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Jeddeloh

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(54) **DATA TRANSFER MANAGEMENT**

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(51) **Int. Cl.**

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G06F 13/28	(2006.01)
G06F 12/00	(2006.01)
G06F 13/14	(2006.01)
G06F 13/16	(2006.01)
G06F 3/06	(2006.01)

(52) **U.S. Cl.**

CPC ***G06F 13/28*** (2013.01); ***G06F 12/00*** (2013.01); ***G06F 13/14*** (2013.01); ***G06F 13/16*** (2013.01)

(58) **Field of Classification Search**

None

See application file for complete search history.

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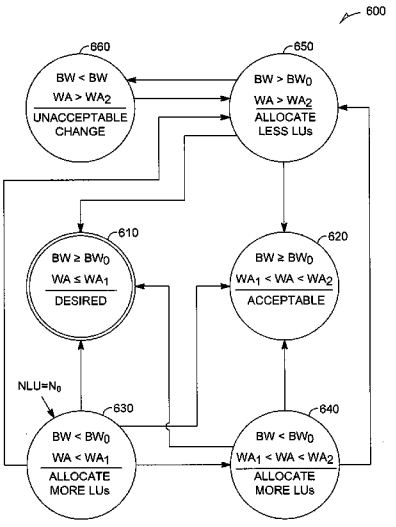
(57) **ABSTRACT**

Methods, controllers, and systems for managing data transfer, such as those in solid state drives (SSDs), are described. In some embodiments, the data transfer between a host and a memory is monitored and then assessed to provide an assessment result. A number of storage units of the memory allocated to service another data transfer is adjusted based on the assessment result. Additional methods and systems are also described.

20 Claims, 7 Drawing Sheets



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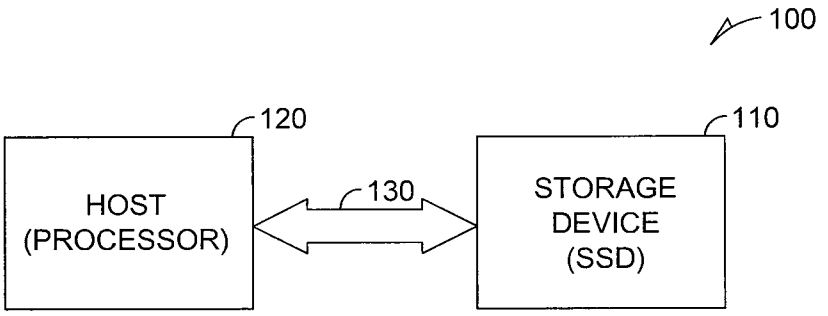
100
120
110
HOST
(PROCESSOR)
130
STORAGE
DEVICE
(SSD)

FIG. 1

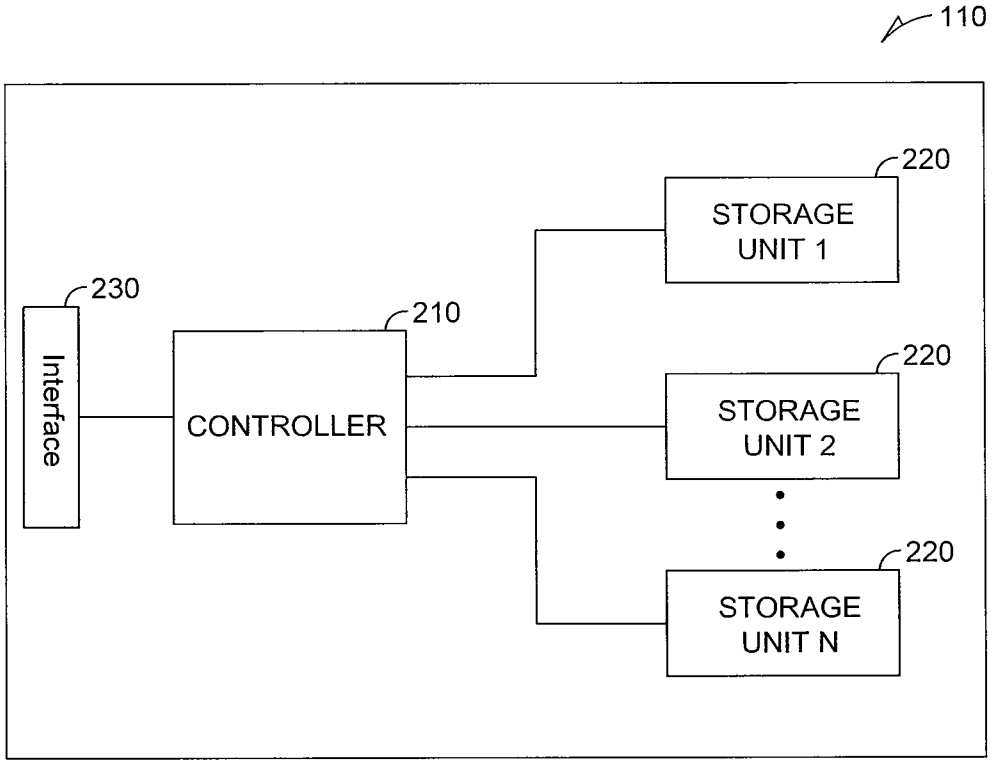
110
220
STORAGE
UNIT 1
230
210
Interface
CONTROLLER
220
STORAGE
UNIT 2
220
STORAGE
UNIT N

FIG. 2

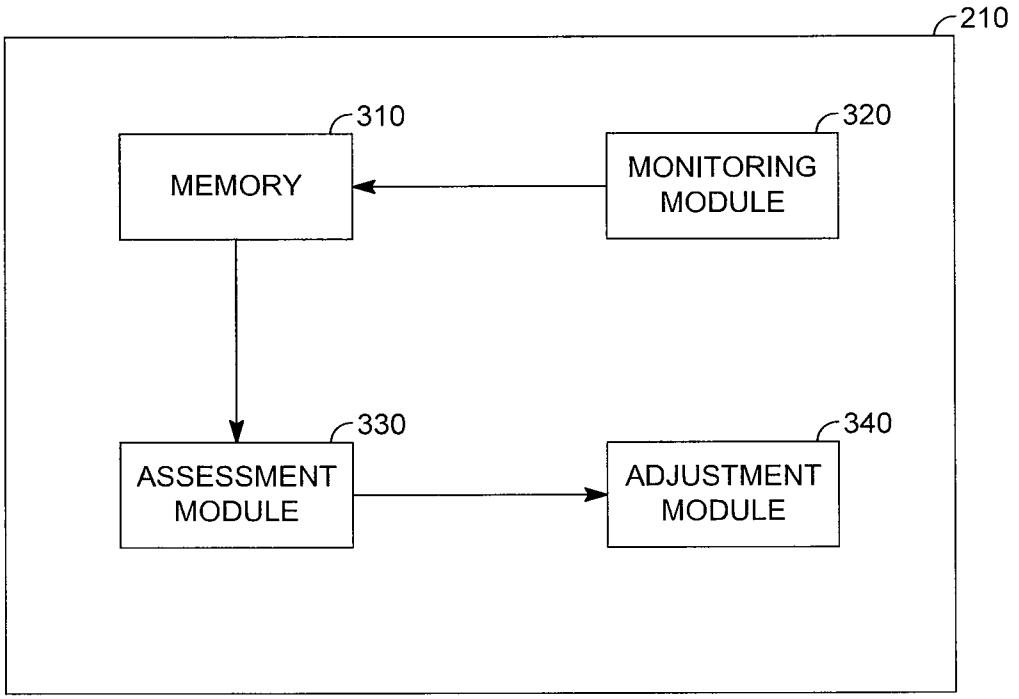
210
310
MEMORY
320
MONITORING MODULE
330
ASSESSMENT MODULE
340
ADJUSTMENT MODULE

FIG. 3

400

LU SIZE = 128KB

420 MEMORY REQUEST SIZE (KB)	430 WA (LUs=2)	440 WA (LUs=4)	450 WA (LUs=6)	460 WA (LUs=10)
2	128.0	256.0	384.0	640.0
4	64.0	128.0	192.0	320.0
20	12.8	25.6	38.4	64.0
50	5.1	10.2	15.4	25.6
120	2.1	4.3	6.4	10.7
180	1.4	2.8	4.3	7.1
200	1.3	2.6	3.8	6.4
240	1.1	2.1	3.2	5.3
500	0.5	1.0	1.5	2.6

FIG. 4

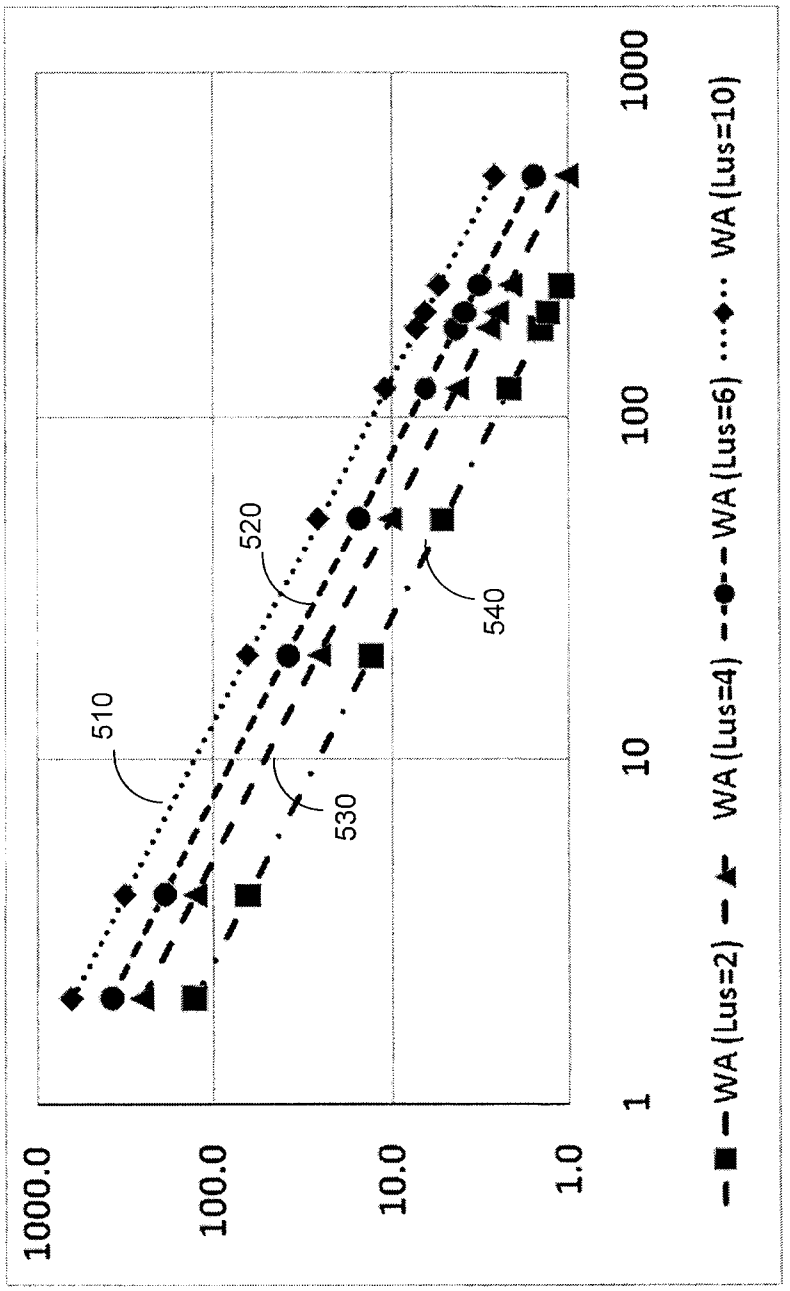
500
1000.0
100.0
10.0
1.0
1
10
100
1000
510
520
530
540
WA (Lus=2)
WA (Lus=4)
WA (Lus=6)
WA (Lus=10)

FIG. 5

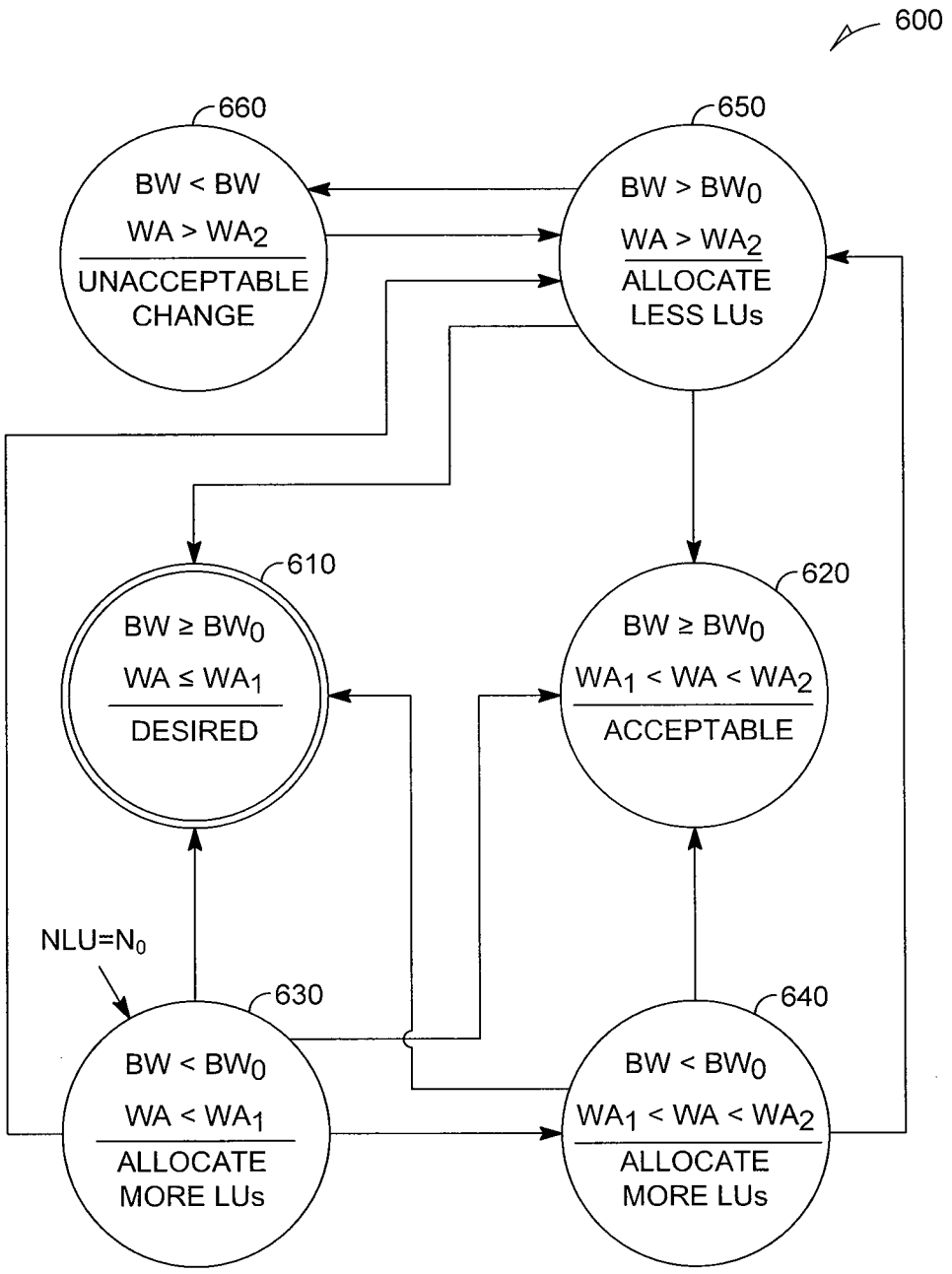
600
660
BW < BW
WA > WA2
UNACCEPTABLE CHANGE
650
BW > BW0
WA > WA2
ALLOCATE LESS LUs
610
BW ≥ BW0
WA ≤ WA1
DESIRED
620
BW ≥ BW0
WA1 < WA < WA2
ACCEPTABLE
NLU=N0
630
BW < BW0
WA < WA1
ALLOCATE MORE LUs
640
BW < BW0
WA1 < WA < WA2
ALLOCATE MORE LUs

FIG. 6

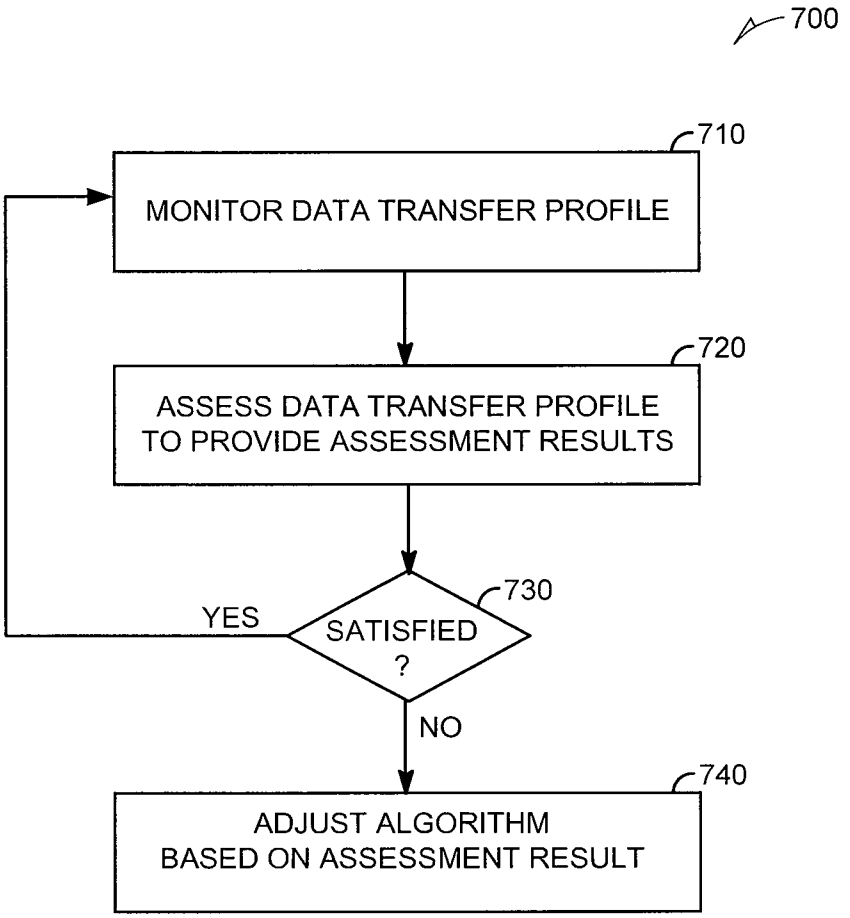
700
710
MONITOR DATA TRANSFER PROFILE
720
ASSESS DATA TRANSFER PROFILE TO PROVIDE ASSESSMENT RESULTS
730
SATISFIED ?
YES
NO
740
ADJUST ALGORITHM BASED ON ASSESSMENT RESULT

FIG. 7

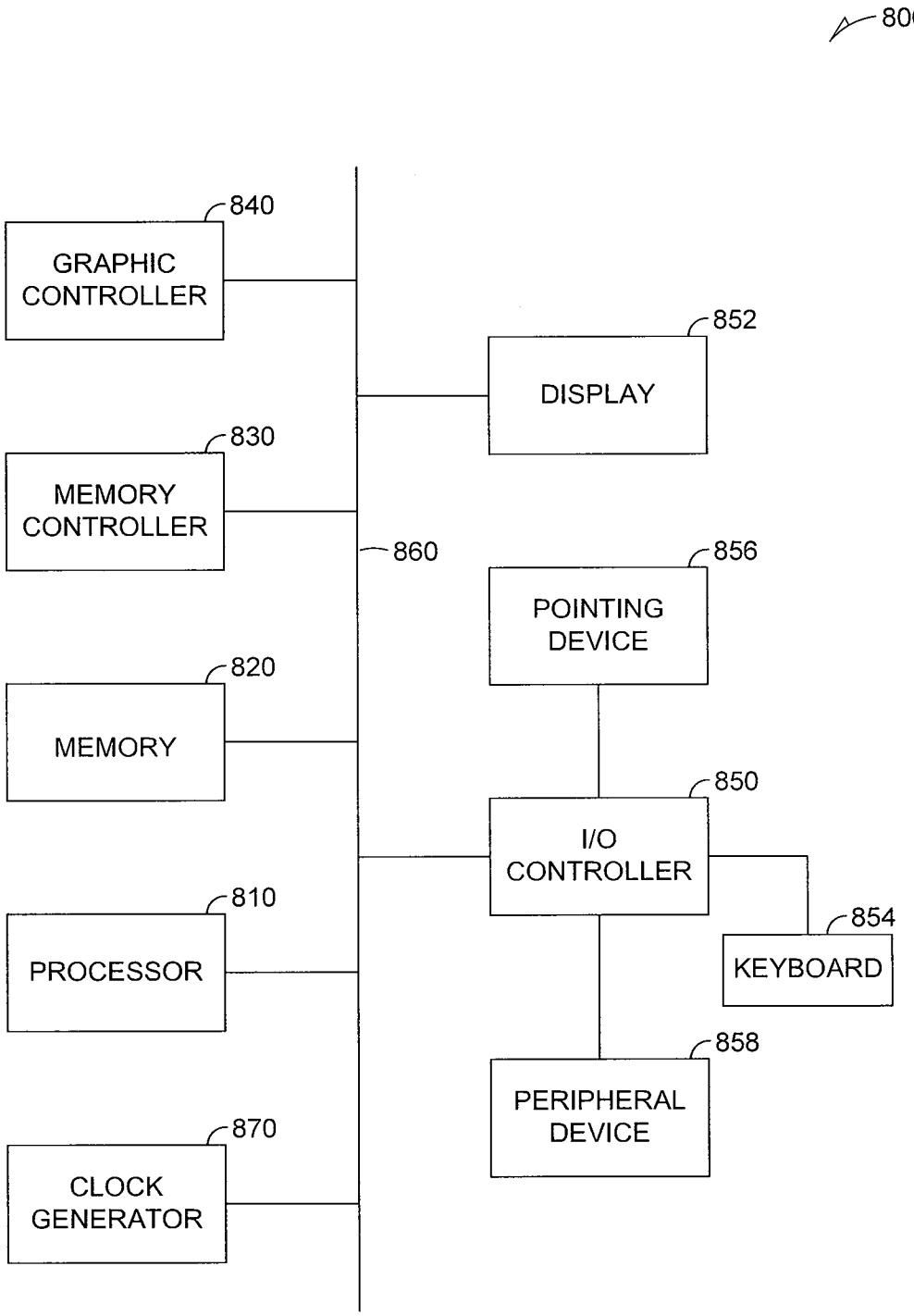
800
840
GRAPHIC CONTROLLER
852
DISPLAY
830
MEMORY CONTROLLER
860
856
POINTING DEVICE
820
MEMORY
850
I/O CONTROLLER
810
PROCESSOR
854
KEYBOARD
858
PERIPHERAL DEVICE
870
CLOCK GENERATOR

FIG. 8

DATA TRANSFER MANAGEMENT

PRIORITY APPLICATION

This application is a continuation of U.S. application Ser. No. 13/651,834, filed Oct. 15, 2012, which is a continuation of U.S. application Ser. No. 12/498,151, filed Jul. 6, 2009, now issued as U.S. Pat. No. 8,291,131 both of which are incorporated herein by reference in their entirety.

BACKGROUND

Electronic devices commonly have some type of memory system, such as a bulk storage device, available to them. A common example is a hard disk drive (HDD). HDDs are capable of large amounts of storage at relatively low cost, with current consumer HDDs available with over one terabyte of capacity. HDDs generally store data on rotating magnetic media or platters. In practice, the resulting data signal is an analog signal whose peaks and valleys are a consequence of the magnetic flux reversals of the data pattern.

HDDs have certain drawbacks due to their mechanical nature. HDDs are susceptible to damage or excessive read write errors due to shock, vibration or strong magnetic fields. In addition, they are relatively large users of power in portable electronic devices. Another example of a bulk storage device is a solid state drive (SSD). Instead of storing data on rotating media, SSDs utilize semiconductor memory devices to store their data, often including an interface and form factor that makes them appear to their host system as if they are a typical HDD. The memory devices of SSDs may comprise non-volatile flash memory devices.

BRIEF DESCRIPTION OF THE DRAWINGS

Some embodiments of the disclosed technology are illustrated by way of example and not limitation in the FIGS. of the accompanying drawings in which:

FIG. **1** is a diagram illustrating a system for managing data writes to a solid state drive (SSD), according to various embodiments of the invention;

FIG. **2** is a diagram illustrating an SSD configured to manage data writes, according to various embodiments of the invention;

FIG. **3** is a diagram illustrating modules of an SSD controller configured to manage data writes, according to various embodiments of the invention;

FIG. **4** is a table illustrating lists of some variables involved in managing data writes to an SSD, according to various embodiments of the invention;

FIG. **5** is a graph illustrating changes in the variables of the table shown in FIG. **4**, according to various embodiments of the invention;

FIG. **6** is a state diagram illustrating implementation of an algorithm for SSD data write management, according to various embodiments of the invention;

FIG. **7** is a flow diagram illustrating a method of SSD data write management, according to various embodiments of the invention; and

FIG. **8** is a diagram illustrating an example system for managing SSD data writes, according to various embodiments of the invention.

DETAILED DESCRIPTION

Example methods and systems for managing data transfers will be described. In the following description for the purpose of explanation, numerous examples having example-specific details are set forth to provide an understanding of example embodiments. It will be evident, however, to one skilled in the art that the present examples may be practiced without these example-specific details, and/or with different combinations of the details than are given here.

Some example embodiments described herein may include monitoring a data transfer (e.g., a data write) between a host and a memory. The data transfer profile may be assessed and an assessment result may be provided. A future data transfer may be adjusted based on the assessment result. For the purposes of this document, a data transfer profile comprises, for example, one or more values associated with, for example, the sizes of the data transfers (e.g., number of bytes transferred), write amplification value, and write bandwidth, as will be described in more detail below.

FIG. **1** is a diagram illustrating a system **100** for managing data writes to a solid state drive (SSD) **110**, according to various embodiments of the invention. The system **100** may include a SSD **110** in communication with a memory access device such as a processor **120**. The system **100** may be considered a host system of the SSD **110** in that it controls the operation of the SSD **110** through, for example, the processor **120**. The system **100** may employ one or more applications (host applications) to access the SSD **110**. Some examples of system **100** may include personal computers, laptop computers, personal digital assistants (PDAs), digital cameras, electronic games, digital media player/recorders, and the like.

The processor **120** may comprise a disk drive controller or other external processor. The processor **120** may communicate with the SSD **110** via a communication bus **130**. The communication bus **130** may employ a known protocol to connect the processor **120** to the SSD **110**. The type of communication bus **130** may depend on the type of drive interface being utilized in the system **100**. Examples of some conventional disk drive interface bus protocols are Integrated Drive Electronics (IDE), Advanced Technology Attachment (ATA), Serial ATA (SATA), Parallel ATA (PATA), Fiber Channel and Small Computer System Interface (SCSI). Other drive interfaces exist and are known in the art. The memory devices of the SSD **110** may include non-volatile flash memory devices. A more detailed description of features of the SSD **110**, as related to data write management, will be given below.

FIG. **2** is a diagram illustrating the SSD **110** configured to manage data writes, according to various embodiments of the invention. The SSD **110** may include a controller **210** coupled to an interface **230** that allows communication with a host (e.g., the processor **120** of FIG. **1**, or a host application) via the communication bus **130** (FIG. **1**). The interface **230** may be one of many connectors commonly known to those of ordinary skill in the art. Some example of theses interface **230** connectors may include IDE, enhanced IDE, ATA, SATA, and Personal Computer Memory Card International Association (PCMCIA) connectors.

The memory in the SSD **110** may be organized into multiple storage units **220** each comprising one or more memory devices, such as flash memory devices. Memory locations in memory devices may be grouped into blocks. In flash memory devices, blocks having storage capacities of 128 kilobytes (KB), 256 KB and 512 KB are common. The host may not be aware of which flash memory devices are currently being addressed for data read/write operations. In other words, the host is not aware of any physical block address (PBA) that is accessed in an actual data transfer (e.g., read or write). The host merely deals with a logical block address (LBA), which is translated into a PBA. This translation from

LBA to PBA may occur within the SSD **110**, after the SSD **110** receives a request for data transfer specifying a logical address.

In flash memory devices, before writing new data, a block of memory is erased to make room for the new data. For example, when the host requests a data write of only 4 KB, a block of memory (e.g., 128 KB) will be erased to accommodate the 4 KB of data. If 2 storage units, each comprising blocks that each have a storage capacity of 128 KB, are allocated to service the request, then two blocks (one from each of the 2 allocated storage units) will be erased to accommodate the 4 KB of data

The erase time may constitute a primary factor in slowing the access time of flash memory devices. In conventional memory devices, data can be written to a memory location within a flash memory a finite number of times. Even though the allowable number of writes is large (typically 10,100 to 100,000), writing to the same location over and over may wear out that location. Therefore, it is prudent to spread data writes evenly across available memory locations to level the amount of wear. Wear leveling is considered one of the tasks that the controller **210** (also called SSD controller) is responsible for managing.

With data being stored in multiple locations in the SSD **110**, and hidden from the host **120**, a large amount of manipulation may be performed to write and retrieve data from the SSD **110**. As data becomes increasingly fragmented, blocks of memory of sufficient size to write data properly may become scarce. At this point, data reclamation (sometime referred to by those of ordinary skill in the art as garbage collection) procedures are used to reclaim blocks of memory in order to allow additional writes to the SSD **110**. The more often data is physically moved within the SDD **110**, the faster the drive can wear out due to program erase cycles.

In an example embodiment, the host may specify a data transfer bandwidth (e.g., a write bandwidth, such as a lowest acceptable write bandwidth, or a read bandwidth) to be considered by the controller **210**, when servicing a host data transfer request. The term "data transfer bandwidth" shall be taken to include a quantity of a data transferred per unit time, commonly expressed in kilobits per second (Kb/s) or megabits per second (Mb/s). A larger data transfer bandwidth can be achieved by allocating more storage units **220** to service a data transfer request, for example, establishing the number of storage units that may have a block used to service the data transfer request.

Another variable of interest to the controller **210** is write amplification (WA). The term "write amplification" refers to, for example, the ratio of the size of memory allocated to be used to service a data write to the size of the actual data written to the memory allocated to service that request. For example, writing 4 KB of data to the SSD **110** may result in the controller **210** servicing a write operation with the lowest usable number of storage units—one storage unit of the memory, which can include a block that can store, for example, 128 KB. In this case of the memory locations in a 128 KB data block of the one allocated storage unit **220** are erased, even though only 4 KB of actual data that make up the write request is written to that block. This results in a WA of 128/4, i.e., 32, which may not be desirable.

Operating by the SSD **110** with lower WA values may result in a more efficient use of the memory devices in the SSD **110**. Operating at lower WA values may also result in more evenly-distributed wear leveling but, at the same time, lower write bandwidths. Therefore, the SSD **110** may be faced with a tradeoff between the write bandwidth specified by the host and the WA resulting from the data write.

The SSD **110** can be configured to address this trade off. For example, the SSD **110** may employ the controller **210** to address the tradeoff by adjusting the allocation of storage units **220**, based on the embodiments described hereinafter.

FIG. **3** is a diagram illustrating modules of an SSD controller **210** to manage data writes, according to various embodiments of the invention. The SSD controller **210** shown in FIG. **3** may, for example, comprise one or more microprocessors. The SSD controller **210** may include a memory **310**, a monitoring module **320**, an assessment module **330**, and an adjustment module **340**. The monitoring module **320**, the assessment module **330**, and the adjustment module **340** may comprise software modules stored in memory **310** or the storage units **220** in FIG. **2**. In an example embodiment, these modules may be implemented as embedded hardware in the SSD controller **210**.

The monitoring module **320** may operate to monitor a data transfer between a host (e.g., the processor **120** of FIG. **1** or a host application) and an memory, such as the memory organized into the storage units **220** of FIG. **2**. For example, the monitoring module **320** may have the capability to monitor one or more characteristics of a data transfer, such as the sizes of the data transfers (e.g., number of bytes transferred), write amplification (WA associated with the data transfer), and write bandwidth (BW) associated with the data transfer. In example embodiments, other characteristics such as reclamation bandwidth (e.g., a processor bandwidth reserved for garbage collection) and read bandwidth may also be monitored.

The monitoring module **320** may store one or more values associated with the monitored characteristic(s) in the memory **310**. The memory **310** may include a static random access memory (SRAM) used in the form of one or more configuration registers. In an example embodiment, the monitoring module **320** may store the characteristic values in flash memory devices organized into the storage units **220** of FIG. **2**. The stored characteristic values may be associated with the host.

In an example embodiment, the assessment module **330** may comprise a processor (e.g. a microprocessor). In some embodiments, it is likely the monitoring module **320** and/or the adjustment module **340** also comprise the same processor. The assessment module **330** may operate to retrieve the characteristic values, stored over a period of time, from the memory **310**. The assessment module **330** may determine statistical values based at least in part on characteristic values. For example, the assessment module **330** may determine an average value of the WA or the write bandwidth stored over a period of time. The assessment module **330** may also determine an average value of a memory request size over a period of time.

The assessment module **330** may compare one or more of the characteristic values and/or the statistical values (collectively and/or individually referred to herein as "the data transfer profile" with one or more target values. The assessment module **330** may, for example, compare the average WA with a target WA and report the results to the adjustment module **340**. In an example embodiment, the assessment module **330** may indicate that a characteristic value (e.g., WA) shows an improvement over previous values.

The adjustment module **340** may implement an algorithm such as one that evaluates the relationship between two or more of the characteristic values and/or statistical values, such as to adjust the number of storage units allocated to service a future data transfer, as described in more detail below, with respect to FIG. **5**.

Additional or different modules could be implemented within the SSD **210** as are known in the art without departing

from the scope of the disclosure. However, the present disclosure may be understood without such additional modules.

FIG. **4** is a table **400** of some variables involved in managing data writes to the SSD **110** of FIG. **1**, according various embodiments of the invention. The variables in table **400** are memory request size in KBs, WA, and the number of storage units (i.e., LUs) allocated to service a data transfer request. Column **420** lists sizes of a number of write requests made by a host (e.g., processor **120** in FIG. **1** or a host application). The controller **210** in FIG. **2** may respond to each request by allocating a number of storage units for the data write. Columns **430**, **440**, **450**, and **460** represent corresponding WA values resulting from allocation of 2, 4, 6, and 10 storage units (i.e., one block from each 2, 4, 6, and 10 LUs), to the data write requests, respectively. In the table **400**, each block of storage units are assumed to have 128 KB of storage capacity. Other capacities may be used. The data in table **400** are also represented in the graph shown in FIG. **5**.

FIG. **5** is a graph **500** illustrating changes in the variables of the table **400** shown in FIG. **4**, according to various embodiments of the invention. The graph **500** shows the variation of WA as a function of data write request sizes for different numbers of the storage units allocated to service the data write requests. The corresponding number of allocated storage units for each of the curves **510**, **520**, **530**, and **540** are shown in the legend of the graph as LUs=2, 4, 6, and 10. The general trend, as seen from the graph **500**, is a decrease in WA as the size of the requested data writes increases. For small data write request sizes (e.g., less than 128 KB), the controller **210** may have no option other than allocating a single storage unit (where a single block from a single storage unit is used to service the request). However, as the size of the requested data writes increases, the controller **210** may have more options.

For example, if the size of a requested data write is equal to 512 KB, the controller **210** in FIG. **2** may allocate 4 storage units, in which case the write amplification would amount to a value of 1. The controller may achieve a write amplification of 2 if it chooses to allocate 8 storage units for that data write request. With the assumed storage capacity of 128 KB for each storage unit, the desired WA values of 1-2 are shown to occur at the data write request sizes of more than 128 KB.

FIG. **6** is a state diagram **600** illustrating implementation of an algorithm for SSD data write management, according to various embodiments of the invention. The adjustment module **340** in FIG. **3** may implement the algorithm shown in the state diagram **600**. The state diagram **600** shows a number of states and transitions between the states. The states may be defined by variables such as write BW and WA associated with data writes by the controller **210** of FIG. **2**, to the SSD **110** of FIG. **1**. In the desired state **610**, the write bandwidth is considered to be greater than BW0, and the desired write amplification is considered to be less than WA1. The lower limit for the write bandwidth BW0 or an acceptable range for the WA (e.g., between WA1 and WA2) maybe also be specified by the host. For example, a user of the host may specify a desired write throughput target, such as 70% reads, 30% of writes and a minimum write BW of 500 Mb/s. The user may store the target values into the SSD controller **210** (e.g. into the memory **310** of FIG. **3**).

In the acceptable state **620**, the write bandwidth is greater than BW0 and the WA is within an acceptable range (e.g., less than WA2 and greater than WA1). The objective of the algorithm is to make transitions from all other states to the desired state **610** or the acceptable state **620**. For example, the initial state considered by the adjustment module **340** in FIG. **3** may be considered to be a state **630**. Before entering this state, the controller **210** in FIG. **2** has serviced a data write request by allocating N_0 storage units on the SSD **110** of FIG. **1** to a data write request. The result of such allocation, as monitored by the monitoring module **320** in FIG. **3**, is a write bandwidth less than BW0, and the WA less than WA1. Based on this scenario, adjustment module **340** may adjust the situation.

The adjustment module **340** may try to adjust the situation by allocating more storage units to the data write request. This may cause a transition to the desired state **610**, the acceptable state **620**, or another state, such as state **640**. At state **640**, the WA is within the acceptable range; however, the write bandwidth is less than BW0. To increase the write bandwidth, the adjustment module **340** may try to remedy the situation by allocating still more storage units. The allocation of more storage units may result in entering the desired state **610**, acceptable state **620**, or state **650**.

State **650** seems to be the result of allocating more than a sufficient number of storage units, because while the bandwidth has increased to an acceptable value, the WA is outside of the acceptable range. The situation may cause the adjustment module **340** in FIG. **3** to decrease the number of storage unit allocated that, in turn, may lead to entering either of the desired state **610** or the acceptable state **620**. It may also lead to entering state **660**, which may be unacceptable, because neither the write bandwidth nor the write amplification is within the acceptable range specified by state **620**. Therefore, the next transition from state **660** would be to go return to state **650** and try another value (e.g., greater than the previous tried values) for the number of storage units to be allocated.

In example embodiments, the adjustment module **340** may store transition information such as the adjustment made to the number of storage units and the monitored results (e.g., write BW, and WA) after each transition in memory **310** of FIG. **3**. The adjustment module **340** may use the transition information to improve its efficiency by trying to avoid intermediate states and directly transitioning to the desired or acceptable states **610** and **620**, respectively, in FIG. **6**. In cases where transitioning to the desired or acceptable states **510** and **520**, respectively, is not possible, the adjustment module **340** may try to lower the WA value to an even greater extent.

FIG. **7** is a flow diagram illustrating a method **700** of SSD data write management, according to various embodiments of the invention. At operations **710** the monitoring module **320** of FIG. **3** may monitor a data transfer between the host and a memory such as a memory organized into storage units **220** of the SSD **110** of FIG. **1**. The monitoring module **320** may monitor the data transfer over a time period. The time period may depend on the volume of traffic associated with the data transfer. The assessments module **330** of FIG. **3**, at operations **720**, may assess the data transfer to provide an assessment result.

The assessment result may comprise a characteristic value including a WA, a data transfer BW, and/or a number of storage units. In some example embodiments, the characteristic value may also or alternatively include a reclamation bandwidth and a read bandwidth. The assessment module **330** may determine a statistical based at least in part on a monitored characteristic value such as an average value of the characteristic value, for example, the average write amplification over a period of time.

The assessment module **330** may provide an assessment result that indicates whether at least one of the characteristic values and/or the statistical values satisfies one or more target values as specified by the host. At a decision block **730**, if the assessment result is satisfactory, the control is passed to the operation **710**. Otherwise, at operations **740**, the adjustment module **340** of FIG. **3**, for example, may be used to adjust a characteristic of a data transfer based on the assessment

result. The algorithm may be implemented by the adjustment module **340** of FIG. **3**, as discussed above. For example, adjustment module **340** may operate to adjust a characteristic such as write bandwidth, write amplification, and the number of storage units.

The adjustment module **340** of FIG. **3** may perform this operation. The adjustment module **340**, as discussed above, may operate to change one or more characteristic of a data transfer or keep the characteristic unchanged based on the assessment result. For example, when the number of storage units allocated to service a data write request is proper (determined when the entered states are one of the desired or acceptable states **610** or **620**, respectively), the adjustment module **340** may keep that number of allocated storage units unchanged. In case where, the number of storage units allocated to the data write request is not proper, for example, as indicated by a write BW smaller that BW0, the adjustment module **340** may remedy the situation by allocating a larger number of storage units to improve the write BW.

FIG. **8** is a diagram illustrating an example system **800** for managing SSD data writes, according to various embodiments of the invention. The system **800** may include a processor **810**, a memory **820**, a memory controller **830**, a graphic controller **840**, and an input and output (I/O) controller **850**, a display **852**, a keyboard **854**, a pointing device **856**, and a peripheral device **858**. A bus **860** couples all of these devices together. A clock generator **870** provides a clock signal to at least one of the devices of system **800** via bus **860**. An example of clock generator **870** may include an oscillator in a circuit board such as a motherboard. Two or more devices shown in system **800** may be formed in a single chip.

Memory **820** may comprise static random access memory (SRAM), dynamic RAM, or non-volatile memory including flash memory. Bus **860** may be interconnect traces on a circuit board or may be one or more cables. Bus **860** may also couple the devices of system **800** by wireless means such as by electromagnetic radiations, for example, radio waves. Peripheral device **858** may comprise a printer, a disk drive unit (e.g., an optical device such as a CD-ROM and a DVD reader and writer, a magnetic device reader and writer such as a floppy disk driver), or an audio device such as a microphone.

System **800** represented by FIG. **8** may include computers (e.g., desktops, laptops, hand-helds, servers, Web appliances, routers, etc.), wireless communication devices (e.g., cellular phones, cordless phones, pagers, personal digital assistants, etc.), computer-related peripherals (e.g., printers, scanners, monitors, etc.), entertainment devices (e.g., televisions, radios, stereos, tape and compact disc players, video cassette recorders/players, camcorders, digital cameras, MP3 (Motion Picture Experts Group, Audio Layer 3) players/recorders, video games, watches, etc.), and the like.

In example embodiments, the peripheral device **858** may include a machine-readable medium on which is stored one or more sets of instructions (e.g., software) embodying any one or more of the methodologies or functions described herein. The instructions may also reside, completely or at least partially, within the memory **820** and/or within the processor **810** during execution thereof by the computer system **800**, with the memory **820** and the processor **810** also constituting machine-readable media.

While the machine-readable medium is shown in an example embodiment to be a single medium, the term "machine-readable medium" should be taken to include a single medium or multiple media (e.g., a centralized or distributed database, and/or associated caches and servers) that store the one or more sets of instructions. The term "machine-readable medium" shall also be taken to include any medium that is capable of storing, encoding, or carrying a set of instructions for execution by the machine and that cause the machine to perform any one or more of the methodologies of the present invention. The term "machine-readable medium" shall accordingly be taken to include, but not be limited to, solid-state memories and optical and magnetic media.

Although specific methods and systems for managing SSD data write operations have been described, it will be evident that various modifications and changes may be made to these embodiments. Accordingly, the specification and drawings are to be regarded in an illustrative rather than a restrictive sense.

The Abstract of the Disclosure is provided to comply with 37 C.F.R. §1.72(b), requiring an abstract that allows the reader to quickly ascertain the nature of the technical disclosure. It is submitted with the understanding that it will not be used to interpret or limit the claims. In addition, in the foregoing Detailed Description, it may be seen that various features are grouped together in a single embodiment for the purpose of streamlining the disclosure. This method of disclosure is not to be interpreted as limiting the claims. Thus, the following claims are hereby incorporated into the Detailed Description, with each claim standing on its own as a separate embodiment.

What is claimed is:

1. A device comprising:

a first module to monitor a data transfer between a host and a memory;

a second module to assess the data transfer to provide at least one of a first assessment result and a second assessment result; and

a third module to adjust a number of storage units allocated to service an additional data transfer between the host and the memory based at least in part on at least one of the first assessment result and the second assessment result and based on at least one of size of requested data associated with the additional data transfer, a host-specified data transfer bandwidth, and a range of write amplifications, wherein the third module adjusts the number of storage units such that the number of storage units remains unchanged if the second module provides the first assessment result, and such that the number of storage units changes to a different number of storage units if the second module provides the second assessment result.

2. The device of claim **1**, wherein the first module is to monitor a characteristic of the data transfer, the characteristic including at least one of a size of the data transfer, a write amplification, a write bandwidth, a reclamation bandwidth, and a read bandwidth.

3. The device of claim **2**, wherein the first module is to store a value associated with the data transfer over a time period.

4. The device of claim **3**, wherein the second module is to retrieve the value to assess the data transfer.

5. The device of claim **3**, wherein the second module is to determine a statistical value based on the value associated with the data transfer.

6. The device of claim **5**, wherein the second module is to determine whether a statistical value satisfies a target value.

7. The device of claim **3**, wherein the second module is to adjust the number of storage units in a number of transitions to achieve a target value.

8. The device of claim **3**, wherein the storage units comprise a block of memory.

9. The device of claim **1**, wherein the second module comprises at least one processor.

10. The device of claim 1, wherein the device comprises a controller to control the memory.

11. A device comprising:

a first module to monitor a data transfer between a host and a memory;

a second module to compare a value associated with the data transfer with a target value with to provide at least one of a first result and a second result; and

a third module to adjust a number of storage units of the memory that are allocated to service an additional data transfer based on at least one of the first result and the second result and based on at least one of a size of requested data associated with the additional data transfer, a host-specified data transfer bandwidth and a range of write amplifications, wherein the third module adjusts the number of storage units such that the number of storage units remains unchanged if the second module provides the first assessment result, and such that the number of storage units changes to a different number of storage units if the second module provides the second assessment result.

12. The device of claim 11, wherein the data transfer comprises writing into the memory.

13. The device of claim 11, wherein the data transfer comprises reading from the memory.

14. The device of claim 11, wherein second module is to assess a characteristic of the data transfer, wherein the characteristic comprises at least one of a size of the data transfer, a write amplification, a write bandwidth, a reclamation bandwidth, and a read bandwidth.

15. The device of claim 14, wherein the first monitor is to store a value associated with the characteristic of the data transfer.

16. The device of claim 15, wherein the second module is to determine a statistical value based at least in part on the value associated with the characteristic of the data transfer.

17. A non-transitory machine readable medium storing instructions, which when executed by one or more processors perform the following operations:

monitoring a data transfer between a host and a memory;

comparing a value associated with the data transfer with a target value with to provide at least one of a first result and a second result; and

adjusting a number of storage units of the memory that are allocated to service an additional data transfer based on at least on one of the first result and the second result and based on at least one of a size of requested data associated with the additional data transfer, a host-specified data transfer bandwidth and a range of write amplifications, wherein adjusting the number of storage units includes keeping the number of storage units unchanged if the comparing provides the first result, and changing the number of storage units to a different number of storage units if the comparing provides the second result.

18. The non-transitory machine readable medium 17, wherein monitoring comprises monitoring a characteristic of the data transfer, the characteristic including at least one of a size of the data transfer, a write amplification, and a write bandwidth.

19. The non-transitory machine readable medium 18, wherein monitoring comprises monitoring at least one of a reclamation bandwidth and a read bandwidth.

20. The non-transitory machine readable medium 17, wherein the data transfer comprises at least one of writing data into the memory and reading data from the memory.

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