g., V_{ss} , V_{bl} , V_{pgm} , V_{inhibit}, and V_{erase}), decoding, sensing and other circuits arranged either adjacent or directly underneath the memory arrays. The substrate circuitry is represented schematically by 206-0 and 206-1 in substrate 201.

In the exemplary fabrication processes described in detail shown in FIG. 21, each active strip includes n⁺ polysilicon 40 above, channel polysilicon 250 is formed prior to forming polarization layer 264. Higher quality channel regions for the FeFETs may be achieved in a process in which the channel region is formed after forming polarization layer 264. In such a "channel-last" approach, the polarization layer and the channel layer may be optimized independently, which improves both endurance and disturb immunity. FIGS. 11a to 11r illustrate a process in which the channel region for an FeFET is formed after forming the polarization layer. The channel region in such a process is achieved by replacing at least in part a previously formed sacrificial layer with a channel material after the polarization layer is formed, according to one embodiment of the present inven-

> FIG. 11a illustrates, for example, at the beginning of processing the array portion of the 3-dimensional array of NOR memory strings (e.g., as described above in conjunction with FIG. 2d above), isolation layer 1105 (e.g., a SiOC layer) and a hard mask layer (e.g., a carbon hard mask) are deposited over the active layers. The hard mask layer is then photo-lithographically patterned, so that its pattern may be transferred to allow etching a first group of trenches (i.e., trenches 1118) in memory structure 200. Trenches 1118 each extend through isolation layers 1105 and 203-7 to 203-0, active layers 202-7 to 202-0 and etch-stop layer 241. Trenches 1118 are each substantially rectilinear in an X-Y plane cross section, each trench being 70.0 nm wide along the X-direction, for example, with corresponding edges of

adjacent trenches separated 190.0 nm from each other. In that embodiment, trenches 1118 are etched at an aspect ratio that is less than 50 (and, in some embodiments, even less than 30).

A series of etching steps then recess silicon nitride layers 5 **204**a and **204**e, n⁺-doped silicon layers **204**b and **204**d and sacrificial oxide layer 204c of each active layer 202 by, for example, 10.0 nm. Thereafter, sacrificial layer 1150 (e.g., amorphous silicon or polysilicon layer, or any other suitable sacrificial material) is conformally deposited to provide a 10 "stand-in" layer for the channel regions to be subsequently formed. Sacrificial layer 1150 is then etched back (i.e., a separation etching step). As described above, n+ silicon layers 204b and 204d in each active layer 202 would provide the common bit line and the common source line for the 15 FeFETs in a NOR memory string to be formed. The hard mask and any excess sacrificial material of sacrificial layer 1105 on top of memory structure 200 are then removed (e.g., by CMP). Resulting memory structure 200 is shown in the X-Z plane cross section of FIG. 11a.

Thereafter, dielectric material 1123 is deposited to fill trenches 1118. Any excess dielectric material is removed from the top surface of semiconductor structure 200, which is then planarization (e.g., by CMP). FIG. 11b shows resulting memory structure 200 in an X-Z plane cross section.

Subsequently, one or more additional groups of trenches may be etched in the same manner as described above for trenches 1118. Forming trenches in multiple groups allow each group of trenches to be etched at a favorable aspect ratio (e.g., less than 50). FIG. 11c shows resulting memory 30 structure 200 in an X-Z plane cross section, after the second group of trenches are formed, according to one embodiment of the present invention. Both groups of trenches are shown in FIG. 11c and each indicated by reference numeral 1118.

In this embodiment, through the unfilled ones of trenches 35 1118, SiN layers 204a and 204e of each active layer 202 may be removed using a wet etch, for example. FIG. 11d shows resulting memory structure 200 in an X-Z plane cross section, after removal of SiN layers 204a and 204e from each of active layers 202-0 to 202-7, according to one 40 embodiment of the present invention.

Conductive material 1129 may then be deposited using, for example, an ALD technique. Conductor material 1129 may be, for example, a liner (e.g., titanium, titanium nitride (TiN), tantalum (Ta), tantalum nitride (TaN), or any com- 45 bination thereof), followed by a refractory metal (e.g., tungsten (W), tungsten nitride, or Molybdenum (Mo)). Conductive material 1129 may be deposited into the cavities resulting from removing SiN layers 204a and 204e from each of active layers 202-0 to 202-7. Thereafter, an etch- 50 back step or an anisotropic etching step removes conductive material 1129 from the bottom of trenches 1118, leaving a substantially conformal layer on their sidewalls. FIG. 11e shows resulting memory structure 200 in an X-Z plane cross section, after replacement of SiN layers 204a and 204e from 55 each of active layers 202-0 to 202-7 by conductive material 1129, according to one embodiment of the present invention.

Conductive material 1129 form optional conductor layers in contact with N⁺ polysilicon layers 204b and 204d to reduce the resistivity in the common bit line and the common source line of a NOR memory string to be formed. Conductive material 1129 may be removed from the sidewalls of trenches 218 and to be further recessed by additional etching. A series of etching steps then recess n⁺-doped silicon layers 204b and 204d, and sacrificial oxide layer 65 204c of each active layer 202 by, for example, 10 nm. FIG. 11f shows resulting memory structure 200 in an X-Z cross

26

section, after recessing conductive material 1129, n⁺-doped silicon layers 204b and 204d and sacrificial oxide layer 204c of each active layer 202, according to one embodiment of the present invention.

Thereafter, channel stand-in sacrificial layer 1150 may be conformally deposited on the sidewalls of trenches 1118, in substantially the same manner as described above in conjunction with FIG. 11a. FIG. 11g shows resulting memory structure 200 in an X-Z cross section, after deposition of sacrificial material 1150 into trenches 1118, according to one embodiment of the present invention.

Trenches 1118 may then be filled by dielectric material 1123 and planarized in the same manner as described above in conjunction with the steps for providing dielectric material 1123 of FIG. 11b. FIG. 11h shows resulting memory structure 200 in an X-Z cross section, after deposition of dielectric material 1123 into trenches 1118, according to one embodiment of the present invention.

After silicon oxide 1123 is deposited into trenches 1118 20 and planarized, the next steps form the gate stacks (i.e., polarization layer 264 and the gate electrode (i.e., the "word line" or the "local word line") for each FeFET of the NOR memory strings. After planarization of dielectric material 1123, hard mask 260 is provided over memory structure 200, photolithographically patterned and developed. Hard mask 260 includes columns of oval openings 261. (In this description, a "column" of objects denotes objects aligned along the Y-direction, whereas a "row" of objects denotes objects aligned along the X-direction.) FIG. 11i(i), oval openings 261 in each column are provide staggered relative to oval openings 261 in the immediately adjacent columns along the X-direction, such that the closest openings in adjacent columns have a greater separation between them than if such openings are aligned along the X-direction. In one embodiment, adjacent openings within each column are located at a 110.0 nm pitch along the Y-direction, while adjacent columns are also provided at 110.0 nm pitch along the X-direction. In this embodiment, the major and minor axes of each opening may be, for example, 100.0 nm and 60.0 nm along the X-direction and the Y-direction, respectively. A series of etchings through openings 261 excavates corresponding shafts 218, removing dielectric material 1123 from trenches 1118 and portions of isolation layers 203-0 to 203-7, and active layers 202-0 to 202-7, reaching down to etch step layer 241. FIGS. 11i(i) and 11i(ii) show, respectively, a top view and an X-Z plane cross section view (along line A-A' of FIG. 11i(i)) of resulting memory structure 200, after shafts 218 are formed, according to one embodiment of the present invention. The aspect ratio of the etch steps creating shafts 218 have an aspect ratio that is still substantially within the desirable range (e.g., less than 50).

Next, silicon nitride liner 284 (e.g., 5 nm thick) is deposited conformally on the sidewalls of shafts 218, which are then filled by sacrificial material 286. Sacrificial material 286 may be, for example, a deposited amorphous silicon. Hard mask 260 is then removed and the surface of memory structure 200 is planarized (e.g., by CMP). FIGS. 11j(i) and 11j(ii) show a top view and an X-Z plane cross section view (along line A-A' of FIG. 11j(i)) of resulting memory structure 200, after shafts 218 are filled by sacrificial material 286, respectively, according to one embodiment of the present invention.

A wet etch then removes dielectric material **1123** from portions **1119** of trenches **1118** between filled shafts **218**. Portions **1119** are also referred to as shafts **1119**. FIGS. **11**k(i) and **11**k(ii) show, respectively, a top view and an X-Z plane cross section view (along line A-A' of FIG. **11**k(i)) of

resulting memory structure 200, after dielectric material 1123 is removed from shafts 1119, according to one embodiment of the present invention.

Thereafter, within each of shafts 1119, gate stacks for FeFETs in the adjacent active strips are provided. As dis-5 cussed above, each gate stack may be formed using any of the techniques illustrated above in conjunction with FIGS. 3a-3b, FIGS. 4a-4c, FIGS. 5a-5c, 5d(i), 5e(i) and 5f(i), or FIGS. 5a-5c, 5d(ii), 5e(ii) and 5f(ii). In some embodiments, as discussed above, each gate stack may include (a) polar- 10 ization layer 264 having, for example, both an interface dielectric layer and the ferroelectric material (e.g., zirconium-doped hafnium oxide) and (b) a gate electrode formed out of conductive material 272 common to FeFET in the adjacent active strips. Conductive material 272 may include 15 a metal liner (e.g., titanium (Ti), titanium nitride (TiN), tantalum (Ta) or tantalum nitride (TaN)) and a refractory metal (e.g., tungsten (W), tungsten nitride (WN) or molybdenum (Mo)). Formed in this manner, the gate electrode may be 60.0 nm×60.0 nm or less. A planarization step (e.g., 20 CMP) may remove excess conductive material 272 from the top surface of memory structure **200**. FIGS. 11l(i) and 11l(ii)show a top view and an X-Z plane cross section view (along line A-A' of FIG. 11l(i)) of resulting memory structure 200, respectively, after the gate stacks to FeFETs have been 25 formed in each of shafts 1119, according to one embodiment of the present invention.

Cap 272t may be provided to facilitate contact to the underlying local word line and to protect polarization layer 264 in subsequent processing steps. Cap 272t may be formed 30 using additional masking, patterning, deposition (e.g., conductive material 272), and planarization steps. FIGS. 11m(i) and 11m(ii) show a top view and an X-Z plane cross section view (along line A-A' of FIG. 11m(i)) of resulting memory structure 200, respectively, after cap 272t has been provided 35 in each of shafts 1119, according to one embodiment of the present invention.

Sacrificial material **286** and silicon nitride liner layer are then removed from shafts **218** to provide access to sacrificial layer **1150**. FIGS. $\mathbf{11}n(i)$ and $\mathbf{11}n(ii)$ show a top view and an 40 X-Z plane cross section view (along line A-A' of FIG. $\mathbf{11}n(i)$) of resulting memory structure **200**, respectively, after sacrificial material **286** and silicon nitride liner layer **284** have been removed from each of shafts **1119**, according to one embodiment of the present invention.

Thereafter, sacrificial material 1150 is removed by a wet etch, for example. FIGS. 11o(i) and 11o(ii) show a top view and an X-Z plane cross section view (along line A-A' of FIG. 11o(i)) of resulting memory structure 200, respectively, after sacrificial material 1150 is removed from each of shafts 218, 50 according to one embodiment of the present invention.

P⁻-doped polysilicon is then deposited as channel polysilicon layer **250** into the cavities resulting from removing sacrificial material **1150** from underneath polarization layer **264**. At the same time, p⁻-doped polysilicon is also deposited up to 10.0 nm on the sidewalls of shafts **218**. FIGS. **11**p(i) and **11**p(ii) show a top view and an X-Z plane cross section view (along line A-A' of FIG. **11**p(i)) of resulting memory structure **200**, respectively, after a wet etch recesses channel polysilicon layer **250** underneath polarization layer **60 264** in shafts **1119**, according to one embodiment of the present invention.

Channel polysilicon layer **250** may be further recessed to provide greater isolation. This etching step also likely removes p^- -doped channel polysilicon layer **250** from the 65 sidewalls of shafts **218**. FIGS. 11q(i) and 11q(ii) show a top view and an X-Z plane cross section view (along line A-A'

28

of FIG. 11q(i)) of resulting memory structure 200, respectively, after channel polysilicon layer 250 is deposited into the cavities resulting from removing sacrificial material 1150 from shafts 1119, according to one embodiment of the present invention.

Thereafter, an ALD silicon oxide liner **285** (e.g., 10 nm) is deposited into the space under polarization layer **264** resulting from recessing channel polysilicon layer **250**. At the same time, silicon oxide liner **285** is also deposited on the sidewalls of shafts **218** to seal access to channel polysilicon **250** from shafts **218**. FIGS. 11r(i) and 11r(ii) show a top view and an X-Z plane cross section view (along line A-A' of FIG. 11r(i)) of resulting memory structure **200**, respectively, after channel polysilicon material **250** is deposited into the cavities resulting from removing sacrificial material **1150** from shafts **1119**, according to one embodiment of the present invention.

FIG. 6c shows memory structure 200, fabricated under a channel-last fabrication process, in accordance with one embodiment of the present invention. As shown in FIG. 6c, channel polysilicon layer 250 is deposited conformally to the walls of the cavity resulting from removal of sacrificial layer 1150. In the example of FIG. 6c, channel polysilicon layer 250 does not completely fill the cavity.

A silicon oxide may then be deposited into shafts **218** to provide further isolation between the FeFETs. Alternatively, shafts **218** may also be left unfilled, allowing the air gaps to serve as the isolation between thin-film storage transistors.

According to yet another embodiment of the present invention, the FeFETs of the present invention need not have a silicon channel region. Instead, the channel region may be formed using an 8.0 to 15.0 nm thick oxide semiconductor material. Suitable oxide semiconductor materials include (i) indium zinc oxide (InZnO, or "IZO") and (ii) indium gallium zinc oxide (InGaZnO, or "IGZO"). An IZO or IGZO channel region has the advantage of a high mobility for greater switching performance and without concern for electron or hole tunneling. For example, a 10.0 nm thick IZO film has an electron mobility of 40.6 cm²/V, relative to an electron mobility of 5.6 cm²/V in an aluminum zirconium oxide (AZO) film of comparable thickness, and an electron mobility of 1.0-4.0 cm²/V in a polysilicon film with the same comparable thickness. Similarly, an IGZO film has an electron mobility of 10.0-100.0 cm²/V, depending on the relative compositions of indium, gallium and zinc.

IZO and IGZO each can provide a junction-less transistor channel, and each may be deposited at a temperature lower than 400° C., even as low as 250° C. Thus, they may be used to provide the channel regions of storage transistors in a 3-dimensional memory array, regardless of whether such storage transistors are FeFETs (i.e., transistors with a polarizable layer) or a quasi-volatile (QV) storage transistors, such as those described in the Related Application. Modifying conventional storage transistors to use IZO or IGZO as channel material imposes only a relatively modest burden on the thermal budget. This is particularly advantageous in applications where planar support logic and analog circuitry is provided at the surface of the substrate beneath the 3-dimensional array of memory strings.

Furthermore, in a junction-less transistor, the source region and the drain region (i.e., bit line) are redundant, and may be replaced by a metal layer (e.g., Mo), thereby dispensing with the need for the n⁺ polysilicon layers to be strapped to metal strapping layers (e.g., metal layers 204f and 204g) in each active strip. These metal layers are referred to as the "metal common bit line" and the "metal source line." The gate stack may be substantially the same

as that described above with respect to the FeFETs with the silicon channel (e.g., an silicon nitride interface dielectric layer, a polarization layer, and the gate electrode). As this alternative FeFET does not have a p/n junction, any leakage current from an FeFET in the programmed state is relatively small. Consequently, such an FeFET is particularly advantageous for high-temperature applications. Such an FeFET may also have a relatively shorter channel length, as there is no need for a margin in the channel length to allow for undesirable dopant diffusion during thermal cycles from the heavily-doped semiconductor common bit line and the common source line. The metal common bit line and metal common source line may each also have a reduced thickness. For example, in one embodiment, a 30.0-40.0 nm thick metal common bit line and a 30.0-40.0 nm thick metal 15 common source line, a 30.0-40.0 nm thick IZO channel layer, and a 30.0 nm thick isolation layer (e.g., SiOC), for a relatively thinner total of 120.0-150.0 nm. The common source line and the common bit line may be formed by replacing a sacrificial material (e.g., carbon) by metal in a 20 metal-replacement step carried out in a later portion of the fabrication process. These thinner layers reduce the aspect ratio for anisotropic etching of trenches through multiple layers of memory strings. Furthermore, the thinner layer of each string results in reduced parasitic capacitance C 25 between local word-lines and the side walls of each string, which reduces the RC delays as well the power dissipation for the 3-dimensional memory array.

Programming, Erasing and Read

An FeFET may be polarized to either (i) a conducting or 30 "erased" state, or (ii) a non-conducting or "programmed" state. In an FeFET, its threshold voltage in the erased state is lower than its threshold voltage in the programmed state. FIG. 7a shows hysteresis in the drain current (I_a) in response to an applied gate voltage (Vg) in a conventional FeFET. 35 (The conventional FeFET is formed at the planar surface of a monocrystalline semiconductor substrate and is not formed as a thin-film field effect transistor.) In FIG. 7a, waveform **1401** traces the drain current of the FeFET in its erased state, as the gate voltage increases from less than -1.0 volts to greater than 1.0 volts, and waveform 1402 traces the drain current of the FeFET in its programmed state, as the gate voltage decreases from greater than 1.0 volts to less than -1.0 volts. As seen in FIG. 7a, the conventional FeFET has a negative threshold voltage (V_t).

It is desirable in some applications, however, for an 45 FeFET (e.g., a thin-film FeFET in a NOR memory string) to have a positive threshold voltage (V_t), such as around 0.5 volts, to prevent undesirable leakage current when it is subject to a disturb condition (e.g., as experienced by a neighboring non-selected FeFET in the NOR memory $_{50}$ string) during a read operation.

FIG. 7b shows a desirable hysteresis in the drain current (I_d) response to an applied gate voltage (V_g) in a thin-film FeFET in a NOR memory array, according to one embodiment of the present invention. In FIG. 7b, waveform 1403 traces drain current I_d of the FeFET in its low-threshold voltage polarization state ("erased state", in this embodiment), as the gate voltage increases from less than -1.0 volts to greater than 1.0 volts, and waveform 1404 traces the drain current of the FeFET in its high-threshold voltage polarization state ("programmed" state, in this embodiment), as the gate voltage decreases from greater than 1.0 volts to less than -1.0 volts. As seen in FIG. 7b, the FeFET has a positive threshold voltage (V_t) of about 0.5 volts, and 1.0-2.5 volts threshold voltage difference ("window") between the polarization states. With a p polysilicon channel region (e.g., 65 boron-doped), this window may be achieved by (i) increasing the boron dopant concentration in the channel region, (ii)

providing a gate electrode that is formed out of a conductive material that has a high work function (e.g., tungsten (W), molybdenum (Mo), aluminum (Al), ruthenium (Ru), tantalum (Ta), or titanium (Ti), (iii) appropriate biasing in the common source region (see below), or (iv) the combination of (i), (ii) and (iii). In an embodiment in which the FeFET channel is an oxide semiconductor material (e.g IGZO), the FeFET's low-threshold voltage state can be permanently shifted to a positive voltage by reducing the IGZO thickness or by a thermal anneal at a specific temperature.

In an array of NOR memory strings, its FeFETs share a common bit line and a common source line. FeFETs in different NOR memory strings may share a common word line. For example, FeFETs at corresponding positions in NOR memory strings in the same active stack, or in NOR memory strings of adjacent active strips, may share a common word line. FIG. 7c shows a circuit schematic diagram of FeFETs in NOR memory strings 1420-1 to 1420-4, according to one embodiment of the present invention. In FIG. 7c, reference numerals 1411, 1412 and 1413 indicate the FeFETs that share common bit line BL, and common source line SL_i. Likewise, reference numerals 1412, 1413 and 1414 indicate FeFETs that share common word line WL_k. Even though not selected for a program, erase or read operation, an FeFET may experience a "disturb" condition affecting the stability of the threshold voltage state it is set. The disturb condition arises because the voltages on the FeFET's word line, source line or bit line may be subject to interference by signal changes in its vicinity as to affect a change in its polarization state during the program, erase or read operation. The signal changes are typically associated with the voltages imposed on the word line or bit line the FeFET shares with the selected FeFET.

To facilitate reference, in this detailed description, the term "affected FeFET" refers to a non-selected FeFET that shares a bit line or a word line with a selected FeFET of an intended program, erase or read operation. Any of the FeFETs indicated by reference numerals **1411**, **1413**, **1414** and **1415** is an affected FeFET. Likewise, the term "onceremoved FeFET" refers to an FeFET that shares a word line or a bit line with an affected FeFET. All the unlabeled FeFETs shown in FIG. **7**b are once-removed FeFETs.

Without mitigation, a disturb condition may cause an affected FeFET to move away from its polarization state and thereby loses the data value it is holding. To mitigate a disturb condition, the non-selected word line or the non-selected bit line of an affected FeFET may be biased to a predetermined voltage ("inhibit voltage") during the program, erase or read operation to counteract the effect of the voltage imposed on the word line or bit line it shares with the selected FeFET.

Table 1 provides exemplary bias voltages for the gate electrodes (or word lines), the source regions (or the common source lines), the drain regions (or the common bit lines) for the selected and affected FeFETs in a 3-dimensional array of NOR memory strings during erase, program and read operations, under a first disturb mitigation scheme ("V/2 scheme"; in which an affected FeFET is subject to one-half the "stress voltage" of a selected FeFET, during a program or erase operation). The stress voltage refers to the difference in voltage between the gate electrode and the bit line in an FeFET, during a program, erase or read operation. The bias voltages of Table 1 are based on the case when (i) an FeFET holding logic value '1' (i.e., in the low-threshold voltage polarization state, with a threshold voltage of approximately -0.25 volts), and (ii) an FeFET holding logic value '0' (i.e., in the high-threshold voltage polarization state, with a threshold voltage of approximately 0.6 volts).

TABLE 1

	Selected FeFET			Non-selected FeFETs that Share a Bit Line or a Gate Electrode with the Selected FeFET			
Operation	Gate/Word Line	Source	Drain/Bit Line	Word Line	Source Line	Drain/Bit Line	
Create Low-threshold Voltage Polarization State ('1')	0.0 volts	3.0 volts	3.0 volts	1.5 volts	1.5 volts	1.5 volts	
Create High-threshold Voltage Polarization State ('0')	3.0 volts	0.0 volts	0.0 volts	1.5 volts	1.5 volts	1.5 volts	
Read	1.0 volts	0.75 volts	1.25 volts	0.0 volts	0.0	0.0	

Under the V/2 scheme, during a program or erase operation, an affected FeFET (e.g., in FIG. 7c, any of the FeFETs indicated by reference numerals 1411, 1413, 1414 and 1415) experiences a stress voltage of -1.5 volts (write '1') or 1.5 volts (write '0'). The once-removed FeFETs each experience 20 no stress voltage (i.e., a stress voltage of 0.0 volts).

In this example, as the threshold voltage of an FeFET in the '0' state is negative, during a read operation, the source voltage is set to 0.75 volts, or any suitable positive voltage sufficient to prevent conduction prior to the gate electrode 25 achieving the read voltage. The higher bias voltage (the "boost") avoids both having to generate a negative voltage and having to provide corresponding additional isolation structures to protect the devices in the silicon substrate. During a read operation, to reduce read disturb under the V/2scheme, the voltage difference between the word line and the common bit line or the common source line in the selected FeFET is maintained at less than the voltage that may change the polarization state of the selected FeFET.

Table 2 provides, during erase, program and read operations under a second disturb mitigation scheme ("V/3 scheme"), exemplary bias voltages for the gate electrodes (or word lines), the source regions (or the common source lines), the drain regions (or the common bit lines) of selected and affected FeFETs in a 3-dimensional array of NOR memory strings. Under the V/3 scheme, an affected FeFET 40 is subject to one-third the stress voltage of a selected FeFET during a program or erase operation. The bias voltages of Table 2 are also based on the case when (i) an FeFET holding logic value '0' (i.e., in the "programmed" state) has a threshold voltage of approximately 0.6 volts, and (ii) an 45 FeFET holding logic value '1' (i.e., in the "erased" state) has a threshold voltage of approximately -0.25 volts.

or 1.0 volts (write '1'). Thus, the V/3 scheme provides a lesser stress voltage in the affected FeFETs. However, because a voltage difference of 1.0 volt (write '1') or -1.0 volt (write '0') exists between the non-selected bit lines and the non-selected word lines, which accounts for the charge separation in the parasitic capacitors between the nonselected word lines and the non-selected bit lines, the V/3 scheme has a greater power dissipation than the V/2 scheme during a program, erase or read operation.

Under either the V/3 scheme or the V/2 scheme, the stress voltage for an affected FeFET sharing a word line with the selected FeFET is 1.0 volts during a read operation, which is equal to the lower stress voltage under the V/3 scheme for a program or erase operation. As a program or erase operation is expected to have a significantly longer duration (e.g., 100.0 ns) than a read operation (e.g., 10.0 ns), limiting the overall effects of the disturb conditions may be achieved by adequately addressing the disturb conditions during the program and erase operations.

Even though the data retention time of an FeFET of the present invention is expected to be significantly longer (e.g., tens of days to tens of years) than the data retention time of a quasi-volatile (QV) storage transistor (e.g., any of those transistors described in the Related Application), the methods for maintaining data integrity beyond the data retention time for QV storage transistors are also applicable to maintaining data retention in the FeFETs of the NOR memory strings of the present invention. For example, Provisional Application VII discloses one or more methods for reducing the effects of disturb conditions by refreshing a subset of the storage transistors that share a word line with a selected storage transistor during a program or erase operation ("re-

TABLE 2

	Selected FeFET			Non-selected FeFETs that Share a Bit Line or a Gate Electrode with the Selected FeFET			
Operation	Gate/Word Line	Source	Drain/Bit Line	Word Line	Source Line	Drain/Bit Line	
Create Low-threshold Voltage Polarization State ('1')	0.0 volts	3.0 volts	3.0 volts	2.0 volts	1.0 volts	1.0 volts	
Create High-threshold Voltage Polarization State ('0')	3.0 volts	0.0 volts	0.0 volts	1.0 volts	2.0 volts	2.0 volts	
Read	1.0 volts	0.75 volts	1.25 volts	0.0 volts	0.0 volts	0.0 volts	

Under the V/3 scheme, during a program or erase operation, an affected FeFET experiences a stress voltage of 1.0 65 activated selected word line in the program or erase operavolts (write '0') or -1.0 volts (write '1'). A once-removed FeFETs experiences a stress voltage of -1.0 volts (write '0')

fresh in-place"). Refresh in-place takes advantage of the tion to write back the data values in the storage transistors in the affected subset.

In some mobile applications, the device's battery life is shorter than the data retention time, periodic refresh may be triggered when the battery is detected to be connected to a charger. In this manner, one can be assured that the data in the memory arrays in a fully charged device will persist until 5 the next charge cycle. In addition, refresh operation is initiated when the time since the last refresh is within a threshold (e.g., 10% of the data retention time).

When the body region of an FeFET in a NOR memory string is floating, its programming speed may be slower than 10 its erase speed. In such a condition, the gate-induced drain leakage (GIDL) effect may be harnessed to improve the programming speed. The GIDL effect may be activated by creating a voltage difference (e.g., 0.5 to 2.0 volts) between the common bit line and the common source line (V_{ds}) 15 during programming, for example. This voltage difference may be achieved by first momentarily pre-charging the common source line to a predetermined source line voltage through the common bit line, and then setting the common bit line to its target voltage, as disclosed in the Related 20 Application.

FIG. 8 illustrates the methods and circuit elements used for setting source voltage V_{ss} in n⁺ silicon layer 204b. Specifically, source voltage V_{ss} may be set through hardwired decoded source line connections 280 (shown in 25 dashed line) or alternatively, by activating pre-charge thin film transistors 300 (which, preferably, is not implemented by a FeFET) and decoded bit line connections 270 to any one of bit line voltages V_{ss} , V_{bl} , V_{pgm} , $V_{inhibit}$ and V_{erase} . Alternatively, source reference voltage V_{ss} may be accessed 30 through a metal or n⁺-doped polysilicon conductor connecting from the top of the memory array through staircase vias. Each of the conductors in hard-wired connections 280 may be independently connected, so that the source voltages for different planes or within planes need not be the same. In 35 some embodiments, it is advantageous to dispense with the hard-wired source voltage V_{ss} connections, by taking advantage of a virtual voltage source in the intrinsic parasitic capacitance of the NOR memory string.

The present invention takes advantage of the cumulative 40 intrinsic parasitic capacitance that is distributed along each NOR memory string to increase the number of FeFETs that can be programmed, read or erased in parallel in a single operation (though, preferably, on separate strings), while also significantly reducing the operating power dissipation. 45 As shown in FIG. 8, local parasitic capacitor 360 (contributing to a cumulative capacitance C) exists at each overlap between a local word line (as one plate) and the n⁺/p⁻/n⁺ active layer (as the other plate). Each local parasitic capacitor may be approximately 0.005 femtofarads (each femto- 50 farad is 1×10^{-15} farad), too small to be of much use for temporary storage of charge. However, since there may be a thousand or more FeFETs contributing capacitance from one or both sides of an active strip, the total distributed capacitance C of n⁺ silicon layer **204***b* (the common source line) 55 and n⁺ silicon layer **204**d (the common bit line) in a long NOR memory string can be in the range of ~10.0 to 100.0 femtofarads. This is also roughly the capacitance at sensing circuitry connected through connections 270 (e.g., voltage source V_{bl} .

During a program or erase operation in a selected FeFET, and their respective inhibit operations on a non-selected FeFETs, both n⁺ silicon layers **204***d* and **204***b* are preferably held at the same voltage, therefore the leakage current between n⁺ silicon layers **204***d* and **204***b* is insignificant. 65 During a program or erase operation, charge leakage from cumulative capacitance C flows primarily to the substrate

through the substrate selection circuitry, which has very little transistor leakage, as it is formed in single crystal or epitaxial silicon. Nevertheless, even a 100-microsecond polarization state retention time is sufficient to complete the sub-100 nanosecond read operation or the sub-100 microsecond program operation (see below) of the selected FeFET on the NOR memory string.

34

An FeFET in a NOR memory string, unlike a DRAM cell, may operate as a non-volatile memory transistor, so that, even if parasitic capacitor C of the NOR memory string is completely discharged, the information stored in the selected FeFET remains intact in the ferroelectric storage material (i.e., polarization layer 264). In a DRAM cell, however, the information would be lost without frequent refreshes. Accordingly, distributed capacitance C of a NOR memory string of the present invention is used to temporarily hold the pre-charge voltage on n⁺ silicon layers 204d and **204**b at one of voltages V_{ss} , V_{bl} , V_{progr} , $V_{inhibit}$ or V_{erase} , and is not used to store actual data for any of the FeFETs in the NOR memory string. Pre-charge transistor 303 of FIG. 3a, controlled by word line 151n (i.e., word line 208-CHG), is activated momentarily immediately preceding each read, program, or erase operation to transfer voltage V_{bl} (e.g., through connections 270) from the substrate circuitry (not shown) to n⁺ silicon layer 204b. For example, under the V/3 scheme described above, voltage V_{bl} can be set at 0.75 V to pre-charge n⁺ silicon layer **204**b to a virtual ground voltage ~0.75 V during a read operation, or to pre-charge both n+ silicon layers 204d and 204b to either 1.0 Volts and 2.0 volts in non-selected bit lines involved in refresh in-place during a program or erase operation.

The value of cumulative capacitance C may be increased by lengthening the NOR string to accommodate the thousands more FeFETs along each side of the active strip, correspondingly increasing the retention time of pre-charge voltage V_{ss} on n⁺ silicon layer 204b. However, a longer NOR memory string suffers from an increased line resistance as well as higher leakage currents between n+ silicon layer **204***d* and n⁺ silicon layer **204***b*. Such leakage currents may interfere with the sensed current when reading the one FeFET being addressed with all other FeFETs of the NOR memory string in their "off" (and somewhat leaky) states. Also, the potentially longer time it takes to pre-charge a larger capacitor during a read operation can conflict with the desirability for a low read latency (i.e., a fast read access time). To speed up the pre-charging of the cumulative capacitance C of a long NOR memory string, pre-charge transistors—which may be constructed separately from the memory FeFETs—may be provided spaced apart along either side of the active strip (e.g., once every 128, 256 or more FeFETs).

Because the variable-threshold FeFETs in a long NOR memory string are connected in parallel, the read operating condition for the NOR memory string should preferably ensure that all FeFETs along both edges of an active strip operate in enhancement mode (i.e., they each have a positive threshold voltage, as applied between control gate 151n and voltage V_{ss} at source 204d). With all FeFETs being in enhancement mode, the leakage current between n⁺ silicon layer 204d and n⁺ silicon layer 204b of the active strip is suppressed when all control gates on both sides of the active strip are held at, or below V_{ss}-0V. This enhancement threshold voltage can be achieved by providing polysilicon channel layer 250 with a suitable dopant concentration (e.g., a boron concentration between 1×10¹⁶ and 1×10¹⁸ per cm³ or higher, which results in an intrinsic FeFET threshold voltage of between ~0.5 V and ~1 V). Alternatively, to substantially

shut off a source-drain leakage current in any of the FeFETs along a NOR memory string that is programmed into a depletion threshold voltage state, the common source region and the common drain region may be boosted into a "virtual enhancement" threshold voltage state by pre-charging the 5 common source line and the common bit line of the NOR memory string to voltages V_{ss} and V_{bl} , respectively. This biasing condition substantially shuts off source-drain conduction currents in the FeFETs of the same NOR memory string. The "virtual enhancement" mode of operation is advantageous because it allows the same voltage polarity to be applied to source, drain and gate terminals for all program, erase and read operations of the FeFETs in the NOR memory string. Furthermore, implementation of the "virtual enhancement" mode of operation provides the flexibility of 15 not having to rely on a factory-based operation to apply a fixed, potentially sub-optimal dopant concentration to the polysilicon channel of the FeFET.

Quasi-Volatile NOR Memory Strings

Endurance is a measure of a storage transistor's performance degradation after some number of write-erase cycles. Endurance of less than around 10¹² cycles is considered too low for some applications. Greater endurance in an FeFET of the present invention may be achieved by, for example, increasing the dielectric constant k in the interface dielectric layer, typically a silicon oxide-nitride or silicon nitride film of thickness 0.1 to 3.0 nm. The interface dielectric layer may also be kept thin (i.e., even less than 0.1 nm), or eliminated altogether or replaced with another dielectric film of a higher dielectric constant k significantly greater than 4 (e.g., aluminum oxide).

A FeFET having a low data retention time (i.e., less than, for example, one year), may be referred to as "quasi-volatile FeFET". Such an FeFET may require refreshes to restore the weakened polarization. Because the quasi-volatile FeFET of 35 the present invention provides a DRAM-like fast read access time with a low latency, the resulting quasi-volatile NOR memory strings may be suitable for use in some applications that currently require DRAMs. The advantages of quasivolatile FeFETs in NOR memory strings, arranged in a 40 planar structure or in 3-dimensional memory arrays, over DRAMs include: (i) a much lower cost-per-bit figure of merit because DRAMs cannot readily be built in 3-dimensional blocks, and (ii) a much lower power dissipation, as the refresh cycles need only be run approximately once every 45 tens of minutes or once every few days, as compared to every ~64 milliseconds, as required in state-of-the-art DRAM technology. 3-dimensional memory arrays of FeFETs may also have advantages over the 3-dimensional NOR memory arrays of the Related Application ("3-D QV 50 NOR memory strings"). For example, (i) FeFETs can be fabricated at significantly lower thermal cycles (e.g., at 400.0-500.0° C.) than the thermal cycle for fabricating direct tunneling, charge-trapping thin-film transistors (e.g. at 600.0-860.0° C.), and (ii) the erase and programming polar- 55 ization states of some FeFETs can be switched at significantly lower voltages (e.g., less than 4.0 volts), as compared to the 6.0-8.0 volts required to inject electrons (i.e., to program) or eject trapped electrons (i.e., to erase) in 3-dimensional QV NOR memory strings.

The NOR memory string of FeFETs in the present invention appropriately adopt the program/read/erase conditions to incorporate the periodic data refreshes. For example, because each NOR memory string of quasi-volatile FeFETs is frequently read-refreshed or program-refreshed, it is not 65 necessary to "hard-program" such FeFETs to open a large threshold voltage window between the '0' and '1' states, as

compared to non-volatile FeFETs where a minimum 10-year data retention time may be required. Quasi-non-volatile threshold voltage window may be as narrow as 0.2 volts to 1.5 volts. The reduced threshold voltage window allows such FeFETs to be programmed and erased at lower programming/erase voltages and by shorter-duration programming pulses, which reduce the stress voltage, thereby extending endurance. Periodic refresh operations can be performed in a background mode so as to minimize collision with erase/program/read operations. During such a refresh operation, it is possible for the system to check the health of each FeFET, to restore its full window of operation or altogether retire the FeFET by replacing it with a redundant fresh FeFET.

36

Source-Drain Leakage in Long NOR Memory Strings

In long NOR memory strings, the current of the one accessed FeFET in a read operation competes with the cumulative subthreshold leakage currents from the thousand or more parallel non-selected FeFETs. Similarly, precharged strip capacitor C has to contend with charge leakage not just of one transistor (as in a DRAM circuit) but the charge leakage through the thousand or more transistors in the NOR memory string. That charge leakage reduces substantially the charge retention time on C to perhaps a few hundred microseconds, potentially requiring counter measures to reduce or neutralize such leakage. However, the leakage from a thousand or so transistors is only significant during read operations. During program, program-inhibit or erase operations, source layer 204b and bit line layer 204d are preferably held at the same voltage, therefore transistor leakage between the two layers is insignificant (the leakage of charge from capacitor C during program, program-inhibit or erase operations is primarily leaked to the substrate through the substrate selection circuitry, which is formed in single-crystal or epitaxial silicon where transistor leakage is very small). For a read operation, even a relatively short 100-microsecond retention time of charge on the source and bit line capacitors may be sufficient time to complete the sub-100 nanosecond read operation (see below) of the FeFETs of the present invention. One difference between a FeFET in a NOR memory string of the present invention and a DRAM cell is that, relative to the latter, the former is a non-volatile memory transistor for DRAM suitable applications, so that even if parasitic capacitor C is completely discharged, the information stored in the selected FeFET is not lost from the ferroelectric or polarizable element, unlike a DRAM cell where the stored information is lost unless refreshed. Capacitor C is used solely to temporarily hold the pre-charge voltage on n⁺ silicon layers 204d and 204b at one of voltages $\mathbf{V}_{ss}, \mathbf{V}_{bl}, \mathbf{V}_{progr}, \mathbf{V}_{inhibit},$ or $\mathbf{V}_{erase};$ capacitor C is not used to store actual data for any of the non-volatile FeFETs in the string. Pre-charge transistor 303, controlled by word line 151n (208-CHG) (FIG. 8) is activated momentarily immediately preceding read, program, program-inhibit or erase operations to transfer through connections 270 the voltage V_{bl} from the substrate circuitry (not shown) to capacitor C of n⁺ silicon layer 204d. For example, voltage V_{bl} can be set at ~0V to pre-charge n⁺ silicon layer 204b to a virtual ground voltage during read, or to pre-charge both 60 n⁺ silicon layers 204d and 204b to their programming or erase bias voltages during program inhibit. The value of cumulative capacitors C may be increased by lengthening the active string to accommodate thousands more FeFETs along each side of the string, correspondingly increasing the retention time of pre-charge voltage V_{ss} on n⁺ silicon layer 204d. To speed up the pre-charging of the capacitance C of a long active strip, several pre-charge thin-film transistor

303 may be provided spaced apart along either side of the active strip (e.g., once every 128, 256 or more FeFETs). Revolving Reference NOR Memory String Address Locations to Extend Cycle Endurance.

In applications requiring a large number of write/erase 5 operations, the threshold-voltage window of operation for the FeFETs in the NOR memory strings may drift with cycling, away from the threshold-voltage window that is programmed into the FeFETs of the reference NOR memory strings at the device's beginning of life. The growing 10 discrepancy between FeFETs on the reference NOR memory strings and FeFETs on the addressed memory NOR memory strings over time, if left unattended, can be detrimental to data integrity. To overcome this drift, reference NOR memory strings in a block need not always be mapped to the 15 same physical address and need not be permanently programmed for the entire life of the device. Since the programmable reference NOR memory strings are practically identical to the memory NOR memory strings sharing the same plane in a block, reference NOR memory strings need 20 not be dedicated for that purpose in any memory array block. In fact, any one of the NOR memory strings can be set aside as a programmable reference NOR memory string. In fact, the physical address locations of the programmable reference NOR memory strings can be rotated periodically (e.g., 25 changed once every 100 times the block is erased) among the sea of memory NOR memory strings, so as to level out the performance degradation of memory NOR memory strings and reference NOR memory strings as a result of extensive program/erase cycles. For example, the reference 30 NOR memory strings may be rotated among the 8 planes (e.g., the reference NOR memory strings may be mapped in a round-robin fashion to one of the 8 planes of an active stack, kept track by a 3-bit counter).

According to the current invention, any NOR memory 35 string can be rotated periodically to be designated as a programmable reference NOR memory string, and its address location may be stored inside or outside the addressed block. The stored address may be retrieved by the system controller when reading the NOR memory string. 40 Under this scheme, rotation of reference NOR memory strings can be done either randomly (e.g., using a random number generator to designate new addresses), or systematically among any of the active memory NOR memory strings. Programming of newly designated reference NOR 45 memory strings can be done as part of the erase sequence when all FeFETs on a slice or a block may be erased together, to be followed by setting anew the reference voltages on the newly designated set of reference NOR memory strings. In this manner, all active memory NOR 50 memory strings and all reference NOR memory strings in a block drift statistically more or less in tandem through extensive cycling.

Programmable Reference Slices.

In some embodiments of the present invention, a block 55 may be partitioned into four equal-size quadrants, as illustrated in FIG. 9a. FIG. 9a show semiconductor structure 600, which is a three-dimensional representation of a NOR memory string array of FeFETs organized into quadrants Q1-Q4. In each quadrant, (i) numerous NOR memory 60 strings are each formed in active strips extending along the Y-direction (e.g., NOR memory string 112), (ii) pages extending along the X-direction (e.g., page 113), each page consisting of one FeFET from each NOR memory string at a corresponding Y-position, the NOR memory strings in the 65 page being of the same corresponding Z-position (i.e., of the same active layer); (iii) slices extending in both the X- and

38

Z-directions (e.g., slice 114), with each slice consisting of the pages of the same corresponding Y-position, one page from each of the planes, and (iv) planes extending along both the X- and Y-directions (e.g., plane 110), each plane consisting of all pages at a given Z-position (i.e., of the same active layer).

FIG. 9b shows structure 600 of FIG. 9a, showing FeFETs in programmable reference NOR memory string 112-Ref in quadrant Q4 and FeFETs in NOR memory string 112 in quadrant Q2 coupled to sense amplifiers SA(a), Q2 and Q4 being "mirror image quadrants." FIG. 9b also shows (i) programmable reference slice 114-Ref (indicated by area B) in quadrant Q3 similarly providing corresponding reference FeFETs for slice 114 in mirror image quadrant Q1, sharing sense amplifiers SA(b), and (ii) programmable reference plane 110-Ref in quadrant Q2 providing corresponding reference FeFETs to plane 110 in mirror image quadrant Q1, sharing sense amplifiers SA(c), and also providing corresponding reference FeFETs for NOR memory strings in the same quadrant (e.g., NOR memory string 112).

As shown in FIG. 9b, programmable reference NOR memory strings 112Ref of FeFETs may be provided in each quadrant to provide reference voltages for the NOR memory strings of FeFETs on the same plane in the same quadrant, in the manner already discussed above. Alternatively, programmable reference slices (e.g., reference slice 114Ref) are provided on mirror-image quadrants for corresponding memory slices. For example, when reading a memory slice in quadrant Q1, programmed reference slice 114Ref (area B) in quadrant Q3 is simultaneously presented to sense amplifiers 206 that are shared between quadrants Q1 and Q3. Similarly, when reading a memory slice in quadrant Q3, reference slice 114Ref (area A) of quadrant Q1 is presented to the shared sense amplifiers 206. There can be more than one reference slice distributed along the length of NOR memory strings 112 to partially accommodate mismatched in RC delay between the slice being read and its reference slice. Alternatively, the system controller can calculate and apply a time delay between the global word line of the addressed slice and that of the reference slice, based on their respective physical locations along their respective NOR memory strings. Where the number of planes is large (e.g., 8 or more planes), one or more planes can be added at the top of the block to serve either as a redundant plane (i.e., to substitute for any defective plane) in the quadrant, or as programmable reference pages, providing reference threshold voltages for the addressed pages sharing the same global word line conductor 208g-a. The sense amplifier at the end of each NOR memory string receives the read signal from the addressed page at the same time as it receives the signal from the reference page at the top of the block, since both pages are activated by the same global word line.

In one embodiment, each memory block consists of two halves, e.g., quadrants Q1 and Q2 constitute an "upper half" and quadrants Q3 and Q4 constitute a "lower half." In this example, each quadrant has 16 planes, 4096 (4K) NOR memory strings in each plane, and 1024 (1K) FeFETs in each NOR memory string. It is customary to use the unit "K" which is 1024. Adjacent quadrants Q1 and Q2 share 1K global word lines (e.g., global word line 208g-a) driving 2048 (2K) local word lines 208W per quadrant (i.e., one local word line for each pair of FeFETs from two adjacent NOR memory strings). 4K FeFETs from quadrant Q1 and 4K FeFETs from quadrant Q2 form an 8K-bit page of FeFETs. 16 pages form a 128K-bit slice, and 1K slices are provided in a half-block, thus providing 256 Mbits of total storage per block. (Here, 1 Mbits is 1K×1 Kbits.) The 4K

strings in each plane of quadrants Q2 and Q4 share substrate circuitry 206, including voltage sources for voltage V_{bl} and sense amplifiers (SA). Also included in each quadrant are redundant NOR memory strings that are used as spares to replace faulty NOR memory strings, as well to store quadrant parameters such as program/erase cycle count, quadrant defect map and quadrant ECC. Such system data are accessible to a system controller. For blocks with high plane counts, it may be desirable to add one or more planes to each block as spares for replacing a defective plane.

High-capacity storage systems based on 3-dimensional arrays of FeFETs organized as NOR memory strings of the present invention require a high-speed system controller for error-free massively parallel erase and erase-inhibit, pro- 15 gram and program-inhibit, and read operations that may span many "chips" including many memory blocks. In fact, in some applications, a hierarchy of controllers (e.g., organized as a tree structure) may be provided, with the controllers at the leaf level each managing the operations of 20 multiple chips, each containing a large number of tiles. Controllers at a higher level in the hierarchy (i.e., closer to the root or roots) control controllers closer to the leaves. At the chip level, each tile is provided logic functions implemented in the memory circuits dedicated to assist the 25 controllers in managing operations carried out by the tiles of the chip. Also, each memory circuit stores system parameters and information related to the files stored in the memory circuit. Such system information is typically accessible to the system controller, but not accessible by the user. 30 It is advantageous for the system controller to quickly read the memory circuit-related information. A block under this organization (i.e., 256 Megabits) provides 2K slices. A terabit memory circuit may be provided by including 4K blocks.

As shown in FIGS. 9a and 9b, the FeFETs in quadrants Q2 and Q4 share voltage source V_{bl} , sense amplifiers SA, data registers, XOR gates and input/output (I/O) terminals to and from substrate circuitry 206. According to one organization, FIG. 9a shows NOR memory strings 112, quarter- 40 planes 110, half-slices 114, and half-pages 113. Also shown are optional pillars 290 supplying back-bias voltage V_{bb} from the substrate. FIG. 9b shows examples of locations of reference strings 112(Ref), reference slices 114(Ref) and reference planes 110 (Ref). In the case of reference strings, 45 reference string 112 (Ref) of quadrant Q4 can serve as a reference string to NOR memory string 112 on the same plane in quadrant Q2, the two NOR memory strings being presented to a shared differential sense amplifier SA in circuitry 206. Similarly, reference slice 114 Ref (area A) in 50 quadrant Q1 can serve as reference for a slice in quadrant Q3, while a reference slice B in quadrant Q1 can serve as reference for slices in quadrant Q3, again sharing differential sense amplifiers SA provided between quadrants Q1 and Q3. Global word lines 208g-a are connected to local word lines 55 208W and local pre-charge word lines 208-CHG. Substrate circuitry and input/output channels 206 are shared between FeFETs in quadrants Q2 and Q4. Under this arrangement, their physical locations allow cutting by half the resistance and capacitance of NOR memory strings 112. Similarly, 60 global word line drivers 262 are shared between quadrants Q1 and Q2 to cut by half the resistance and capacitance of the global word lines, and pillars 290 (optional) connect p layers of NOR memory strings 112 to the substrate voltage.

Since silicon real estate on an integrated circuit is a 65 valuable resource, rather than adding reference strings or reference pages to each plane, it may be advantageous to

40

have some or all reference strings or reference pages provided in one or more additional planes. The additional plane or planes consume minimal additional silicon real-estate and the reference plane has the advantage that the addressed global word line 208g-a accesses a reference page at the same time it accesses an addressed page on any of the planes at the same address location along the active strings in the same quadrant. For example, in FIG. 9b, reference string 112Ref, which is shown as dashed line in quadrant Q2, resides in reference plane 110Ref in this example. NOR memory string 112Ref tracks memory NOR memory string 112 being selected for read in the same quadrant and the read signals from the two NOR memory strings reach the differential sense amplifiers SA for that quadrant practically at the same time. Although reference plane 110Ref is shown in FIG. 9b as being provided in the top plane, any plane in the quadrant can be designated a reference plane. In fact, it is not necessary for every NOR memory string on the reference plane to be a reference string: e.g., every one in eight NOR memory strings can be designated as a reference NOR memory string that is shared by eight NOR memory strings in other planes. The remainder of NOR memory strings in the reference plane may serve as spare strings to substitute for defective strings on the other planes in the block.

Alternatively, one or more additional planes (e.g., plane 117 in FIG. 9c) can be set aside to serve as spare memory resources to substitute for defective NOR memory strings, defective pages or defective planes in the same quadrant.

As related to electrically programmable reference strings, slices, pages or planes, once set in their designated threshold voltage states, care must be exercised at all times to inhibit their inadvertent programming or erasing during programming, erasing or reading the non-reference strings.

A very large storage system of 1 petabyte (8×10¹⁵ bits) 35 requires 8,000 1-terabit memory circuits ("chips"), involving 32M blocks or 64 G slices. (1 Gbits is 1K×1 Mbits). This is a large amount of data to be written (i.e., programmed) or read. Therefore, it is advantageous to be able to program or erase, and read in parallel a great many blocks, slices or pages on numerous chips at once, and to do so with minimum power dissipation at the system level. It is also advantageous for a terabit capacity memory chip to have many input/output channels such that requested data can be streamed in and out in parallel concurrently from and to a large number of blocks. The time required to track down the physical location of the most current version of any given stored file or data set would require a significant amount of time for the system controller to maintain, such as the translation the logical address into the most current physical addresses. The translation between logical to physical addresses would require, for example, a large look-up file allocation table (FAT) to correctly access the slice in the right block on the right chip. Such a search could add considerable read latency (e.g., in the range of 50-100 microseconds) which would impede a fast-read access goal (e.g., under 100 nanoseconds). Accordingly, one aspect of the present invention significantly reduces the search time by introducing a system-wide parallel on-chip rapid file searches, so as to significantly reduce the latency associated with a large FAT, as described below.

Fast Reads: Pipelined Streaming and Random Access

At system initiation of a multi-chip storage system of the present invention, all chips may be erased and reference strings, reference slices or reference planes may be uploaded to their reference states. In one embodiment, one such multi-chip storage system includes a memory module that includes multiple memory chips connected to a module

controller which interfaces with the host system controller and manages the multiple memory chips without participation by the host system controller. The memory circuits in the memory module may communicate with each other and with the module controller through an interposer substrate, 5 through-silicon vias (when the memory circuits are stacked one on top of another), hybrid bonding or any other suitable multi-chip assembly technique.

The module controller (sometime also referred to as a "chiplet") designates as cache storage the memory slices 10 (e.g., slice 116 in FIG. 9c) that are physically closest to the sense amplifiers and voltage sources 206. Because of the RC delays along the length of each NOR memory string, the FeFETs in each NOR memory string that are physically closest to substrate circuitry 206 may have their voltages V_{bl} 15 established a few nanoseconds sooner than the FeFETs furthest from substrate circuitry 206. For example, the first ~50 slices or so (shown as slice 116 in FIG. 9c) out of the 1K slices in each quadrant have the shortest latency and can be designated as a cache memory or storage, to be used for 20 storing quadrant operational parameters, as well as information regarding the files or data set stored in the quadrant. For example, each memory page (2×4 Kbits) or slice (2×4 Kbits×16=128 Kbits) written into the upper half-block (i.e., quadrants Q1 and Q2) can have a unique identifier number 25 assigned to it by the module controller, together with an index number that identifies the type of file that is stored.

The cache storage may be used to store on-chip resource management data, such as file management data. A file can be identified, for example, as "hot file" (i.e., associated with 30 a large number of accesses, or a "high cycle count"), "cold file" (i.e., has not been altered for a long time, and is ready to be moved to slower storage or archival memory at a future time), "delete file" (i.e., ready for future erase in background mode), "defective file" (i.e., to be skipped over), or "sub- 35 stitute file" (i.e., replacing a defective file). Also included in the identifier may be a time stamp representing the last time and date the file associated with the identifier was written into the quadrant. Such unique identifier, typically between 32-bit and 128-bit long can be written into one or more of 40 the cache slices as part of the writing of the file itself into the other memory slices in the same half-tile. Files are written sequentially into available erased space, and the identifiers can be assigned by incrementing the previous unique identifier by one for each new file written into memory. If 45 desired, new files can be written into partial slices and the unwritten part of the slice can be used for writing part or whole of the next file, to avoid wasting storage space. Writing sequentially until the entire memory space of the system is used helps level out the wear-out of FeFETs 50 throughout the system. Other on-chip resource management data may include chip, tile, plane, slice, page and string parameters, address locations of faulty strings and their replacement strings, defective pages, defective planes, defective slices and defective tiles and their substitute 55 replacements, file identifiers for all files resident in the tile, look up tables and link lists for skipping over unusable memory, tile-erase cycle counts, optimum voltages and pulse shape and durations for erase, program, programinhibit, program scrub, read, margin read, read refresh, read 60 scrub operations, error correcting codes, and data recovery modes, and other system parameters.

Because of the modularity of each chip at the tile level and their low-power operations, it is possible to execute simultaneously erase of some tiles, programming at some other 65 tiles, and reading one or more of remaining tiles. The module controller can use that parallelism of operations at

the tile level to work in background mode; for example, the module controller may delete (i.e., erase, so as to free up space) some tiles or entire chips, de-fragment fragmented files into consolidated files, move files, tiles or chips that have been inactive for longer than a predetermined time to slower or archival storage, or to chips that group together files with close dates and time stamps, while rewriting the original file identifier with the latest time stamp into cache storage 116 of the next available physical tile.

42

To facilitate high-speed searches for the location of the most current version of any one file out of the many millions such files in a petabyte storage system, it is important that the unique identifier for each file, wherever it has been physically relocated to, be accessed quickly by the module controller. According to one embodiment of the present invention, a module controller broadcasts the unique identifier (i.e., the 32-128 bits word) for the file being searched simultaneously to some or all the chips in the system. Each chip is provided with a buffer memory to temporarily store that identifier and, using on-chip Exclusive-Or (XOR) circuits, compare the identifier in the buffer memory with all the identifiers stored on cache 116 of each tile and report to the module controller when a match has been found, together with the location where the corresponding file is located. If more than one match is found, the module controller picks the identifier with the most recent timestamp. The search can be narrowed to just a few chips if the file being searched has been written within a known time period. For a 1-terabit chip, just one 128-Kbit slice or 16×8 Kb pages would be sufficient to store all the 64-bit identifiers for all 2K slices of each tile.

FeFET Pairs for Fast Read Cache Memory.

To reduce read latency for cache storage 116, FeFETs in NOR memory strings that are physically nearest to sense amplifiers 206 can be arranged in pairs. For example, in adjacent NOR memory strings, two FeFETs sharing a common local word line may be used to store a single data bit between them. In a typical programming operation, FeFETs on both NOR memory strings are initialized to the erased state, then one of the FeFETs is programmed to the programmed state, while the other FeFET is program-inhibited, so as to remain in the erased state. Both FeFETs on the two adjacent active strips are read simultaneously by a differential sense amplifier in substrate circuitry when their shared local word line is raised to the read voltage, the first FeFET that start to conduct tips the sense amplifier into state '0' or state '1', depending on which FeFET is in the programmed

This FeFET-pair scheme has the advantage of high-speed sensing and higher endurance because FeFETs of two adjacent NOR memory strings are almost perfectly matched, so that at the sense amplifier even a small programmed voltage differential between the two FeFETs being read will suffice to correctly trip the sense amplifier. In other words, among other roles, the FeFET-pair acts as its own reference, neutralizing any leakage current or noise along the FeFET-pair's respective NOR memory strings. In addition, as the threshold voltage of a programmable reference FeFET may drift over many write/erase cycle over the life of the device, under this scheme the reference FeFET and the read FeFET are both reset with each new cycle. In fact, either one of the two FeFETs in the pair can serve as the reference FeFET. If the two FeFETs making the pair are randomly scrambled to invert or not invert the data written in each cycle, to ensure that statistically each FeFET in each pair serves as the reference FeFET for approximately the same number of cycles as the other FeFET. (The invert/not invert code can be

stored in the same page as the page being programmed, to assist in the descrambling during a read operation). Because the paired FeFETs are in close proximity to each other, i.e., on two adjacent active strips on the same plane, the FeFETs can best track each other for local variations in the manu- 5 facturing process or to best neutralize (i.e., cancel out) the strip leakage during a read operation. In one embodiment the first and second FeFETs of a pair share a common local word line and separate pre-charge transistors. The respective common bit lines of the first and second FeFETs of the pair function as the input or output terminals of a quasi-volatile data latch. Compared to an SRAM circuit with its low memory density and high leakage, or a DRAM circuit with its very short data retention time, a FeFET-pair in a 3-dimensional NOR memory strings of the present invention 15 have a longer data retention time, lower power dissipation, and very high memory density as data latches.

Alternatively, the FeFET pairing scheme may be applied to FeFETs on different planes where the pair shares a common vertical local word line. The cost of this scheme is 20 in silicon efficiency, as the two FeFETs are required to store one bit between them. For this reason, each tile can be organized such that only a small percentage (e.g., 1% to 10%) of the tile is used as high-speed dual FeFET pairs, while the rest of the tile is operated as regular NOR memory 25 strings and programmable reference FeFET strings. The actual percentage set aside for the FeFET-pair scheme can be altered dynamically by the system controller, depending on the specific usage application. For example, the FeFET pairs can be used to store critical system parameters values and 30 may be periodically refreshed, since they are quasi-volatile. Thus, these FeFET-pairs can maintain their polarization states for extended periods of time (e.g., hours to days) without consuming much power. In fact, such FeFET-pair implemented in 3-dimensional NOR memory arrays of the 35 current invention (which are highly silicon efficient) can be employed as high-density data caches or buffer memory, to back up or altogether replace SRAM or DRAM cache memories that are manufactured on planar (i.e., two-dimensional) silicon substrate and are therefore costly and power 40 consuming, even when they are in their standby modes. Such FeFET-pairs could be particularly useful in artificial intelligence and machine learning applications. The high level of flexibility for operating the NOR memory strings of FeFETs of the present invention result in part from the fact 45 that the FeFETs in a NOR memory string are randomly addressable and operate independently of each other, or of FeFETs in other NOR memory strings, unlike conventional NAND strings.

Numerous applications of data storage, such as video or 50 high-resolution imaging require data files that occupy many pages or even many slices. Such files can be accessed rapidly in a pipelined fashion, i.e., the system controller stores the first page or first slice of the file in the cache memory while storing the remaining pages or slices of the 55 file in a low-cost memory and streaming out the data in a pipeline sequence. The pages or slices may thus be linked into a continuous stream, such that the first page of the file is read quickly into the sense amplifiers and transferred to a data buffer shift register to clock the first page out of the tile 60 while pre-charging and reading the next, slower page in a pipeline sequence, thereby hiding the read access time of each page following the first page. For example, if the first page of 8 Kbits stored in the cache memory is read in 10.0 nanoseconds and then clocked out at 1 Gbit per second, the 65 entire 8K bits would take approximately 1 microsecond to complete clocking out, which is more than sufficient time for

44

the second page to be read from the slower, lower-cost pages. The flexibility afforded by pre-charging randomly selected FeFET strings makes it possible for one or more data files from one or more tiles to be written or read concurrently, with their data streams routed on-chip to one or more data input/output ports.

Random Access Reads

The pre-charging scheme of the current invention allows data to be programmed to be serially clocked into, or randomly accessed, and likewise read out serially in a stream or randomly accessed by words. For example, an addressed page in one plane can be read in one or more operations into the sense amplifiers, registers or latches of the addressed plane, after which it can be randomly accessed in 32-bit, 64-bit, 128-bit or 512-bit or more words, one word at a time, for routing to the input/output pads of the chip. In this manner, the delay attendant to streaming the entire page sequentially is avoided.

Reading FeFETs of the present invention is much faster than reading conventional NAND flash memory cells because, in a NOR memory string, only the FeFET to be read is required to be "on", as compared to a NAND string, in which all the FeFETs in series with the one FeFET being read must also be "on". To reduce read latency, some or all the planes in selected memory tiles may be kept pre-charged to their read voltages V_{ss} (source line) and V_{bl} (bit line), thereby rendering them ready to immediately sense the addressed FeFET (i.e., eliminating the time required for pre-charge immediately before the read operation). Such ready-standby requires very little standby power because the current required to periodically re-charge capacitor 360 to compensate for charge leakage is small. Within each tile, all NOR memory strings on all eight or more planes can be pre-charged to be ready for fast read; for example, after reading FeFETs in NOR memory strings of plane 207-0 (FIG. 2a), FeFETs in NOR memory strings of plane 207-1 can be read in short order because its source and bit line voltages V_{ss} and V_{bl} are already previously set for a read operation. The pre-charge operation of a selected FeFET in a NOR string can be performed as the final step following completion of a read operation so that it does not adversely impact the read latency.

In memory tile 100, only one FeFET per NOR memory string can be read in a single operation. In a plane with eight thousand side by side NOR memory strings, the eight thousand FeFETs that share a common global word line may all be read concurrently, provided that each NOR memory string is connected to its own sense amplifier. If each sense amplifier is shared among, for example, four NOR memory strings in the same plane using a string decode circuit, then four read operations are required to take place in four successive steps, with each read operation involving two thousand FeFETs. Each plane can be provided its own set of dedicated sense amplifiers or, alternatively one set of sense amplifiers can be shared among NOR memory strings in the eight or more planes through a plane-decoding selector. Additionally, one or more sets of sense amplifiers can be shared between NOR memory strings in quadrants and their mirror image quadrants (see, e.g., sense amplifiers (SA) 206 in FIGS. 6a, 6b, and 6c). Providing separate sense amplifiers for each plane allows concurrent read operations of NOR memory strings of all planes, which correspondingly improves the read operation throughput. However, such higher data throughput comes at the expense of greater power dissipation and the extra chip area needed for the additional sense amplifiers (unless they can be laid out in substrate 201 underneath tile 100). In practice, just one set

of sense amplifiers per stack of NOR memory strings may suffice because of the pipeline clocking or data in and out of the memory tile, so that while a first page in one plane is being transferred out of its sense amplifiers to a high speed shift register or buffer memory, the first page of the second 5 plane is being read into the second set of sense amplifiers, with the two sets sharing one set of input/output shift registers.

Parallel operations may also create excessive electrical noise through ground voltage bounces when too many 10 FeFETs are read all at once. This ground bounce is substantially suppressed in all embodiments that rely on precharging capacitor 360 to set and temporarily hold the virtual V_{ss} voltage for each active strip. In this case, source voltage V_{ss} of all NOR memory strings is not connected to 15 the chip's V_{ss} ground line, allowing any number of active strips to be sensed simultaneously without drawing charge from the chip ground supply.

Read, Program, Margin Read, Refresh and Erase Operations for Quasi-Volatile 3-Dimensional NOR Memory Strings of 20 FeFETs.

The quasi-volatile NOR memory strings or slices of the current invention may be used as alternatives to some or all DRAMs in many memory applications, e.g., the memory devices for supporting central processing unit (CPU) or 25 microprocessor operations on the main board ("motherboard") of a computer. The memory devices in those applications are typically required to be capable of fast random read access and to have very high cycle-endurance. In that capacity, the quasi-volatile NOR memory strings of FeFETs 30 of the present invention employ similar read/program/inhibit/erase sequences as the non-volatile NOR implementation. In addition, since the polarization stored on programmed FeFETs slowly depolarizes, the weakened polarization state may need to be restored by reprogramming 35 the FeFETs in advance of a read error. To avoid the read error, one may employ "margin read" conditions to determine if a program-refresh operation is required, as are well known to a person skilled in the art. Margin read is an early-detection mechanism for identifying which FeFET 40 will soon fail, to trigger a restoration to the correct programmed state. Quasi-volatile FeFETs typically may be programmed, program-inhibited or erased at reduced programming voltage (V_{pgm}) , program inhibit voltage $(V_{inhibit})$ or erase voltage (V_{erase}) or are programmed using shorter 45 pulse durations. The reduced voltages or shorter pulse durations result in a reduced stress on the ferroelectric material and, hence, improvement in endurance. All slices in a tile may require periodic reads under margin conditions to early-detect excessive threshold voltage shifts of the pro- 50 grammed FeFETs due to damage in the ferroelectric storage material. For example, the erase threshold voltage may be 0.5V±0.2 V and the programmed threshold voltage may be $1.5V \pm 0.2V$, so that a normal read voltage may be set at ~1V while the margin-read may be set at ~1.2V. Any slice that 55 requires a program-refresh needs to be read and then correctly reprogrammed into the same slice or into an erased slice in the same tile or in another previously erased tile. Multiple reads of quasi-volatile FeFETs can result in disturbing the erase or program threshold voltages, and may 60 require rewriting the slice into another, erased slice. Read disturbs are suppressed by lowering the voltages applied to the control gate, and the source and drain regions during reads. However, repetitive reads may cumulatively cause read errors. Such errors can be recovered by requiring the 65 data to be encoded with error correcting codes ("ECC"). Typically, each read operation involves decoding the

46

retrieved encoded data. The decoding process computes a syndrome, which indicates if the retrieved data has been corrupted. Based on the syndrome, the decoding process determines where the corrupted bits ("error bits") are and makes the necessary correction, if possible. In one embodiment, for a 512-bit data word, one suitable ECC scheme detects up to 6 error bits and corrects up to 5 error bits. An FeFET that has been subject to a billion cycles or more of erase or programming operations may have one or both of its polarization states drifted outside of its designated operational threshold voltage window. This condition may result in an error bit. The system controller may replace the data word associated with the FeFET with a spare word and may adjust the erase/program bias conditions to achieve an operational window for the FeFET. This same repair/replace sequence may be included within the refresh operation (typically, operating in a background mode), as every FeFET in the entire memory circuit is read by the refresh process according to its schedule. The ECC encoded word has 10% or more bits than the data word, and thus imposes an overhead cost in word size and may increase the read latency (e.g., by a few nanoseconds). However, a suitable ECC scheme greatly enhances memory reliability over the life of the memory circuit.

One requirement for the proper operation of the quasivolatile memory of the present invention is the ability to read and program-refresh a large number of FeFETs, NOR memory strings, pages or slices. For example, a quasivolatile 1-terabit chip has ~8,000,000 slices of 128K bits each. Assuming that 8 slices (~1 million) of FeFETs can be program-refreshed in parallel (e.g., one slice in each of 8 tiles), and assuming a program-refresh time of 100.0 microseconds, then program-refreshing the entire chip may be achieved in ~100 seconds. This massive parallelism is made possible in memory devices of the present invention primarily because of two key factors; 1) the programming or erase mechanism requires very low current per FeFET, allowing 1 million or more FeFETs to be programmed together without expanding excessive power; and 2) the parasitic capacitor intrinsic to a long NOR memory string enables pre-charging and temporarily holding the pre-charged voltage on multiple NOR memory strings. These characteristics allow a multitude of pages or slices on different tiles to be first read in margin-read mode (or in ECC-check mode) to determine if a refresh is required, and if so, the pages or slices are individually pre-charged for program or program-inhibit and then program-refreshed in a single parallel operation. A quasi-volatile memory with average retention time of ~10 minutes or longer will allow the system controller to have adequate time for properly program-refresh, and to maintain a low error rate that is well within the ECC recovery capability.

FIG. 10a shows in simplified form prior art storage system 800 in which microprocessor (CPU) 801 communicates with system controller 803 in a flash solid state drive (SSD) that employs NAND flash chips 804. The SSD emulates a hard disk drive and NAND flash chips 804 do not communicate directly with CPU 801 and have relatively long read latency. FIG. 10b shows in simplified form system architecture 850 using the memory devices of the current invention, in which 3-dimensional non-volatile FeFET NOR memory string arrays 854, or quasi-volatile NOR memory string arrays 855 (or both) are accessed directly by CPU 801 through one or more of input and output (I/O) ports 861. I/O ports 861 may be one or more high speed serial ports for data streaming in or out of NOR memory string arrays 854 and 855, or they may be 8-bit, 16-bit, 32-bit, 64-bit, 128-bit, or

any suitably sized wide words that are randomly accessed, one word at a time. Such access may be provided, for example, using a DRAM-compatible protocol (e.g., DDR4), and future higher speed industry standard memory interface protocols, or other protocols for DRAM, SRAM or NOR 5 flash memories. I/O ports 862 handle storage system management commands, with flash memory controller 853 translating CPU commands for memory chip management operations and for data input to be programmed into the memory chips. In addition, CPU 801 may use I/O ports 862 to write and read stored files using one of several standard formats (e.g., CXL, PCIe, HBM, NVMe, eMMC, SD, USB, SAS, or multi-Gbit high data-rate ports). I/O ports 862 communicate between system controller 853 and FeFET NOR memory string arrays in the memory chips.

It is advantageous to keep the system controller or chiplet (e.g., system controller 853 of FIG. 8b) off the memory chips, as each system controller typically manages a large number of memory chips, so that it is disengaged as much as possible from the continuous ongoing margin-read/pro- 20 gram-refresh operations, which can be more efficiently controlled by simple on-chip state machines, sequencers or dedicated microcontrollers. For example, parity-check bit (1-bit) or more powerful ECC words (typically, between a few bits to 70 bits or more) can be generated for the 25 incoming data by the off-chip controller or on-chip by dedicated logic or state machines and stored with the page or slice being programmed. During a margin-read operation the parity bit generated on-chip for the addressed page is compared with the stored parity bit. If the two bits do not 30 match, the controller reads again the addressed page under a standard read (i.e., non-margin). If that gives a parity bit match, the controller reprograms the correct data into the page, even though it is not yet fully corrupted. If the parity bits do not match, then on-chip dedicated ECC logic or the 35 off-chip controller intervenes to detect and corrects the bad bits and rewrites the correct data preferably into another available page or slice, and permanently retiring the errant page or slice. To speed up the on-chip ECC operations, it is advantageous to have on-chip Exclusive-Or, or other logic 40 circuitry to find ECC matches quickly without having to go off-chip. Alternatively, a memory chip can have one or more high-speed I/O ports dedicated for communication with the controller for ECC and other system management chores (e.g., dynamic defect management), so as not to interfere 45 with the low latency data I/O ports. As the frequency of read or program-refresh operations may vary over the life of the memory chip due to FeFET wear-out after excessive program/erase cycling, the controller may store in each tile (preferably in the high-speed cache slices) a value indicating 50 the time interval between refresh operations. Such a time interval tracks the cycle count of the tile. Additionally, the chip or the system may have a temperature monitoring circuit whose output data is used to modulate the frequency of refreshes according to the measured chip temperature.

In the example of a 1-terabit chip having only 8 tiles out of 4,000 tiles, or 0.2% or less of all tiles are being refreshed at any one time, program-refresh operations can be performed in a background mode, while all other tiles can proceed in parallel with their pre-charge, read, program and 60 erase operations. In the event of an address collision between the 0.2% and the 99.8% of tiles, the system controller arbitrates one of the accesses is more urgent. For example, the system controller can pre-empt a program-refresh, storing away temporarily the state of the interrupted 65 operation, to yield priority to a fast read, then return to complete the program-refresh.

48

In summary, in the integrated circuit memory chip of the present invention, each active strip and its multiple associated conductive word lines are architected as a single-port isolated capacitor that can be charged to pre-determined voltages which are held semi-floating (i.e., subject to charge leaking out through the string-select transistor in the substrate circuitry) during read, program, program-inhibit or erase operations. That isolated semi-floating capacitor of each active strip, coupled with the low current required to program or erase the FeFETs in a 3-dimensional NOR memory string associated with the active strip, makes it possible to program, erase or read a massive number of randomly selected tiles, sequentially or concurrently. Within the integrated circuit memory chip, the NOR memory strings of one or more of a first group of tiles are first pre-charged and then erased together, while the NOR memory strings of one or more other groups of tiles are first pre-charged and then programmed or read together. Furthermore, erasing of the first group of tiles and programming or reading of a second group of tiles can take place sequentially or concurrently. Tiles that are dormant (e.g., tiles that store archival data) are preferably held at a semi-floating state, preferably isolated from the substrate circuits after having their NOR memory strings and conductors set at ground potential. To take advantage of the massively parallel read and program bandwidths of these quasi-floating NOR memory strings, it is advantageous for the integrated circuit memory chip to incorporate therein multiple high-speed I/O ports such as high-bandwidth memory (HBM) employing dense through-silicon vias (TSVs)) or the much denser heterogeneous integration connectors (e.g., in die-to-die or die-to-wafer format). Data can be routed on-chip to and from these I/O ports, for example, to provide multiple channels for word-wide random access, or for serial data streams out of the chip (reading) or into the chip (programming or

FeFETs of the present invention has the added advantage of being selectively configurable to a depletion threshold voltage state ("depletion mode"). When a depletion mode FeFET has negligible or insignificant source-drain leakage current. A NOR memory string that is not selected for read, write or erase operation may have its common drain region and its common source region biased to a predetermined voltage difference to prevent source-drain leakage currents in the FeFETs that are not in the depletion mode. While held in this bias condition, the NOR memory string has insignificant or negligible leakage current.

In some memory applications, FeFETs may be organized as differential pairs, in which one FeFET in each differential pair holds a first binary value and the other FeFET holds in the differential pair holds the opposite binary value. The differential pair may be used as input or output terminal of a latch. The differential pair may be formed by storage transistors sharing a common gate electrode.

In this detailed description, various embodiments or examples of the present invention may be implemented in numerous ways, including as a process; an apparatus; a system; and a composition of matter. A detailed description of one or more embodiments of the invention is provided above along with accompanying figures that illustrate the principles of the invention. The invention is described in connection with such embodiments, but the invention is not limited to any embodiment. Numerous modifications and variations within the scope of the present invention are possible. The scope of the invention is limited only by the claims and the invention encompasses numerous alternatives, modifications, and equivalents. Numerous specific

details are set forth in the description in order to provide a thorough understanding of the invention. These details are provided for the purpose of example and the invention may be practiced according to the claims without some or all of these specific details. For the purpose of clarity, technical material that is known in the technical fields related to the invention has not been described in detail so that the invention is not unnecessarily obscured. The present invention is defined by the appended claims.

We claim:

- 1. A memory structure formed above a planar surface of a semiconductor substrate, organized as a 3-dimensional array of NOR memory strings, each NOR memory string 15 comprising a plurality of thin-film ferroelectric transistors (FeFETs) sharing a common drain region and a common source region, the common source region and the common drain region each extending lengthwise along a first direction substantially parallel to the planar surface of the semi- 20 conductor substrate, the common source region and the common drain region being spaced apart in a second direction substantially perpendicular to the planar surface of the semiconductor substrate, wherein each FeFET further comprises (i) a channel region in contact along the second 25 direction with both the common source region and the common drain region at a first side of the channel region and, (ii) an electrically polarizable layer in contact with the channel region at a second side, opposite the first side, of the channel region so that the channel region is provided 30 between the electrically polarizable layer and the common source region and between the electrically polarizable layer and the common drain region, and (iii) a conductive layer serving as a gate electrode provided adjacent the electrically polarizable layer.
- 2. The memory structure of claim 1, wherein the common source region, the common drain region and the channel region of each FeFET comprise crystalline silicon or polycrystalline silicon.
- 3. The memory structure of claim 1, wherein each FeFET 40 comprises a junction-less field-effect transistor without a p/n junction between the channel region and the common source region and without a p/n junction between the channel region and the common drain region.
- **4**. The memory structure of claim **3**, wherein the channel 45 region of the FeFET comprises an oxide semiconductor.
- 5. The memory structure of claim 4, wherein the oxide semiconductor comprises one or more of: indium zinc oxide and indium gallium zinc oxide.
- **6**. The memory structure of claim **4**, wherein the common 50 drain region and the common source region each comprise a metal.
- 7. The memory structure of claim 6, wherein the metal comprises one or more of: molybdenum, titanium, tungsten, lanthanum, tantalum, ruthenium, any silicide thereof, any 55 nitride thereof and any combination thereof.
- 8. The memory structure of claim 1, wherein the polarizable layer comprises one or more of: laminated layers of hafnium oxide (HfO2) and zirconium oxide (ZrO2), zirconium-doped hafnium oxide (HfO:Zr), aluminum zirconium-doped hafnium oxide (HfZrAlO), aluminum-doped hafnium oxide (HfO2:Si), lanthanum-doped hafnium oxide (HfO2:La), hafnium zirconium oxide (HfZrON), hafnium zirconium oxide (HfZrON), hafnium zirconium oxide (HfZrON), hafnium zirconium impurities, and combinations thereof.

50

- 9. The memory structure of claim 4, wherein the oxide semiconductor has a charge carrier mobility exceeding 4.0 cm²/V.
- **10**. The memory structure of claim **4**, wherein the oxide semiconductor in the channel region is formed by deposition at a temperature between 200.0-450.0° C.
- 11. The memory structure of claim 1, wherein the polarizable layers of the FeFETs are electrically isolated from each other.
- 12. The memory structure of claim 1, wherein each FeFET further comprises an interface dielectric layer between the channel region and the polarizable layer, the interface dielectric layer having a dielectric constant greater than stoichiometric silicon oxide's dielectric constant.
- 13. The memory structure of claim 12, wherein the interface dielectric layer comprises one or more of a silicon nitride layer and an aluminum oxide layer.
- **14**. The memory structure of claim **12**, wherein the interface dielectric layer is formed by nitridation of a silicon oxide or silicon oxynitride.
- 15. The memory structure of claim 12, wherein the interface dielectric layer is converted by nitridation from an oxide formed on the surface of the channel region.
- **16**. The memory structure of claim **15**, wherein the nitridation is carried out using ammonia anneal.
- 17. The memory structure of claim 12, wherein the interface dielectric layer has a dielectric constant greater than 3.9.
- 18. The memory structure of claim 12, wherein the interface dielectric layer is treated using one or more of: pulsed ozone and thermal annealing in a hydrogen or deuterium ambient.
- 19. The memory structure of claim 12, wherein the interface dielectric layer comprises one or more of: aluminum oxide, zirconium oxide (ZrO2), silicon oxynitride (SiON), silicon nitride (Si3N4), and silicon oxide (SiO2).
 - 20. The memory structure of claim 1, wherein the 3-dimensional array comprises stacks of NOR memory strings wherein, within each stack, the NOR memory strings are provided at successively greater predetermined distances along the second direction from the planar surface of the semiconductor substrate.
 - 21. The memory structure of claim 20, wherein each stack of NOR memory strings has a predetermined number of NOR memory strings.
 - 22. The memory structure of claim 20, wherein the stacks are arranged in rows substantially along a third direction parallel to the planar surface of the semiconductor substrate and orthogonal to both the first and second directions.
 - 23. The memory structure of claim 22, wherein the conductive layers of a first group of FeFETs in the stacks of NOR memory strings are connected to a first conductor extending along the third direction, each FeFET belonging to a different NOR memory string within the stacks.
 - 24. The memory structure of claim 23, wherein the conductive layers of a second group of FeFETs provided on an opposite side of the same stacks are connected to a second conductor extending along the third direction, each FeFET belonging to a different NOR memory string within the stacks.
 - 25. The memory structure of claim 24, wherein the first conductor and the second conductor are offset from each other along the first direction.
 - **26**. The memory structure of claim **1**, wherein current flows in the channel region of each FeFET substantially along the second direction.

- **27**. A method for operating a selected FeFET within the memory structure of claim **24**, comprising, during a program or erase operation:
 - (i) applying a first voltage on the gate electrode associated with the selected FeFET, and applying a second voltage on both the common source region and the common drain region of the selected FeFET, wherein the first and second voltages each have a magnitude difference sufficient to carrying out programming or erasing on the selected FeFET;
 - (ii) for each FeFET sharing the first conductor with the selected FeFET, applying on the common drain region of the FeFET a third voltage such that the first voltage and the third voltage is different in magnitude by either one-half the magnitude difference of the selected 15 FeFET, or one-third the magnitude difference of the selected FeFET; and
 - (ii) for each FeFET sharing a common drain region with the selected FeFET, applying on the common gate electrode associated with the FeFET a fourth voltage 20 such that the second voltage and the fourth voltage is different in magnitude by either one-half the magnitude difference of the selected FeFET or one-third the magnitude difference of the selected FeFET.
- **28**. The method of claim **27**, wherein (i) the selected 25 FeFET has a threshold voltage that represents either a programmed state or an erased state, and (ii) when the selected FeFET is in a programmed state, the threshold voltage is a negative value, the method further comprising, during a read operation, applying a boost voltage to the 30 common source region of the selected FeFET that is greater than the negative value.
- **29**. The method of claim **28**, further comprising, during a program or erase operation on the selected FeFET, refreshing one or more FeFETs that share a common gate electrode 35 of the selected FeFET.
- **30**. The memory structure of claim **20**, wherein the NOR memory strings within each stack are isolated one from another by an isolation layer.
- **31**. The memory structure of claim **30**, wherein the 40 polarizable layer is deposited using a selective deposition technique that prevents the polarizable layer from being deposited on the isolation layer.
- **32**. The memory structure of claim **31**, wherein the selective deposition is carried out in an ozone ambient at 45 temperature between 200.0° C. to 350.0° C.
- 33. The memory structure of claim 31, wherein the selective deposition includes a post-deposition annealing step carried out at a temperature between 400° C. and 1000° C.
- **34**. The memory structure of claim **31**, wherein the isolation layer comprises silicon oxycarbide (SiOC).
- **35**. The memory structure of claim **1**, wherein each NOR memory string further comprises one or more pre-charge transistors, each configured, when activated, for pre-charging the common source region to a voltage provided on the common drain region.
- **36**. The memory structure of claim **1**, wherein circuitry for supporting memory operation is formed at the planar surface of the semiconductor substrate underneath the FeFETs.
- **37**. The memory structure of claim **36**, wherein the circuitry for supporting memory operations includes both analog and digital circuits.
- **38**. The memory structure of claim **36**, further comprising a layer of interconnect conductors formed above, and in 65 electrically connection with, the array of NOR memory strings, the layer of interconnect conductors being provided

52

for routing control and data signals among the NOR memory strings and the circuitry for supporting memory operations.

- 39. The memory structure of claim 38, wherein the circuitry for supporting memory operations comprises two or more of: shift registers, latches, sense amplifiers, reference cells, power supply lines, bias and reference voltage generators, inverters, NAND, NOR, Exclusive-Or and other logic gates, input/output drivers, address decoders, other memory elements, sequencers, state machines, and data processing circuitry that receives control and data signal from an external control circuit.
- **40**. The memory structure of claim **39**, wherein the external control circuit is formed on a separate semiconductor substrate.
- **41**. The memory structure of claim **39**, wherein the semiconductor substrates are provided in an integrated circuit package.
- **42**. The memory structure of claim **40**, wherein the semiconductor substrates are wafer-bonded.
- 43. The memory structure of claim 36, further comprising interconnect conductors formed between the array of NOR memory strings and the planar surface, the interconnect conductors electrically connect among the support circuitry, the common drain regions of the FeFETs transistors, and the conductive layers of the FeFETs.
- **44**. The memory structure of claim **1**, wherein each FeFET, upon application of a first voltage difference between the gate electrode and the common source region or the common drain region, is electrically placed in a first of two polarization states and, upon application of a second voltage difference between the gate electrode and the common source region or the common drain region is electrically placed in a second of the two polarization states.
- **45**. The memory structure of claim 1, wherein the polarizable layer contacts the channel region directly.
- **46**. The memory structure of claim **1**, wherein the polarizable layer is deposited using chemical or physical vapor deposition, ALD, or evaporation.
- **47**. The memory structure of claim **1**, wherein a portion of the polarizable layer is provided in a recess between the common source region and the common drain region.
- **48**. The memory structure of claim **47**, wherein the polarizable layer is formed using a sacrificial material that protects a first portion of the polarizable layer inside the recess but allows a second portion of the polarizable layer to be removed from the isolation layers.
- 49. The memory structure of claim 47, wherein, within the for recess, the ferroelectric material is embedded in a dielectric material.
 - **50**. The memory structure of claim **49**, wherein the dielectric material comprises silicon oxide.
 - **51**. The memory structure of claim **47**, wherein the polarizable layer is embedded with a dielectric material within the recess.
 - **52**. The memory structure of claim **51**, wherein the dielectric material comprises silicon oxide.
- 53. The memory structure of claim 1, wherein the channel region is formed subsequent to forming the polarizable layer.
 - **54**. The memory structure of claim 1, further comprising a metal layer adjacent and in electrical contact with the common drain region.
 - **55**. The memory structure of claim **1**, wherein the FeFET polarization for erase and program states are achieved with single voltage polarity.

- **56**. The memory structure of claim **1**, wherein programming of each FeFET takes advantage of a gate-induced drain leakage (GIDL) effect.
- 57. The memory structure of claim 1, wherein programming of each FeFET comprises applying a voltage differsence between the common drain region and the common source region, thereby generating charge carriers in the channel region.
- **58**. The memory structure of claim **57**, wherein the voltage difference is between 0.5 volts and 2.0 volts.
- **59**. The memory structure of claim **1**, wherein each FeFET is selectively configurable to a depletion threshold voltage state.
- **60**. The memory structure of claim **59**, wherein one of the FeFET in one of the NOR memory strings is configured to 15 the depletion threshold voltage state, and wherein, during operation, the common drain region and the common source region of that FeFET are pre-charged respectively to a first predetermined voltage and a second predetermined voltage.
- **61**. The memory structure of claim **60**, wherein the 20 FeFETs in the NOR memory string, other than the FeFET in the depletion threshold voltage state, each have a substantially zero conduction current between their respective common source regions and common drain regions during the operation.
- **62**. The memory structure of claim **61**, wherein the FeFET in the depletion threshold voltage state has a substantially zero source-drain leakage current.
- **63**. The memory structure of claim **1**, wherein a first FeFET of a first NOR memory string and a second FeFET

54

of a second NOR memory string operate as a differential pair, in which the first FeFET holds a first binary value and the second FeFET holds a binary value different from the first binary value.

- **64**. The memory structure of claim **63**, wherein the gate electrode of the first FeFET and the gate electrode of the second FeFET are electrically connected.
- **65**. The memory structure of claim **63**, wherein the common drain region of the first FeFET and the common drain region of the second FeFET function as input or output terminals of a latch.
- 66. A memory structure formed above a planar surface of a semiconductor substrate, organized as a 3-dimensional array of NOR memory strings, each NOR memory string comprising a plurality of thin-film ferroelectric transistors (FeFETs) sharing a common drain region and a common source region, the common source region and the common drain region each extending lengthwise along a first direction substantially parallel the planar surface of the semiconductor substrate, wherein each FeFET further comprises (i) a channel region in contact with both the common source region and the common drain region, (ii) a conductive layer serving as a gate electrode, and (iii) an electrically polarizable layer in contact with the channel region, and wherein each NOR memory string further comprises one or more pre-charge transistors, each configured, when activated, for pre-charging the common source region to a voltage provided on the common drain region.

* * * * *