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(54) ADAPTIVE DELAY OF TIMING CONTROL SIGNALS

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(57)ABSTRACT

Embodiments of the present invention relate to a circuit that varies the delay time of a clock signal in response to a frequency of the clock signal. The circuit may include a sensor and a delay circuit. The sensor may output a determination of the frequency of the clock signal. The delay time of the delay circuit may be responsive to the determination of a frequency of the clock signal. Accordingly, when the frequency of a clock signal varies, a delay of that clock signal can be varied to accommodate for the change in frequency.











FIG. 5



FIG. 6











FIG. 9









ADAPTIVE DELAY OF TIMING CONTROL SIGNALS

[0001] This application is a continuation of U.S. patent application Ser. No. 10/402,137 filed Mar. 31, 2003.

BACKGROUND OF THE INVENTION

[0002] 1. Field of the Invention

[0003] The field of the invention generally relates to electronics.

[0004] 2. Background of the Related Art

[0005] Electronics are very important in the lives of many people. In fact, electronics are present in almost all electrical devices (e.g., radios, televisions, toasters, and computers). Many times electronics are virtually invisible to a user because they can be made up of very small devices inside a case. Although electronics may not be readily visible, they can be very complicated. It may be desirable in many devices for electronics to operate at fast speeds. This may be important, as many devices utilize electronics to process data or information. Accordingly, the faster a device can operate, the more data or information that can be processed in a given amount of time. For example, a DVD player processes data at a faster rate than a CD player, because reproducing moving images stored on a DVD is more intensive than reproducing audio stored on a CD.

BRIEF DESCRIPTION OF THE DRAWINGS

[0006] FIG. 1 is an exemplary global diagram of a portion of a computer.

[0007] FIG. 2 is an exemplary diagram illustrating a sensor and a delay.

[0008] FIGS. 3 and 4 are exemplary diagrams illustrating dependent logic operation circuits implemented in a single clock cycle.

[0009] FIG. 5 is an exemplary illustration of a sensor.

[0010] FIG. 6 is an exemplary block diagram of a delay.

[0011] FIG. 7 is an exemplary illustration of delay circuitry.

[0012] FIGS. 8-12 illustrate exemplary embodiments of the present invention.

DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

[0013] Electrical hardware (e.g., a computer) may include many electrical devices. In fact, a computer may include thousands of electrical devices (e.g., transistors, resistors, and capacitors). These electrical devices must work together in order for hardware to operate correctly. Accordingly, electrical devices of hardware may be electrically coupled together. This coupling may be either direct coupling (e.g., direct electrical connection) or indirect coupling (e.g., electrical communication through a series of components).

[0014] FIG. 1 is an exemplary global illustration of a computer. The computer may include processor 4, which acts as a brain of the computer. Processor 4 may be formed on a die. Processor 4 may include an Arithmetic Logic Unit (ALU) 8 and may be included on the same die as processor

4. ALU 8 may be able to perform continuous calculations in order for processor 4 to operate. Processor 4 may include cache memory 6, which may be for temporarily storing information. Cache memory 6 may be included on the same die as processor 4. The information stored in cache memory 6 may be readily available to ALU 8 for performing calculations. A computer may also include an external cache memory 2 to supplement internal cache memory 6. Power supply 7 may be provided to supply energy to processor 4 and other components of a computer. A computer may include chip set 12 coupled to processor 4. Chip set 12 may intermediately couple processor 4 to other components of the computer (e.g., graphical interface 10, Random Access Memory (RAM) 14, and/or a network interface 16). One exemplary purpose of chip set 12 is to manage communication between processor 4 and these other components. For example, graphical interface 10, RAM 14, and/or network interface 16 may be coupled to chip set 12.

[0015] FIG. 2 is an exemplary block diagram illustrating how a clock signal can be dynamically delayed. A clock signal, which is to be delayed, may be input into both sensor 18 and delay 20. Sensor 18 may be configured to sense the frequency of the input clock signal and/or the operating conditions of a circuit. The output of sensor 18 may be input into delay 20. The clock signal input into delay 20 may be delayed by delay 20 according to the output of sensor 18. The output of sensor 18 may be a signal indicative of a sensed frequency of the clock signal. The signal output from sensor 18 may be utilized by delay 20 to vary the amount of delay of the delay clock signal. In other words, in accordance with at least one embodiment of the present invention, a delay time of a clock signal may be varied in response to a sensed frequency of the clock signal.

[0016] At least one embodiment of the present invention may be implemented in a microprocessor. In microprocessors, a large amount of logic operations may be performed during each clock cycle of a clock signal. Some of these logic operations are dependent. In accordance with at least one embodiment, a logic operation may be implemented as a circuit. In the case when the circuit is an electrical circuit, a logic operation may be implemented using electrical components, such as transistors, resistors, capacitors, and inductors. In each logic operation circuit, the logic operation may be triggered by a clock signal. A clock signal is typically a periodic signal. A clock cycle may be one period of the periodic clock signal. In every clock cycle, a logic operation may be performed in a logic operation circuit. In other words, each clock cycle may drive a logic operation circuit.

[0017] FIG. 3 is an exemplary illustration of an implementation of two logic functions that are dependent. For example, during a first clock cycle, logic operation circuit 33 (logic 33) may receive a Data Input. During the first clock cycle, logic 33 may perform a logical function of the Input Data and then output the result of that logical function to logic 17. During a second clock cycle, which is subsequent to the first clock cycle, logic 17 may perform a logical function of the output from logic 17. Accordingly, the input to logic 17 is dependent on the output from logic 33. At least for this reason, logic 33 and logic 17 cannot be triggered at the same time. This is evident and apparent, as logic 33 must perform its logic operation prior to logic 17 are triggered at the

same time by a clock signal, logic **17** will not be able to utilize the output from logic **33**. Accordingly, logic **33** and logic **17** perform their logical operations during different clock cycles.

[0018] FIG. 4 is an exemplary illustration of the implementation of two logic operations performed in a single clock cycle. In at least one embodiment, logic 50 and logic 52 of FIG. 4 may have the same logical operations as logic 33 and logic 17 of FIG. 3, respectively. In other words, logic 50 may perform the same logic operation as logic 33 and logic 52 may perform the same logic operation as logic 17. However, logic 50 and logic 52 may be implemented in a single clock cycle by utilizing delay 62. Assuming that the time period of a clock cycle applied in FIG. 3 and FIG. 4 are identical, the cumulative logic operations of logic 50 and logic 52 of FIG. 4 may be processed in half the time as the cumulative logic operations of logic 33 and logic 17 of FIG. 3. In other words, the circuit structure illustrated in FIG. 4 takes only one clock cycle while the circuit structure implemented in FIG. 3 takes two clock cycles. Accordingly, the circuit structure implemented in FIG. 4 would take half the time as the circuit structure illustrated in FIG. 3 to perform the same sequence of logical operations. One of ordinary skill in the art would appreciate that implementing more dependent logic operations in the same clock cycle is advantageous, as more logic operations can be accomplished in a shorter amount of time. One of ordinary skill in the art would also appreciate that more than two dependent logic operations could be implemented in a single clock cycle.

[0019] In one or more embodiments of the present invention, there is a race between the clock signal that drives logic 50 and the delayed clock signal that drives logic 52. The delay between the clock signal that drives logic 50 and the delayed clock signal that drives logic 52 may be referred to as a race margin. The race margin is preferably long enough, such that the operation of logic 50 is completed before the operation of logic 52 begins. However, the race margin may be short enough such that both the operation of logic 50 and the operation of logic 52 are completed in the same clock cycle.

[0020] During implementation of multiple dependent logical operations in a single clock cycle, it may be desirable to minimize the delay of the clock signal between dependent logic circuits. For example, in FIG. 4, it may be desirable for delay 62 to delay the clock signal just enough such that when logic 50 outputs data, logic 52 is triggered. However, if a clock signal is not delayed enough, logic 52 would be triggered before receiving the output from logic 50. In this circumstance, the logic of the circuitry would fail. Accordingly, if delay 62 is a static delay device, then the delay must be adequate such that a failure will not occur. Further, the delay between the performance of logic 52 and logic 50 may be minimized so that multiple logic operations can be performed in a single clock cycle. One of ordinary skill in the art would appreciate that the total time consumed by logic 50, delay 62, and logic 52 cannot exceed the time period of a single clock cycle of the clock signal.

[0021] There may be variances in the frequency of a clock signal. In accordance with at least one embodiment, delay 62 must accommodate for these variances. For example, when the frequency of a clock signal is relatively high, the total time consumed by logic 50, delay 62, and logic 52 must not

exceed the time period of a single clock cycle at this relatively high frequency. Accordingly, in order to satisfy this tolerance, a relatively short static delay may be implemented. However, this relatively short static delay may not be adequate, if the frequency of the clock signal shifts to a lower frequency during normal variation of the clock frequency. Accordingly, dynamic delay circuitry, as illustrated in embodiments exemplified in **FIG. 2**, may be implemented to allow dependent logic to be implemented, regardless of normal variances in clock frequency. Accordingly, a circuit with more than one logic operation implemented in the same clock cycle will be relatively efficient and effective.

[0022] FIG. 5 is an exemplary illustration of a sensor for sensing a frequency of a clock signal. In at least one embodiment, sensor 19 of FIG. 5 may be implemented as sensor 18 of FIG. 2, and sensor 19 may be a clock chopper circuit. Clock divider 22 may reduce the frequency of the input clock signal by a factor N. In at least one embodiment, N may be two and the output of divider 22 will be a clock signal with half the frequency of the input clock signal. AND gate 24, AND gate 26, and AND gate 28 may have both a delay attribute and a logical function attribute. For example, the input into AND gate 24 may be the output of divider 22 and the clock signal. The output of AND gate 24 may be a delay of the input output from divider 22 if the clock signal is at a high level. AND gate 26 may receive the output of AND gate 24 and the clock signal. Likewise, AND gate 28 may receive the output of AND gate 26 and the clock signal. Accordingly, the outputs of AND gate 24, AND gate 26, and AND gate 28 are outputs 30a, 30b, and 30c. Accordingly, the number of consecutive lines 30a, 30b, and 30c which have a high level voltage signal may be indicative of a frequency of the clock signal. In embodiments, the output over line 30a, line 30b, and line 30c is a plurality of binary digits. In embodiments, the output over line 30a, line 30b, and line **30***c* are delay control bits.

[0023] Sensor **19** of **FIG. 5** illustrates only three AND gates for exemplary purposes and simplification. One of ordinary skill in the art would appreciate that any number of AND gates could be used depending on the required resolution and according to design preferences. Additionally, one of ordinary skill in the art would appreciate that frequency divider **22** can be divided by any multiple to accommodate for a required resolution according to design preferences.

[0024] FIG. 6 is an exemplary illustration of a delay. Delay 21 of FIG. 6 may be implemented as delay 20 of FIG. 2. Input into delay 21 are lines 31*a*, 31*b*, 31*c* and a clock signal. In accordance with at least one embodiment, lines 31*a*, 31*b*, and 31*c* may be the output of sensor 18 of FIG. 2. In accordance with at least one embodiment, these lines may be the output of sensor 19 of FIG. 5. Line 31*a* of FIG. 6 may correspond to line 30*a* of FIG. 5, line 31*b* of FIG. 6 may correspond to line 30*b* of FIG. 5, and line 31*c* of FIG. 6 may correspond to line 30*c* of FIG. 5. The clock signal input into delay 21 may be input into delay circuitry 34. According to the parameters of delay circuitry 34, delay circuitry 34 may output a delayed version of the input clock signal. Delay 21 may be a time delay register, although other circuit elements are possible.

[0025] Translation table 32 may receive signals from line 31*a*, line 31*b*, and line 31*c*. Translation table 32 may then output at least one signal 36 to delay circuitry 34. The at

least one signal 36 may control the amount of delay of delay circuitry 34. In other words, the delayed clock signal output from delay circuitry 34 may be delayed by an amount according to the plurality of signals 36 output from translation table 32. Translation table 32 may delay a clock signal by correlating a predetermined time delay as a function of a sensed frequency of the clock signal. In at least one embodiment, this function is a linear relationship between the frequency of a clock signal and the delay of the clock signal. Translation table 32 may be constructed either empirically or theoretically and may comprise logical circuitry. The logical circuitry may correlate the output of a frequency sensor to control signals that drive delay circuitry 34. Control circuitry 34 are driven according to a predetermined relationship between the frequency of a clock signal and the necessary amount of delay of the clock signal.

[0026] Only three lines [e.g., line 31a, line 31b, and line 31c] are illustrated for simplification. However, one of ordinary skill in the art would appreciate that more than three lines may be implemented. The amount of sensor lines may be equivalent to the number of outputs from sensor 19. One of ordinary skill in the art would appreciate the at least one signal 36 output from translation table 32 may be a plurality of parallel wire lines or a single wire line communicating a signal. The at least one signal 36 may be the signal that directly controls the delay of delay circuitry 34.

[0027] FIG. 7 is an exemplary illustration of delay circuitry. This circuitry may correspond to delay circuitry 34 of FIG. 6. Delay circuitry 35 may receive line 37a, line 37b, line 37c, and a clock signal. Line 37a, line 37b, and line 37c may correspond to the at least one signal 36 of FIG. 6. Capacitor 46 may be connected between node 48 and ground. Capacitor 46 may be precharged by power supply voltage Vcc through transistor 42, when the clock signal is at a low voltage level. Transistor 42 may have an inverted gate input such that Vcc is connected to capacitor 46 when the clock signal is at a low voltage level. When the clock signal is at a high voltage level, transistors 40a, 40b, and 40c may be in a highly conductive state. Likewise, when input 37a is at a high level, transistor 38a is in a highly conductive state; when input 37b is at a high voltage level, transistor 38b is in a highly conductive state; and when input 37c is at a high voltage level, transistor 38c is in a highly conductive state.

[0028] When the clock signal is at a high voltage level during each clock cycle, node 48 of capacitor 46 is connected to ground through a plurality of electrically resistant paths. The number of resistant paths which are connected to node 48 is depending on how many of line 37a, line 37b, and line 37c are at a high voltage level. Accordingly, the number of these lines which are at a high voltage level at a given time, determines how fast capacitor 46 discharges. The speed at which capacitor 46 is discharged effects the output of buffer 44. Consequently, the output of buffer 44 is a delayed version of the input clock signal.

[0029] In accordance with at least one embodiment, the number of lines (e.g., line 37a, line 37b, and line 37c) which are at a high voltage level is dependent on the frequency of the clock signal detected at the sensor 18. Transistor 38a and 40a may form cell 41a. Likewise, transistors 38b and 40b may form cell 41b. Further, transistors 38c and 40c may form cell 41c. Only three exemplary cells are illustrated for

simplification. However, one of ordinary skill in the art would appreciate that a plurality of cells may be implemented. In some embodiments, 64 cells are implemented. For each cell, a corresponding parallel line may be connected from a translation table. The number of cells which are activated during a given cycle may have a relationship to the delay of the clock signal output from delay circuitry **35**.

[0030] Embodiments of the present invention relate to races in synchronous systems that occurs when two signals start propagating from one common point and the functionality of receiving circuitry requires that one signal become valid before the other signal. If this requirement is violated, then there may be a functional failure. This failure may not be able to be resolved by reducing the clock frequency or by moving to faster operating conditions, since the two signals are starting to propagate from the same edge of the clock.

[0031] Races in high-speed digital circuits may be implemented to achieve improved performance. Races may be utilized to improve performance when the gap between the valid-time of two signals is smaller than clock phase time. Accordingly, converting a race to a synchronous path may make use of an unused timing window. An example of a race in a high-speed CPU is a race between a cache senseamplifier enabling signal and a cache memory array wordline. Another example is a race between activation of an AND-plane and evaluation of an OR-plane in dynamic Programmable-Logic-Arrays.

[0032] Embodiments of the present invention relate to a race in high speed embedded cache memories between a Sense-Amp Enable (SAE) signal and a memory array Word-Line (WL). SAE activation may occur after WL assertion in order to develop a minimum required differential signal between memory array bit-lines. When SAE and WL are activated by the same edge of the clock, there may be a functional race.

[0033] In embodiments, circuitry generates a delay (Dsae) between WL and SAE that may behave like a synchronous path. Equation 1 illustrates an exemplary required function of Dsae verses the clock cycle-time (Tcyc). Equation 1 is representation of an exemplary expression for the delay between two signals that are generated by different edges of the clock.

Equation (1):

[0034] In regular synchronous paths, K=1 if two signals are from two adjacent cycles. Likewise K=0.5 if two signals are from two adjacent phases. In embodiments where a synchronous solution for two signals is within the same phase, K may be less than 0.5.

[0035] Embodiments of the present invention may include a Speed Estimator (SPE) 51 and a Linear Time-Delay-Register (LTDR) 53 as the two basic compounds illustrated in FIG. 8. SPE 51 and LTDR 53 may include CMOS devices. The SPE 51 may generate a fix width vector (S) in which the number of bits that are at a high state (S1) is linearly dependent on the clock frequency at given operating conditions. Equation 2 is an exemplary expression for S1, where Dtap is the delay of one tap in SPE 51.

S1=Tcyc/Dtap

Dsae=D0+K*Tcvc

Equation (2):

[0036] LTDR 53 may be a programmable delay element, wherein the delay (Dtdr) is linearly controlled by a digital

word. Equation 3 is an exemplary representation of the delay of LTDR **53**, where S1 is the number of bits that are at high state in the input data word. Dc is an exemplary representation of fixed delay and it may be a parameter of LTDR **53**.

Dtdr=D0+S1*Dc

Equation (3):

[0037] Accordingly, the connection of SPE 51 to LTDR 53 to output SAE, as shown in FIG. 8, can be represented by the exemplary expression of Equation 4 for the delay of SAE, after combining Equation 2 and Equation 3.

Dsae=D0+(Dc/Dtap)*Tcyc

Equation (4):

[0038] Dtap and Dc may be delay parameters of the SPE 51 and the LTDR 53. Dtap and Dc may respond to operating conditions (i.e., supply voltage, temperature, and process parameters). The ratio of Dc/Dtap may be a fixed number. Accordingly, the delay between the WL and SAE may behave like a regular synchronous path, as described by Equation 1, with K=Dc/Dtap.

[0039] FIG. 9 illustrates an exemplary implementation of SPE 51. This implementation may include a clock divider 55 that generates a half frequency clock (CLK2) with a phase time that is equal to Tcyc of the main clock. CLK2 may feed a simple chain of N delay elements 57 that are connected in a clock chopper configuration. The delay of each element 59 of N delay elements 57 is Dtap. The clock is ANDed with the outputs of each element 59 along the chain. The output of clock-AND[i] is a signal that rises i*Dtap [pS] after the clock rising edge, and falls immediately after the clock falling edge, as illustrated in FIG. 10. The taps may be sampled by falling edge of triggered flip flops 58. The outputs of these flip flops 58 may be the required vector S defined by Equation 2.

[0040] LTDR 53 may be implemented, in embodiments of the present invention, as shown in FIG. 11. TDR 61 may be similar to a programmable delay component. TDR 61 may have a set of small devices 63 that are connected through a set of sixty-four switches 65 to a capacitor 67. Sixty-four switches 65 control how many elements will discharge capacitor 67. Therefore, the delay of TDR 61 may be proportional to an inverse of the number of driving elements. Equation 5 is an exemplary representation of the delay of TDR 61, wherein N is the number elements (i.e., sixty-four) that are enabled to discharge capacitor 67.

Dtdr=K3/N

Equation (5):

[0041] In order to obtain a linear ratio between the delay of LTDR 53 and the input data word that is coming from the SPE 51, the number of driving cells may be proportional to 1/S1. This may be achieved by connecting each bit from SPE 51 to a variable number of cells in TDR 53. When S1=0, TDR 53 is at maximum speed and the number of driving cells may be at maximum value -n(0). When S1 increases to 1, the number of driving cells may be reduced by Dn(0). When S1 increases from 1 to 2, the number of driving cells may be reduced by Dn(0), since now less cells are driving so the weight of each eliminated cell is greater.

[0042] An exemplary mathematical solution of Equations 6-10 illustrate that the number of cells that should be disabled (Dn) is a function of the current number of driving cells (n) and the number of high bits in the input vector (S1).

To get a linear response, the derivative of the function Dtdr(S1) should be constant.

DDtdr/DS1=Dc Equation (6):

[0043] The following set of equations extracts the derivative Dn/DS1 (i.e., how many cells should be disconnected when S1 increases by 1).

 $DDtdr/DS1=(DDtrd/Dn)^*(Dn/DS1)$ Equation (7): $DDtdr/Dn=-K3/n^2$ (from Equation 5)tm Equation (8):

[0044] Combining Equation 6, Equation 7 and Equation 8 gives:

 $Dn/DS1=-(Dc/K3)*n^2$ Equation (9):

[0045] Therefore, if the current value of driving cells is n, then when S1 increase by 1, the next value of n is given by:

 $n(S1+1)=n(S1)B(Dc/K3)^*[n(S1)]^2$ Equation (10): [0046] Below is an example of how many cells should be disabled as S1 increases from 0. Assume that the delay coefficient K3 of the TDR is 630 pS and that the required fixed delay step Dc is 3 pS. The initial value n(0) is 41 cells.

 $n(1)=41B(3/630)*41^2=33$

n(2)=33B(3/630)*33^2=28

[0047] Table 1 is an exemplary full set of n(S1) when S1 grows from zero to twelve. The values are rounded to the closest integer number. Beyond twelve, $_n(S1)$ is smaller than 0.5 and the LTDR may become non-linear. However, the available linear range may be wide enough for practical design purposes.

TABLE 1

S1	n(S1)	$\Delta n(S1)$
0	41.0	8.0
1	33.0	5.2
2	27.8	3.7
3	24.1	2.8
4	21.4	2.2
5	19.2	1.8
6	17.4	1.4
7	16.0	1.2
8	14.8	1.0
9	13.7	0.9
10	12.8	0.8
11	12.1	0.7
12	11.4	0.6

[0048] FIG. 12 is an exemplary graph of Dsae vs. Tcyc. The delay parameters in this design are Dc=3 pS and Dtap=21 pS. Therefore, an expected slope of the curve is 0.14. It may be appreciated that SAE behaves almost like a synchronous circuit with a delay that is linearly dependent on Tcyc within a wide range of the frequency. Therefore, the race conditions between the WL and SAE may be removed.

[0049] The foregoing embodiments and advantages are merely exemplary and are not to be construed as limiting the present invention. The present teaching can be readily applied to other types of apparatuses. The description of the present invention is intended to be illustrative, and not to limit the scope of the claims. Many alternatives, modifications, and variations will be apparent to those skilled in the art.

What is claimed is:

1. An apparatus, comprising:

a circuit to apply a clock signal to a block that performs a first logical operation;

- a circuit to delay the clock signal for input into a block that performs a second logical operation, wherein the second logical operation depends on a result of the first logical operation, and wherein a race margin between the clock signal and the delayed clock signal is (a) long enough so that the first logical operation is completed before the second logical operation begins and (b) short enough so that the first and second logical operations are performed within a same clock cycle.
- 2. The apparatus of claim 1, further comprising:
- a sensor to detect a frequency of the clock signal,
- wherein the delay circuit sets the delay of the clock signal based on the frequency of the clock signal detected by the sensor.

3. The apparatus of claim 2, wherein the sensor generates a plurality of signals indicative of the frequency of the clock signal.

4. The apparatus of claim 3, wherein the delay circuit delays the clock signal based on a number of the signals at a predetermined voltage level.

5. The apparatus of claim 3, wherein the delay circuit includes:

- a translation table to store delay information for a plurality of clock frequencies, said translation table outputting a delay signal based on the signals from the sensor; and
- a combination of delay elements to delay the clock signal in response to the delay signal output from the translation table.

6. The apparatus of claim 5, wherein each of the delay elements includes:

- a first transistor having a gate coupled to the clock signal; and
- a second transistor having a gate coupled to a respective one of a plurality of signal lines that output the delay signal from the translation table, wherein voltage signals from the first transistors of the delay elements combine to charge and discharge a capacitor in order to delay the clock signal by an amount which satisfies the race margin.

7. The apparatus of claim 2, wherein the delay circuit sets the delay of the clock signal as a linear function of the frequency of the clock signal.

8. The apparatus of claim 2, wherein the delay circuit changes a delay of the clock signal to cause the race margin to continue to satisfy (a) and (b) when the sensor detects a change in the clock signal frequency.

9. A control method, comprising:

driving a first logical operation with a clock signal;

delaying the clock signal;

- driving a second logical operation with the delayed clock signal,
- wherein the second logical operation depends on a result of the first logical operation, and wherein a race margin between the clock signal and the delayed clock signal is (a) long enough so that the first logical operation is completed before the second logical operation begins and (b) short enough so that the first and second logical operations are performed within a same clock cycle.

- 10. The method of claim 9, further comprising:
- setting delay of the clock signal based on a frequency of the clock signal.

11. The method of claim 10, wherein the delay of the clock signal is set as a linear function of the frequency of the clock signal

12. The method of claim 10, further comprising:

generating a plurality of binary digits from a frequency sensor indicative of the frequency of the clock signal, wherein the delay of the clock signal is set based on a number of the binary digits which have a predetermined voltage level.

13. The method of claim 12, wherein delaying the clock signal includes:

- generating a control signal from a translation table in response to the number of binary digits at said predetermined voltage level; and
- activating at least one combination of delay elements based on the control signal.

14. The method of claim 13, wherein a relationship between a number of activated delay elements and the delay of the clock signal is a nonlinear relationship.

15. The method of claim 9, further comprising:

- detecting a change in frequency of the clock signal; and
- changing delay of the clock signal in response to the detected frequency change,
- wherein the delay change causes the race margin to continue to satisfy (a) and (b).
- 16. A system, comprising:
- a first circuit having blocks to perform first and second logical operations;
- a second circuit to apply a clock signal to the block performing the first logical operation;
- a third circuit to delay the clock signal for input into the block that performs the second logical operation, wherein the second logical operation depends on a result of the first logical operation and wherein a race margin between the clock signal and the delayed clock signal is (a) long enough so that the first logical operation is completed before the second logical operation begins and (b) short enough so that the first and second logical operations are performed within a same clock cycle.

17. The system of claim 16, wherein the first circuit is one of a processor, graphical interface, cache, arithmetic logic unit, network interface, chipset, and memory.

18. The system of claim 16, further comprising:

a sensor to detect a frequency of the clock signal,

wherein the delay circuit sets the delay of the clock signal based on the frequency of the clock signal detected by the sensor.

19. The system of claim 16, wherein the delay circuit changes a delay of the clock signal to cause the race margin to continue to satisfy (a) and (b) when the sensor detects a change in the clock signal frequency.

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